

Introduction to Parallel Programming for Multicore/Manycore Clusters

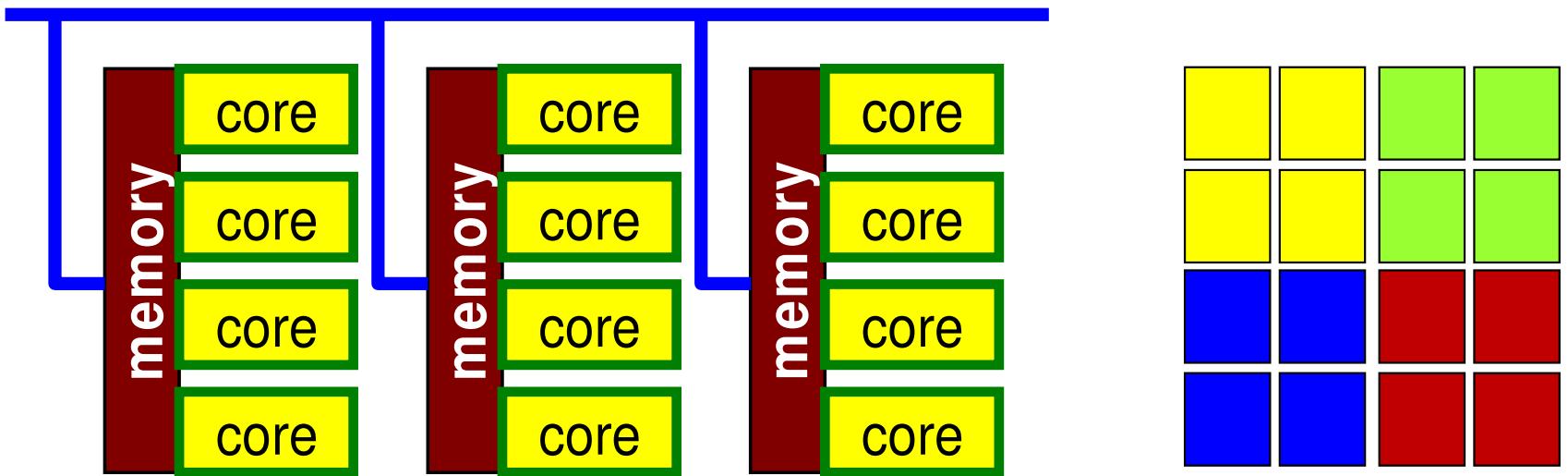
Part A2: Introduction to OpenMP

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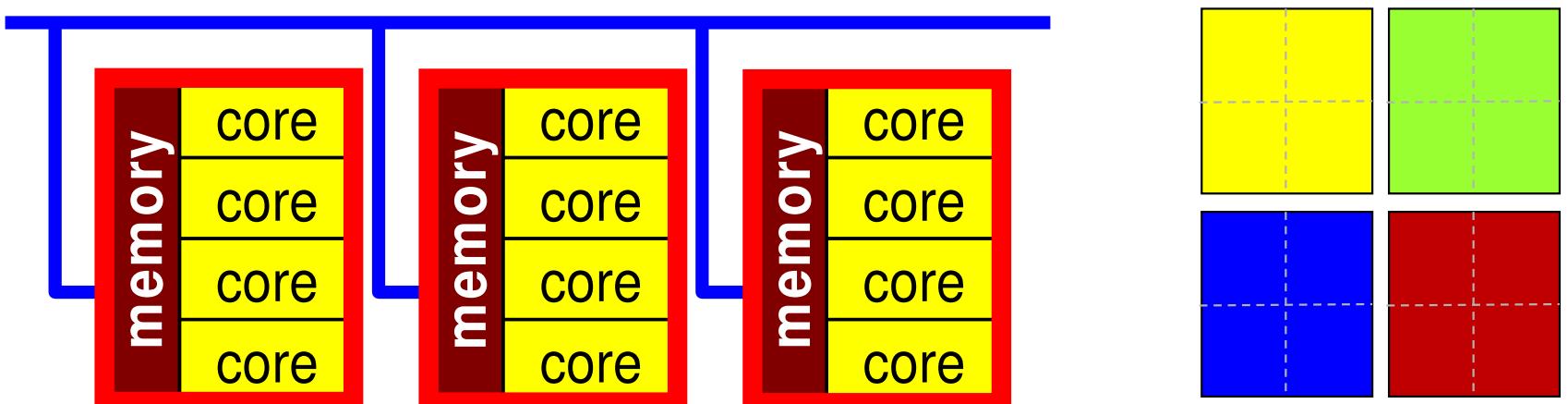
- OpenMP
- Login to Wisteria/BDEC-01
- Parallel Code by OpenMP (0): up to 12 cores
- Parallel Code by OpenMP (1): First Touch
- Parallel Code by OpenMP (2): +ELL
- Parallel Code by OpenMP (3): reduced omp-parallel
- Parallel Code by OpenMP (4): Further Optimization (Fortran only)

Flat MPI vs. Hybrid

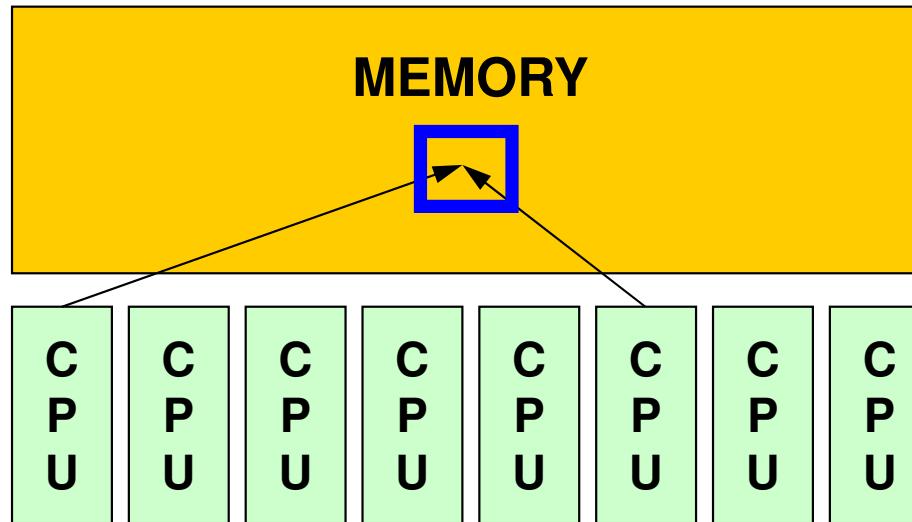
Flat-MPI: Each Core -> Independent



Hybrid : Hierarchical Structure



SMP



- **SMP**
 - Symmetric Multi Processors
 - Multiple CPU's (cores) share a single memory space

What is OpenMP ? (1/2)

<http://www.openmp.org>

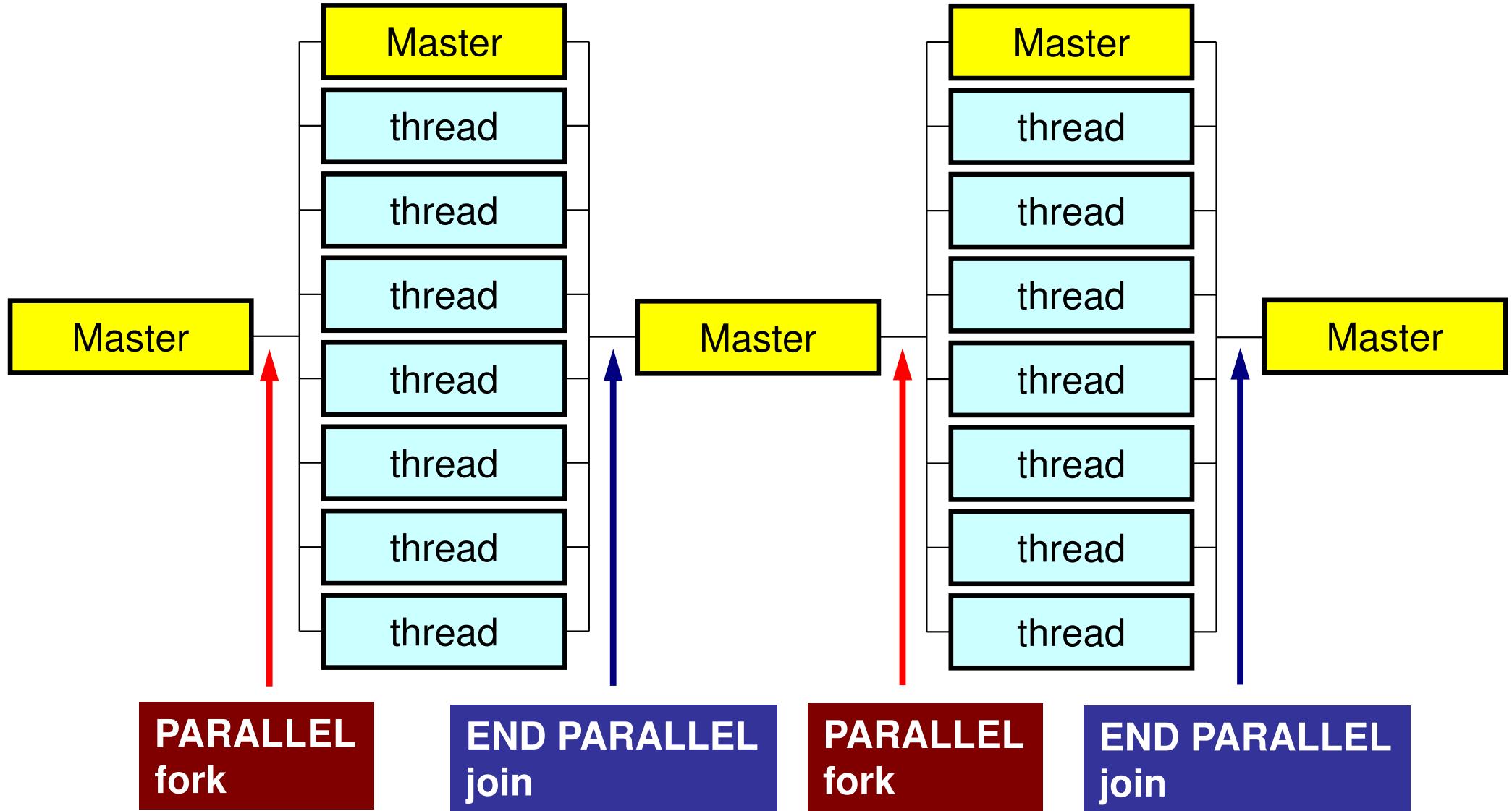
- An API (Application Programming Interface) for multi-platform shared-memory parallel programming in C/C++ and Fortran
 - Current version: 4.X (5.0 is already announced)
 - GPU, Accelerators: close to OpenACC
- Background
 - Merger of Cray and SGI in 1996
 - Separated later, ... but both are now merged into HPE
 - ASCI project (US-DOE (Dept. of Energy)) started in 1995
 - Accelerated Strategic Computing Initiative (ASCI) -> Advanced Simulation and Computing Program (ASC)
 - The goal of ASCI is to simulate the results of new weapons designs as well as the effects of aging on existing and new designs, all in the absence of additional data from underground nuclear tests.
 - Development of Supercomputers & Software/Applications
 - SMP Clusters: Intel ASCI Red, IBM Power (Blue, White, Purple)/Blue Gene, SGI
 - Common API for SMP Clusters needed

What is OpenMP ? (2/2)

<http://www.openmp.org>

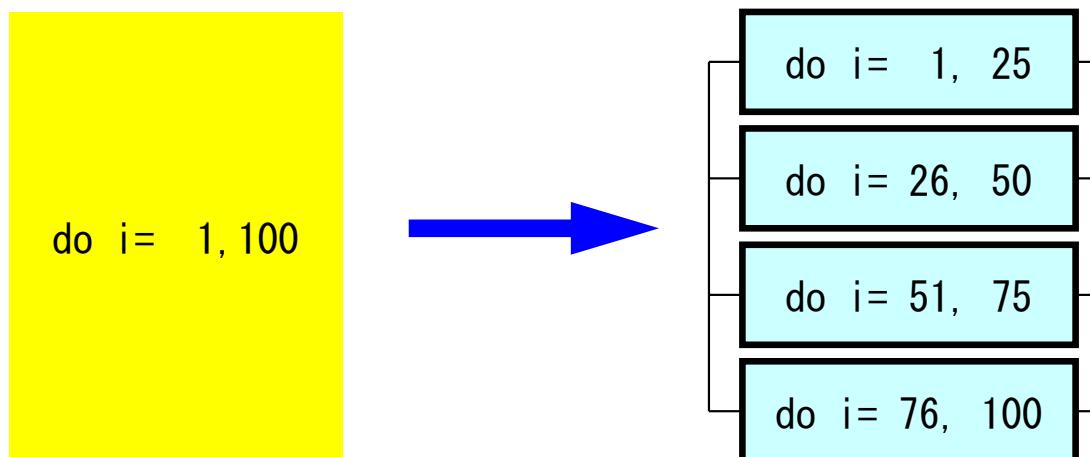
- C/C++ version and Fortran version have been separately developed until ver.2.5.
- Fork-Join Parallel Execution Model (Next Page)
 - Directives: Parallel, End Parallel
 - Serial Execution: Master Thread
 - Parallel Execution: Master Thread/Thread Team
- Users have to specify everything by directives.
 - Nothing happen, if there are no directives

Fork-Join Parallel Execution Model



Number of Threads

- **OMP_NUM_THREADS**
 - How to change ?
 - bash(.bashrc)
 - csh(.cshrc)
- **OMP_NUM_THREADS=4**



Information about OpenMP

- OpenMP Architecture Review Board (ARB)
 - <http://www.openmp.org>
 - Spec. of OpenMP is available
- References
 - Chandra, R. et al. 「Parallel Programming in OpenMP」
(Morgan Kaufmann)
 - Quinn, M.J. 「Parallel Programming in C with MPI and OpenMP」(McGrawHill)
 - Mattson, T.G. et al. 「Patterns for Parallel Programming」
(Addison Wesley)
 - 牛島「OpenMPによる並列プログラミングと数値計算法」(丸善)
 - Chapman, B. et al. 「Using OpenMP」(MIT Press)
- Japanese Version of OpenMP 3.0 Spec. (Fujitsu etc.)
 - <http://www.openmp.org/mp-documents/OpenMP30spec-ja.pdf>

Features of OpenMP

- Directives
 - Loops right after the directives are parallelized.
 - If the compiler does not support OpenMP, directives are considered as just comments.

OpenMP/Directives Array Operations

Simple Substitution

```
!$omp parallel do
  do i= 1, N
    W(i, 1)= 0. d0
    W(i, 2)= 0. d0
  enddo
 !$omp end parallel do
```

Dot Products

```
!$omp parallel do private(i)
!$omp&           reduction(+:RH0)
  do i= 1, N
    RH0= RH0 + W(i, R)*W(i, Z)
  enddo
 !$omp end parallel do
```

DAXPY

```
!$omp parallel do
  do i= 1, N
    Y(i)= ALPHA*X(i) + Y(i)
  enddo
 !$omp end parallel do
```

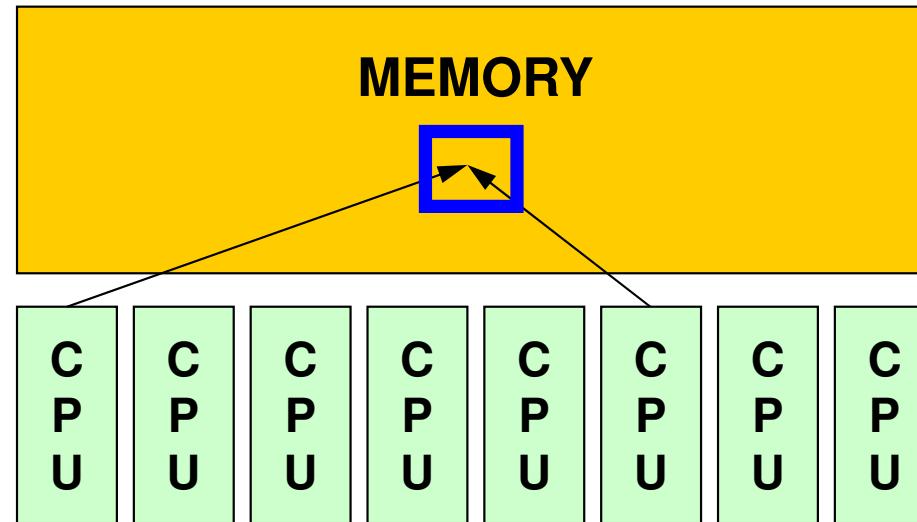
OpenMP/Direceives Matrix/Vector Products

```
!$omp parallel do private(i, k)
    do i= 1, N
        W(i, Q)= D(i)*W(i, P)
        do k= indexLU(i-1)+1, indexLU(i)
            W(i, Q)= W(i, Q) + AMAT(k)*W(itemLU(k), P)
        enddo
    enddo
 !$omp end parallel do
```

Features of OpenMP

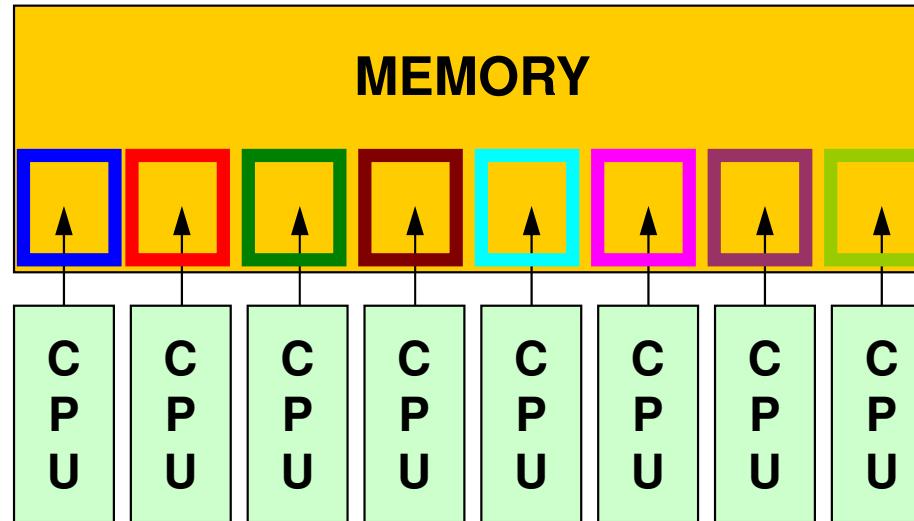
- Directives
 - Loops right after the directives are parallelized.
 - If the compiler does not support OpenMP, directives are considered as just comments.
- Nothing happen without explicit directives
 - Different from “automatic parallelization/vectorization”
 - Something wrong may happen by un-proper way of usage
 - Data configuration, ordering etc. are done under users’ responsibility
- “Threads” are created according to the number of cores on the node
 - Thread: “Process” in MPI
 - Generally, “# threads = # cores”: Xeon Phi supports 4 threads per core (Hyper Multithreading)

Memory Contention: メモリ競合



- During a complicated process, multiple threads may simultaneously try to update the data in same address on the memory.
 - e.g.: Multiple cores update a single component of an array.
 - This situation is possible.
 - Answers may change compared to serial cases with a single core (thread).

Memory Contention (cont.)



- In this lecture, any such case does not happen by reordering etc.
 - In OpenMP, users are responsible for such issues (e.g. proper data configuration, reordering etc.)
- Data Dependency
- Performance per core reduces as number of used cores (thread #) increases (Memory Saturation)

Features of OpenMP (cont.)

- “`!omp parallel do`”-“`!omp end parallel do`”
- Global (Shared) Variables, Private Variables
 - Default: Global (Shared)
 - Dot Products: reduction

```
!$omp parallel do private(i)
!$omp&           reduction(+:RHO)
do i= 1, N
    RHO= RHO + W(i, R)*W(i, Z)
enddo
 !$omp end parallel do
```

W(:,:, R, Z
global (shared)

FORTRAN & C

```
use omp_lib

```
 !$omp parallel do shared(n, x, y) private(i)
 do i= 1, n
 x(i)= x(i) + y(i)
 enddo
 !$omp end parallel do
```

```
#include <omp.h>

```
 {
    #pragma omp parallel for default(none) shared(n, x, y) private(i)

        for (i=0; i<n; i++)
            x[i] += y[i];
}
```

In this class ...

- There are many capabilities of OpenMP.
- In this class, only several functions are shown for parallelization of PCG/ICCG solver.

First things to be done

- use omp_lib Fortran
- #include <omp.h> C

OpenMP Directives (Fortran)

```
sentinel directive_name [clause[, ] clause]...]
```

- NO distinctions between upper and lower cases.
- sentinel
 - Fortran: !\$OMP, C\$OMP, *\$OMP
 - !\$OMP only for free format
 - Continuation Lines (Same rule as that of Fortran compiler is applied)
 - Example for !\$OMP PARALLEL DO SHARED (A, B, C)

```
!$OMP PARALLEL DO  
!$OMP+SHARED (A, B, C)
```

```
!$OMP PARALLEL DO &  
!$OMP SHARED (A, B, C)
```

OpenMP Directives (C)

```
#pragma omp directive_name [clause[,] clause]...
```

- “\” for continuation lines
- Only lower case (except names of variables)

```
#pragma omp parallel for shared (a,b,c)
```

PARALLEL DO/for

```
!$OMP PARALLEL DO [clause[,] clause] ... ]  
  (do_loop)  
 !$OMP END PARALLEL DO
```

```
#pragma parallel for [clause[,] clause] ... ]  
  (for_loop)
```

- Parallelize DO/for Loops
- Examples of “clause”
 - private(list)
 - shared(list)
 - default(private|shared|none)
 - reduction({operation|intrinsic}: list)

REDUCTION

```
REDUCTION ({operator|intrinsic}: list)
```

```
reduction ({operator|intrinsic}: list)
```

- Similar to “MPI_Reduce”
- Operator
 - +, *, -, .AND., .OR., .EQV., .NEQV.
- Intrinsic
 - MAX, MIN, IAND, IOR, IEQR

Example-1: A Simple Loop

```
!$OMP PARALLEL DO
    do i= 1, N
        B(i)= (A(i) + B(i)) * 0.50
    enddo
!$OMP END PARALLEL DO
```

- Default status of loop variables (“i” in this case) is private. Therefore, explicit declaration is not needed.
- “END PARALLEL DO” is not required
 - In C, there are no definitions of “end parallel do”

Example-1: REDUCTION

```
!$OMP PARALLEL DO DEFAULT(PRIVATE) REDUCTION(+:A,B)
    do i= 1, N
        call WORK (Alocal, Blocl)
        A= A + Alocal
        B= B + Blocl
    enddo
!$OMP END PARALLEL DO
```

- “END PARALLEL DO” is not required

Functions in OpenMP

functions	description
<code>int omp_get_num_threads (void)</code>	Thread #
<code>int omp_get_thread_num (void)</code>	Thread ID
<code>double omp_get_wtime (void)</code>	Timer
<code>void omp_set_num_threads (int num_threads)</code> <code>call omp_set_num_threads (num_threads)</code>	Specifying Thread #

OpenMP for Dot Products

```
VAL= 0. d0
do i= 1, N
    VAL= VAL + W(i, R) * W(i, Z)
enddo
```

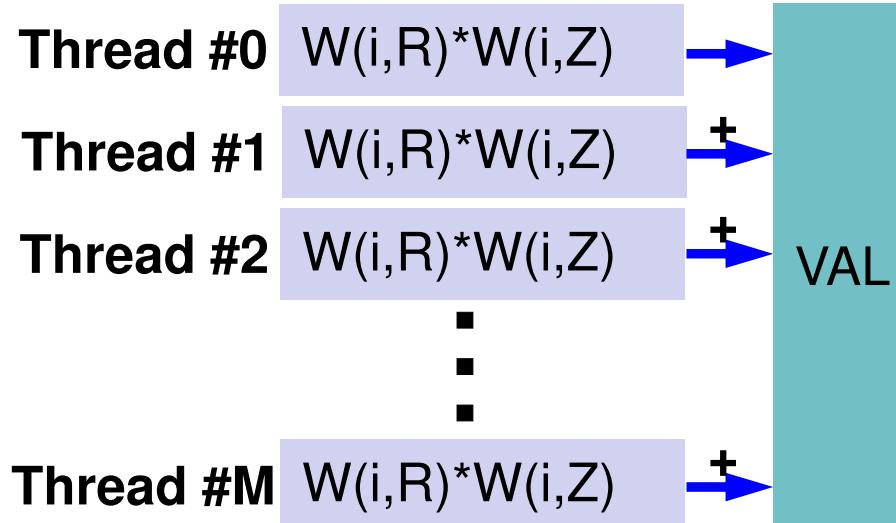
OpenMP for Dot Products

```
VAL= 0. d0
do i= 1, N
    VAL= VAL + W(i, R) * W(i, Z)
enddo
```



```
VAL= 0. d0
!$OMP PARALLEL DO PRIVATE(i) REDUCTION(+:VAL)
do i= 1, N
    VAL= VAL + W(i, R) * W(i, Z)
enddo
!$OMP END PARALLEL DO
```

Directives are just inserted.



OpenMP for Dot Products

```
VAL= 0. d0
do i= 1, N
    VAL= VAL + W(i, R) * W(i, Z)
enddo
```



```
VAL= 0. d0
!$OMP PARALLEL DO PRIVATE(i) REDUCTION(+:VAL)
do i= 1, N
    VAL= VAL + W(i, R) * W(i, Z)
enddo
!$OMP END PARALLEL DO
```



```
VAL= 0. d0
!$OMP PARALLEL DO PRIVATE(ip, i) REDUCTION(+:VAL)
do ip= 1, PEsmptOT
    do i= index(ip-1)+1, index(ip)
        VAL= VAL + W(i, R) * W(i, Z)
    enddo
enddo
!$OMP END PARALLEL DO
```

Directives are just inserted.

Multiple Loop
PEsmptOT: Number of threads

Additional array **INDEX(:)** is needed.
 Efficiency is not necessarily good, but users can specify thread for each component of data.

OpenMP for Dot Products

```

VAL= 0. d0
!$OMP PARALLEL DO PRIVATE(ip, i) REDUCTION(+:VAL)
do ip= 1, PEsmptOT
    do i= index(ip-1)+1, index(ip)
        VAL= VAL + W(i,R) * W(i,Z)
    enddo
enddo
!$OMP END PARALLEL DO

```

Multiple Loop

PEsmptOT: Number of threads

Additional array **INDEX(:)** is needed.

Efficiency is not necessarily good, but users can specify thread for each component of data.

e.g.: N=100, PEsmptOT=4

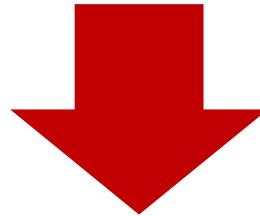
```

INDEX(0)= 0
INDEX(1)= 25
INDEX(2)= 50
INDEX(3)= 75
INDEX(4)= 100

```

Matrix-Vector Multiply

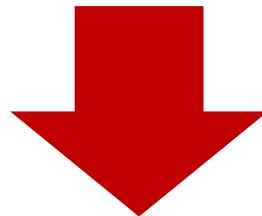
```
do i= 1, N
    W(i, Q) = D(i)*W(i, P)
    do k= indexLU(i-1)+1, indexLU(i)
        W(i, Q) = W(i, Q) + AMAT(k)*W(itemLU(k), P)
    enddo
enddo
```



```
!$omp parallel do private(i, k)
do i= 1, N
    W(i, Q) = D(i)*W(i, P)
    do k= indexLU(i-1)+1, indexLU(i)
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    enddo
enddo
 !$omp end parallel do
```

Matrix-Vector Multiply

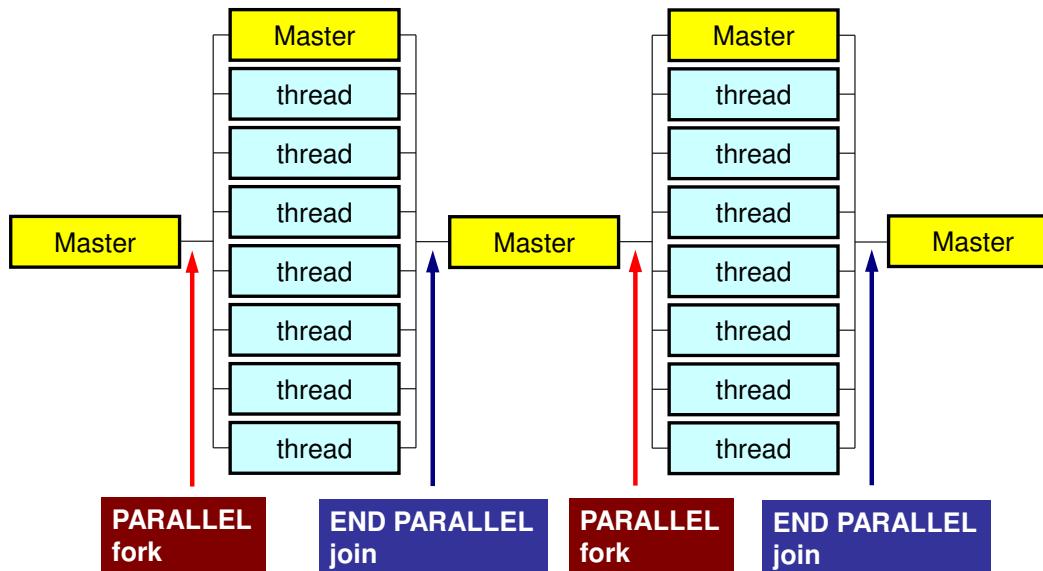
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    do k= indexLU(i-1)+1, indexLU(i)
        W(i, Q) = W(i, Q) + AMAT(k)*W(itemLU(k), P)
    enddo
enddo
```



```
!$omp parallel do private(ip, i, k)
do ip= 1, PEsmptOT
do i= index(ip-1)+1, index(ip)
    W(i, Q) = D(i)*W(i, P)
    do k= indexLU(i-1)+1, indexLU(i)
        W(i, Q) = W(i, Q) + AMAT(k)*W(itemLU(k), P)
    enddo
enddo
enddo
 !$omp end parallel do
```

omp parallel (do)

- “omp parallel-omp end parallel” = “fork-join”
- If you have many loops, these “fork-join’s” cause overheads
- omp parallel + omp do/omp for



```
#pragma omp parallel ...
```

```
#pragma omp for {
```

...

```
#pragma omp for {
```

```
!$omp parallel ...
```

```
!$omp do  
do i= 1, N
```

...

```
!$omp do  
do i= 1, N
```

...

```
!$omp end parallel required
```

- OpenMP
- **Login to Wisteria/BDEC-01**
- Parallel Code by OpenMP (0): up to 12 cores
- Parallel Code by OpenMP (1): First Touch
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Target for Parallelization

- “Parallelization” of the FVM code
- Just inserting OpenMP directives
 - “poi_gen.f/c”(poi_gen),
 - “solver_PCG.f/c” (solve_PCG)

Preconditioned Conjugate Gradient Method (PCG)

```

Compute  $r^{(0)} = b - [A]x^{(0)}$ 
for  $i = 1, 2, \dots$ 
    solve  $[M]z^{(i-1)} = r^{(i-1)}$ 
     $\rho_{i-1} = r^{(i-1)} \cdot z^{(i-1)}$ 
    if  $i=1$ 
         $p^{(1)} = z^{(0)}$ 
    else
         $\beta_{i-1} = \rho_{i-1}/\rho_{i-2}$ 
         $p^{(i)} = z^{(i-1)} + \beta_{i-1} p^{(i-1)}$ 
    endif
     $q^{(i)} = [A]p^{(i)}$ 
     $\alpha_i = \rho_{i-1}/p^{(i)}q^{(i)}$ 
     $x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$ 
     $r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$ 
    check convergence  $|r|$ 
end

```

Solving the following equation:

$$\{z\} = [M]^{-1}\{r\}$$

“Approximate Inverse Matrix”

$$[M]^{-1} \approx [A]^{-1}, \quad [M] \approx [A]$$

Ultimate Preconditioning:
Inverse Matrix

$$[M]^{-1} = [A]^{-1}, \quad [M] = [A]$$

Diagonal Scaling: Simple but weak

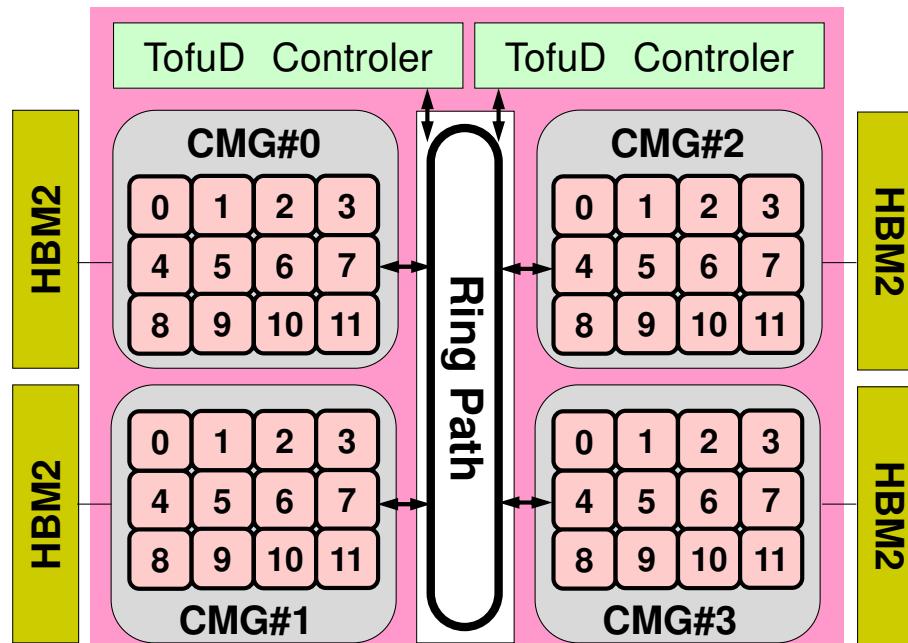
$$[M]^{-1} = [D]^{-1}, \quad [M] = [D]$$

Diagonal Scaling, Point-Jacobi

$$[M] = \begin{bmatrix} D_1 & 0 & \dots & 0 & 0 \\ 0 & D_2 & & 0 & 0 \\ \dots & & \dots & & \dots \\ 0 & 0 & & D_{N-1} & 0 \\ 0 & 0 & \dots & 0 & D_N \end{bmatrix}$$

- **solve** $[M] z^{(i-1)} = r^{(i-1)}$ is very easy.
- Provides fast convergence for simple problems.

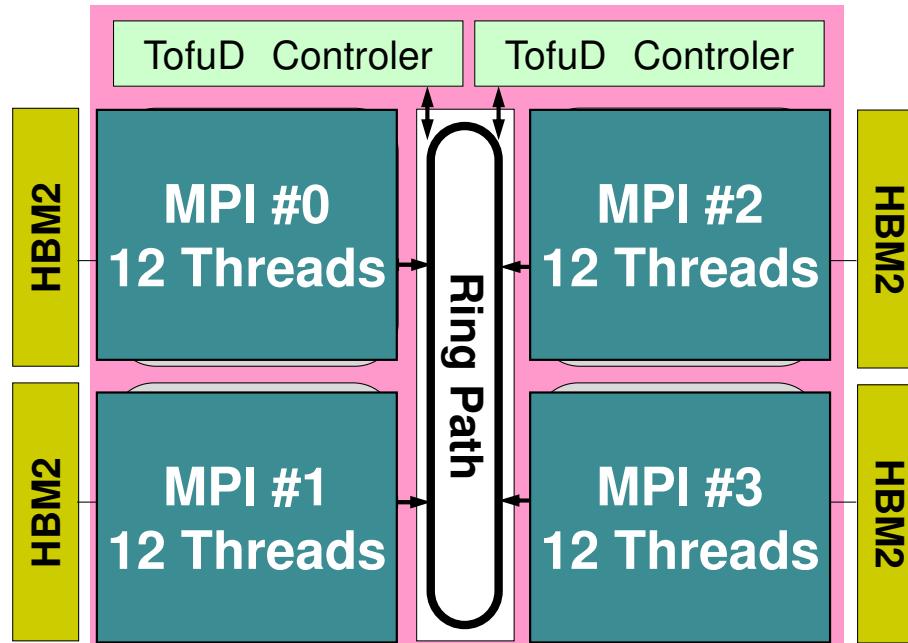
A64FX Processor on Odyssey



Name	A64FX
Processor # (Core #)	1 (48+ 2or4 Assistant Cores)
Frequency	2.2 GHz
Peak Performance	3.3792 TFLOPS
Memory Size	32 GiB
Memory Bandwidth	1,024 GB/s
L1 Cache	64 KiB/core (Inst/Data)
L2 Cache	8 MiB/CMG

- 4 CMG's (Core Memory Group), 12 cores/CMG
 - 48 Cores/Node (Processor)
 - $2.2\text{GHz} \times 32\text{DP} \times 48 = 3379.2 \text{ GFLOPS} = 3.3792 \text{ TFLOPS}$
- NUMA Architecture (Non-Uniform Memory Access)
 - Each core of a CMG can access to the memory on other CMG's
 - Utilization of the local memory is more efficient

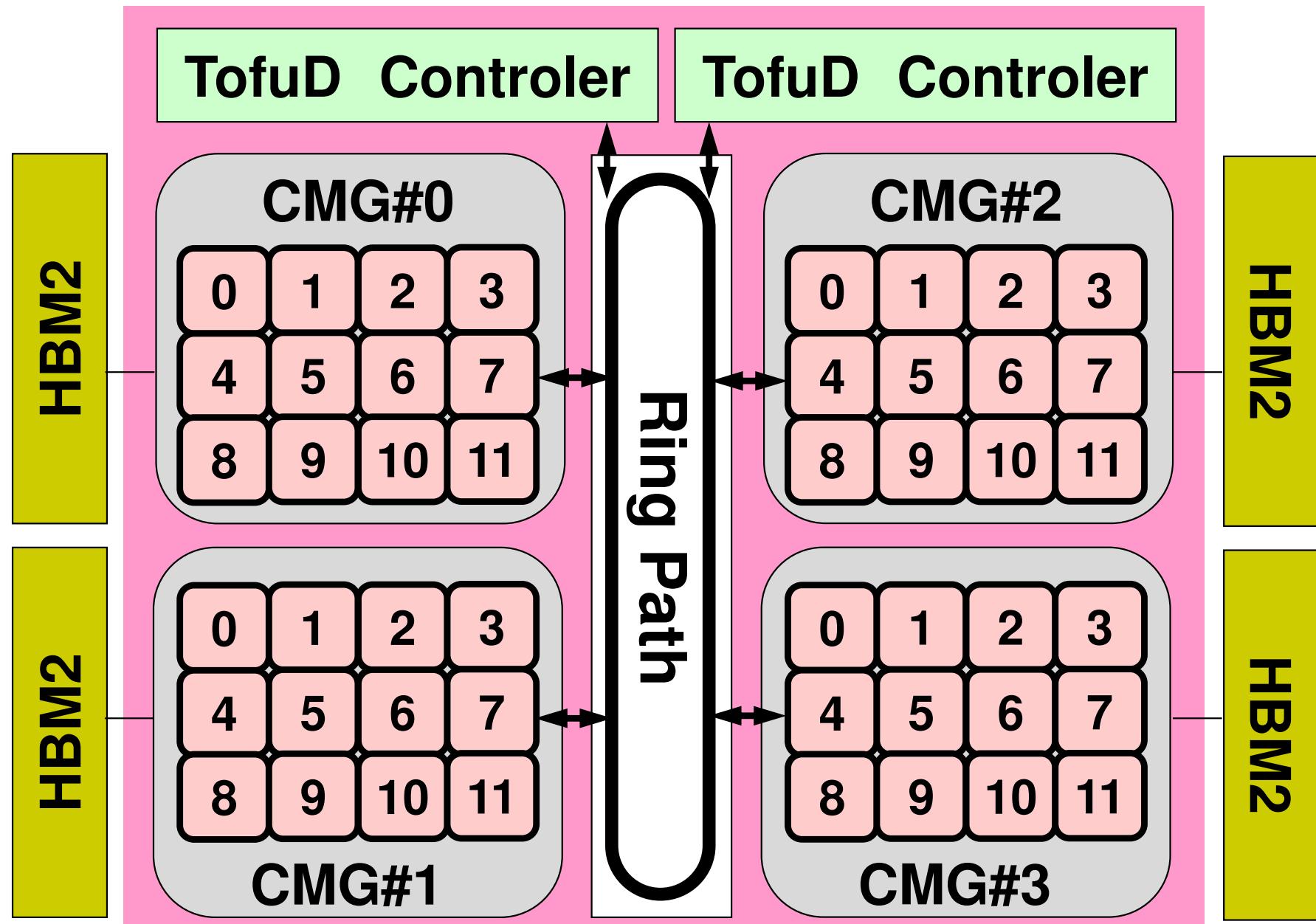
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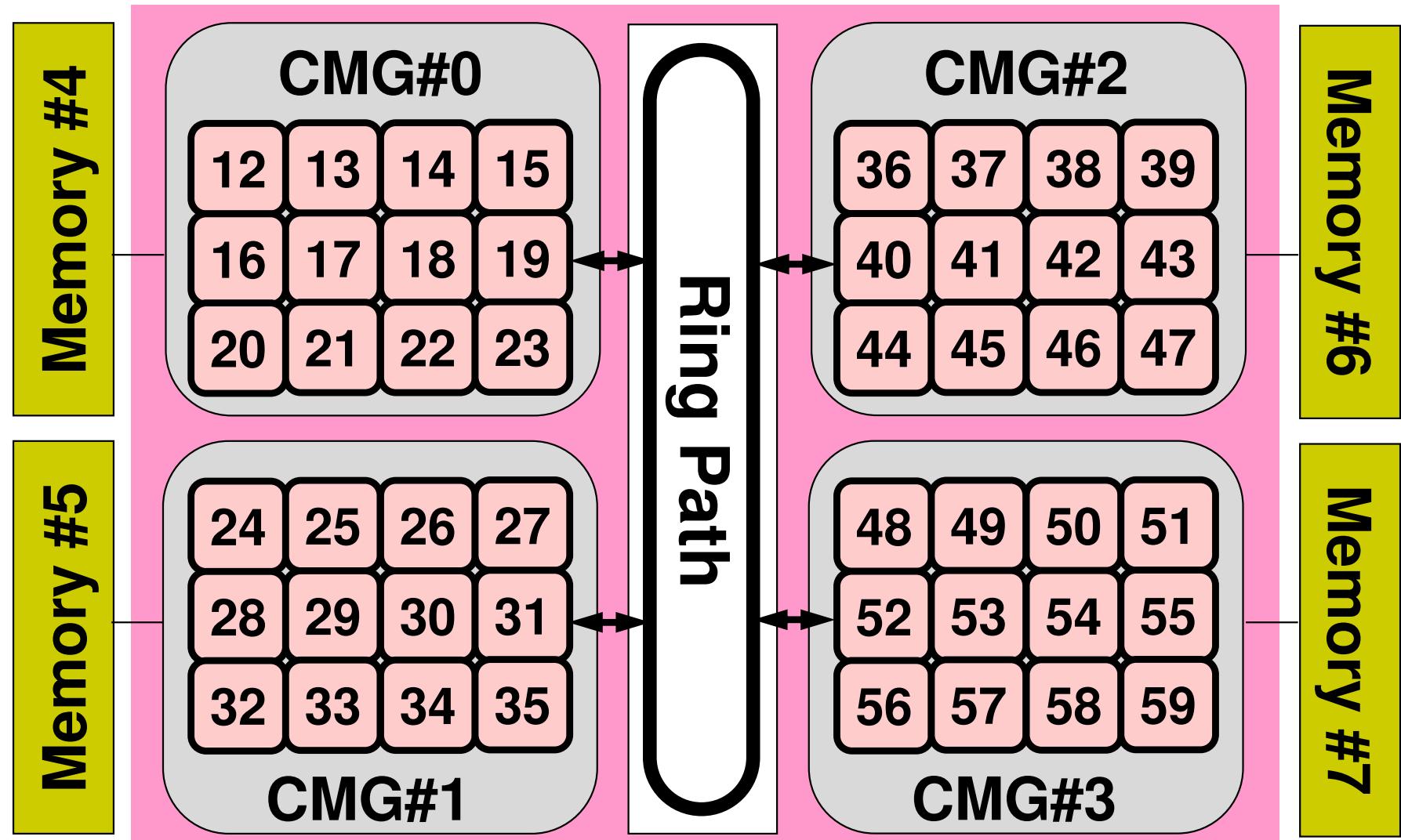
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A64FX : CMG (Core Memory Group)



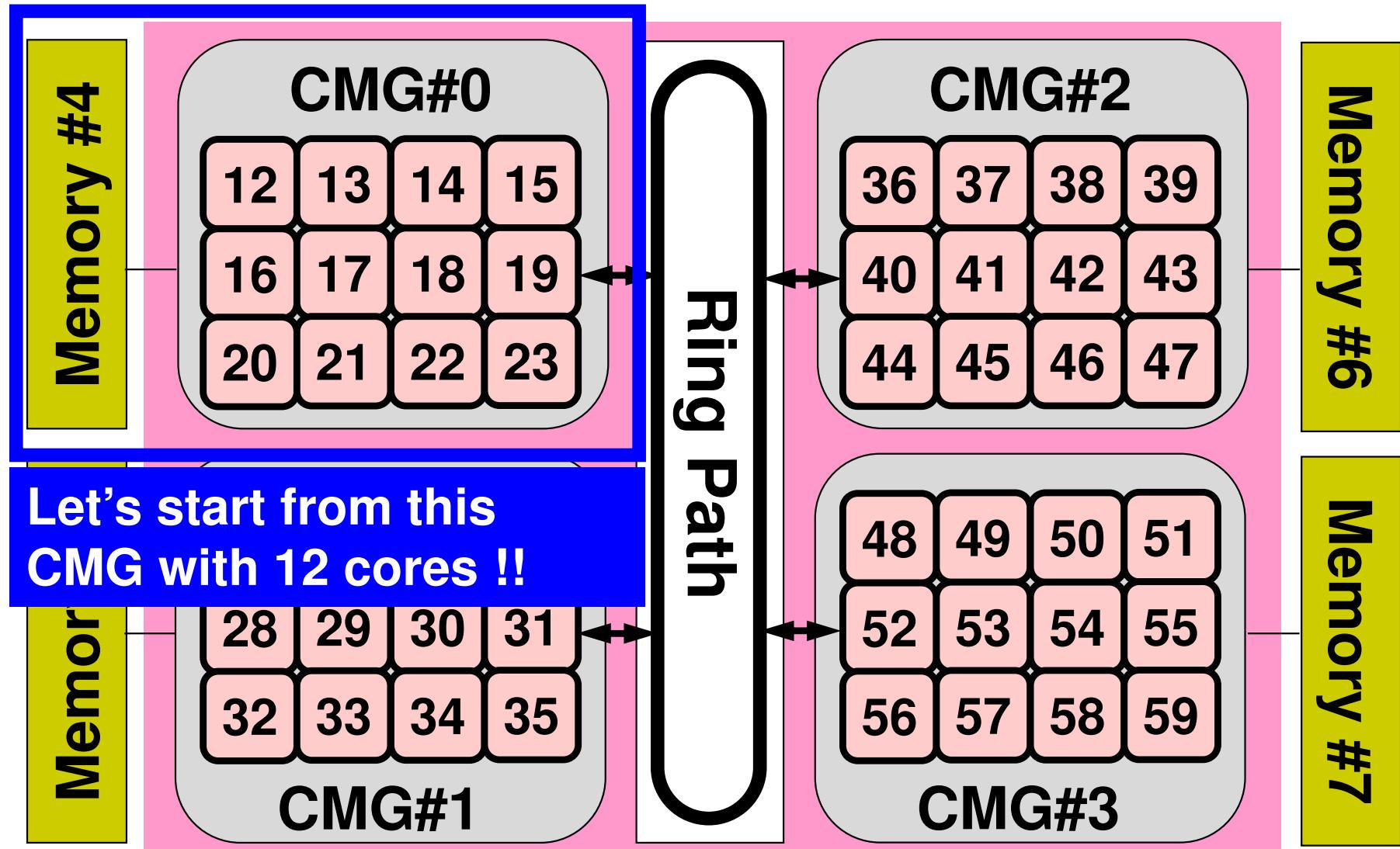
ID of CMGs, Cores, Memory's (1/2)

CMG:#0-#3, Core:#12-59, Memory:#4-#7



ID of CMGs, Cores, Memory's (1/2)

CMG:#0-#3, Core:#12-59, Memory:#4-#7



Files on Wisteria/BDEC-01 (Odyssey)

```
>$ cd /work/gt89/t89XYZ  
>$ cp /work/gt00/z30088/ompw.tar .
```

```
>$ tar xvf ompw.tar
```

```
>$ cd ompw <$O-ompw>
```

Please make sure that following directories are created

```
run src-c0 src-c1 src-c2 src-c3 src-c3b src-f0 src-f1 src-f2 src-f3 src-f4
```

```
>$ module load fj Please type this at every login !
```

```
>$ make -f makef
```

```
>$ ls run/solf0  
      solf0
```

```
>$ cd run
```

```
<modify "INPUT.DAT", "f12.sh">
```

```
>$ pbsub f12.sh
```

<\$O-ompw>/makef

parallel computing by OpenMP

```
default:  
    (cd src-f0 ; make )  
    (cd src-f1 ; make )  
    (cd src-f2 ; make )  
    (cd src-f3 ; make )  
    (cd src-f4 ; make )  
  
clean:  
    (cd src-f0 ; make clean)  
    (cd src-f1 ; make clean)  
    (cd src-f2 ; make clean)  
    (cd src-f3 ; make clean)  
    (cd src-f4 ; make clean)
```

<\$O-ompw>/src-f0/Makefile

parallel computing by OpenMP

```

F90 = frtpx
F90OPTFLAGS= -Kfast,openmp

F90FLAGS = $(F90OPTFLAGS)

.SUFFIXES:
.SUFFIXES: .o .f .f90 .c
#
.f90.o:; $(F90) -c $(F90FLAGS) $(F90OPTFLAG) $<
.f.o:; $(F90) -c $(F90FLAGS) $(F90OPTFLAG) $<
#
OBJS = \
solver_PCG.o struct.o pcg.o \
boundary_cell.o cell_metrics.o \
input.o main.o poi_gen.o pointer_init.o outucl.o

TARGET = ../run/solf0

all: $(TARGET)

$(TARGET): $(OBJS)
    $(F90) $(F90FLAGS) -o $(TARGET) \
    $(OBJS) \
    $(F90FLAGS)

clean:
    rm -f *.o $(TARGET) *.mod *~ PI* *.lst

```

C-Compiler : 2-modes

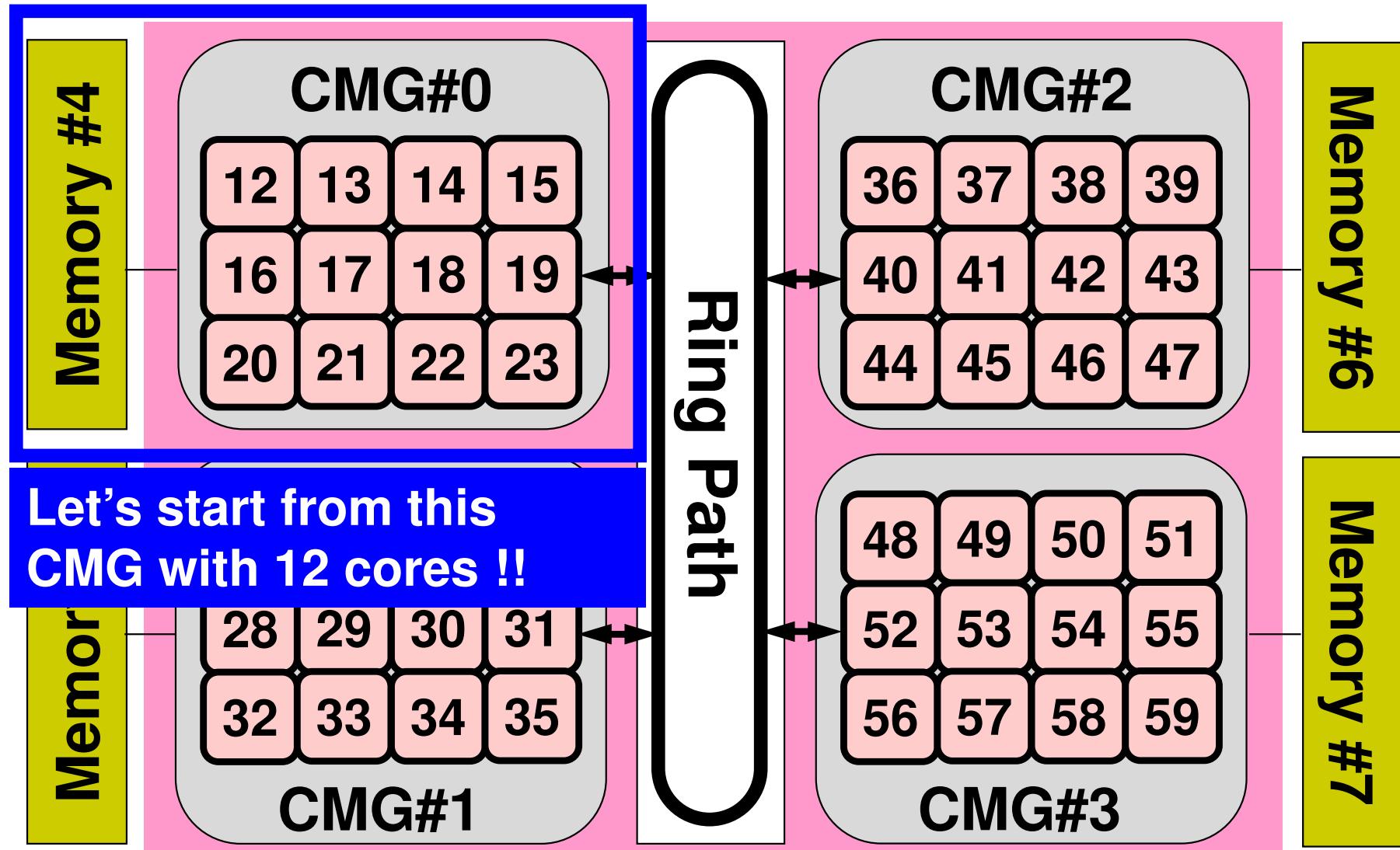
trad (-Nnoclang) (default)	<ul style="list-style-type: none">• Based on Fujitsu's compiler developed for K and PRIMEHPC FX100 or older• Compatible with Fujitsu's Traditional Compilers• C89/C99/C11, OpenMP 3.1/OpenMP 4.5 (partially)• Default (-Nnoclang)• Generally slow for the materials in this class• make -f make-org (make-o)
clang (-Nclang)	<ul style="list-style-type: none">• Based on Clang/LLVM Compilers (Open Source)• Suitable for using Most Updated Capability's, and for using OSS (Open Source Software)• C89/C99/C11, OpenMP 4.5/OpenMP 5.0 (partially)• Generally faster than "trad" modes, difference between "trad" and "clang" is smaller for optimized codes• In this class, default is "clang" mode• make -f makec (Makefile)

Running Job

- Batch Jobs
 - Only batch jobs are allowed.
 - Interactive executions of jobs are not allowed.
- How to run
 - writing job script
 - submitting job
 - checking job status
 - checking results
- Utilization of computational resources
 - 1-node (48 cores) is occupied by each job.
 - Your node is not shared by other jobs.

ID of CMGs, Cores, Memory's (1/2)

CMG:#0-#3, Core:#12-59, Memory:#4-#7



Job Script (1/3): f12.sh

- `/work/gt89/t89xxx/ompw/run/f12.sh`
- Scheduling + Shell Script

```

#!/bin/sh
#PJM -N "f12"
#PJM -L rscgrp=lecture9-o
#PJM -L node=1
#PJM --omp thread=12
#PJM -L elapse=00:15:00
#PJM -g gt89
#PJM -j
#PJM -e err
#PJM -o f12.lst

module load fj
export OMP_NUM_THREADS=12           Thread # (--omp thread=XX)
export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

numactl ./solf0
numactl -C 12-23 -m 4 ./solf0

```

#!/bin/sh
#PJM -N "f12" Job Name (not required)
#PJM -L rscgrp=lecture9-o Name of Queue (Resource Group)
#PJM -L node=1 Node # (=1)
#PJM --omp thread=12 Thread # (1-48, ~12 for a while)
#PJM -L elapse=00:15:00 Elapsed Computation Time
#PJM -g gt89 Group Name (Wallet)
#PJM -j
#PJM -e err Standard Error
#PJM -o f12.lst Standard Output

module load fj
export OMP_NUM_THREADS=12 Thread # (--omp thread=XX)
export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

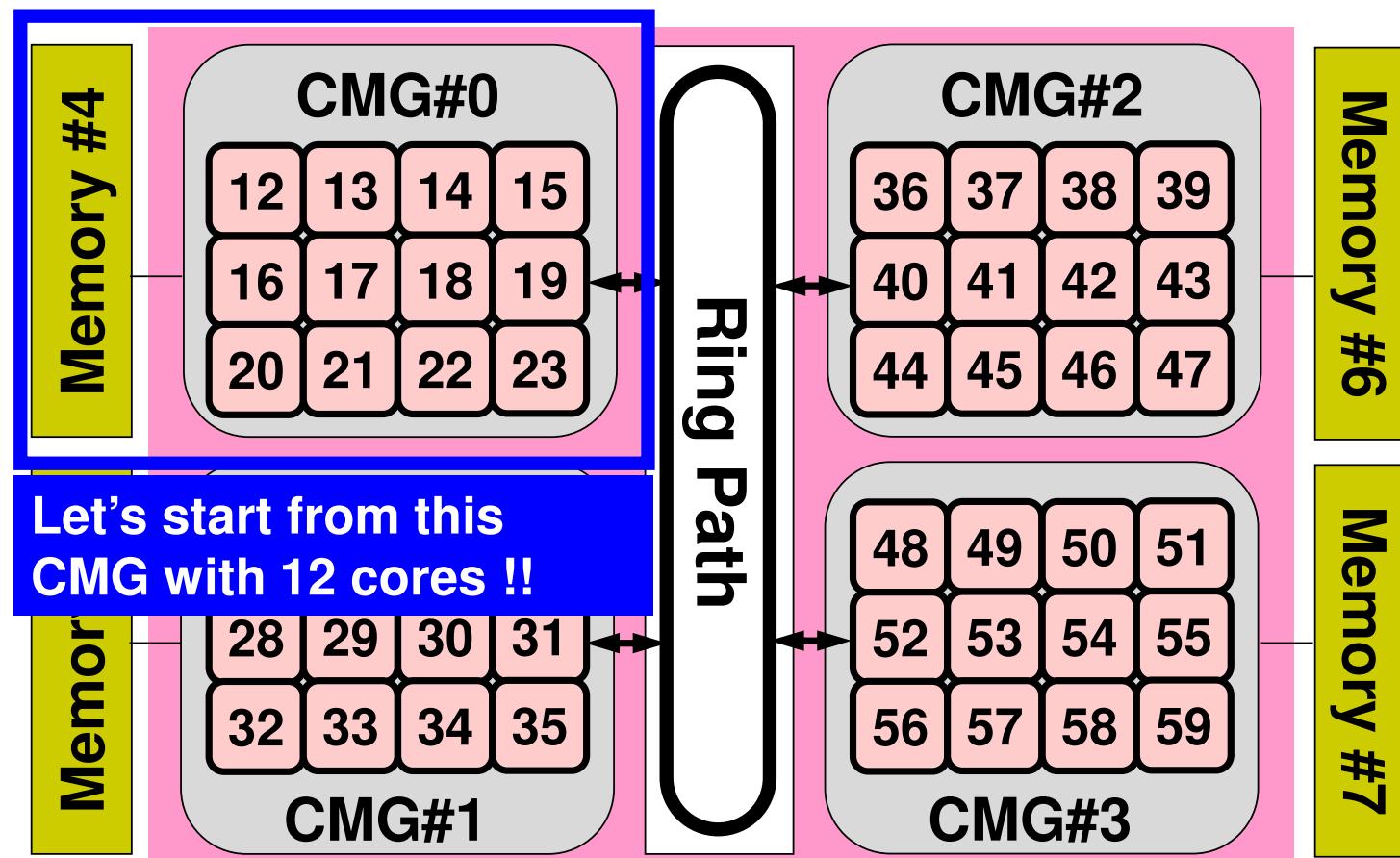
numactl ./solf0
numactl -C 12-23 -m 4 ./solf0

Job Script (2/3): f12.sh

- `/work/gt89/t89xxx/ompw/run/f12.sh`
- `numactl`: utilization of local resources as much as possible

```
numactl ./solf0
```

```
numactl -C 12-23 -m 4 ./solf0
```



Job Script (3/3) : f12.sh

```
export XOS_MMM_L_PAGING_POLICY=
demand:demand:demand
```

Parameters	Values (Underline: Default)	Description
XOS_MMM_L_PAGING _POLICY	[demand <u>prepage</u>] [<u>demand</u> prepage] [demand <u>prepage</u>]	<p>Paging policy (page allocation trigger) of each memory unit</p> <ul style="list-style-type: none"> ✓ demand: Demand Paging Method ✓ prepage: Prepaging Method <p>3 Items are defined</p> <ul style="list-style-type: none"> ✓ 1st Item: .bss area of static data (.data area of static data is always “prepage”) ✓ 2nd Item: Stack Area, Thread Stack Area ✓ 3rd Item: Area for Dynamic Memory Allocation <p>If a value other than the specified value (demand/prepage), the configuration is considered as “prepage:demand:prepage”</p> <p>“demand:demand:demand” is recommended for using multiple CMG’s</p>

Running Jobs

```
>$ cd /work/gt89/t89XYZ  
>$ cd ompw/run  
>$ pbsub f12.sh  
  
>$ cat f12.lst
```

INPUT.DAT

128 128 128	NX NY NZ
1.00e-0 1.00e-00 1.00e-00	DX/DY/DZ
1.0e-08	EPSICCG

Output (1/2)

```
[t00XYZ@wisteria01 run]$ cat f12.1st
```

```
128    128    128
      1    8.958216E+00
     101    8.313496E+00
    201    2.090443E+00
   301    3.811029E-01
  401    3.769653E-02
 501    9.429978E-04
 601    4.940783E-05
 701    1.888611E-06
 801    2.243179E-08
 826    9.818026E-09
      5.275418E+00 sec. (solver)
##ANSWER    2097152    1.459831E+04
```

```
128    128    128
      1    8.958216E+00
     101    8.313496E+00
    201    2.090443E+00
   301    3.811029E-01
  401    3.769653E-02
 501    9.429978E-04
 601    4.940783E-05
 701    1.888611E-06
 801    2.243179E-08
 826    9.818026E-09
      5.270398E+00 sec. (solver)
##ANSWER    2097152    1.459831E+04
```

Output (2/2): 5-times

```
[t00xyz@wisteria01 run]$ grep "(sol" f12.1st
```

```
5.275418E+00 sec. (solver)  
5.270398E+00 sec. (solver)  
5.270445E+00 sec. (solver)  
5.271554E+00 sec. (solver)  
5.270543E+00 sec. (solver)
```

```
numactl ./solf0
```

```
5.272427E+00 sec. (solver)  
5.272081E+00 sec. (solver)  
5.270522E+00 sec. (solver)  
5.271067E+00 sec. (solver)  
5.271237E+00 sec. (solver)
```

```
numactl -C 12-23 -m 4 ./solf0
```

Available Resource Groups (Queue's)

- Following 2 resource groups are available
- Up to 12 nodes are available, while you need a single node in this class
 - **lecture-o**
 - 12 nodes (576 cores), 15 min., valid until the end of Sep. 2022
 - Shared by all “educational” users
 - **lecture9-o**
 - 12 nodes (576 cores), 15 min., active during class time
 - More jobs (compared to **lecture-o**) can be processed up on availability.

Submitting & Checking Jobs

- Submitting Jobs `pjsub SCRIPT NAME`
- Checking status of jobs `pjstat`
- Deleting/aborting `pjdel JOB ID`
- Checking status of queues `pjstat --rsc`
- Detailed info. of queues `pjstat --rsc -x`
- Number of running jobs `pjstat -a`
- History of Submission `pjstat -H`
- Limitation of submission `pjstat --limit`

[t00470@wisteria01 run]\$ pjsub f2_48.sh
[INFO] PJM 0000 pjsub Job 15713 submitted.

[t00470@wisteria01 run]\$ pjsub f3_48.sh
[INFO] PJM 0000 pjsub Job 15714 submitted.

[t00470@wisteria01 run]\$ pjstat
Wisteria/BDEC-01 scheduled stop time: 2021/05/28(Fri) 09:00:00 (Remain: 4days 1:25:56)

JOB_ID	JOB_NAME	STATUS	PROJECT	RSCGROUP	START_DATE	ELAPSE	TOKEN	NODE	GPU
15713	f2_48	RUNNING	gt00	lecture-o	05/24 07:34:03	00:00:02	-	1	-
15714	f3_48	QUEUED	gt00	lecture-o	--/-- --:--:--	00:00:00	-	1	-

[t00470@wisteria01 run]\$ pjstat
Wisteria/BDEC-01 scheduled stop time: 2021/05/28(Fri) 09:00:00 (Remain: 4days 1:25:56)

JOB_ID	JOB_NAME	STATUS	PROJECT	RSCGROUP	START_DATE	ELAPSE	TOKEN	NODE	GPU
15713	f2_48	RUNNING	gt00	lecture-o	05/24 07:34:03	00:00:02	-	1	-
15714	f3_48	RUNNING	gt00	lecture-o	(05/24 07:34)	00:00:00	-	1	-

[t00XYZ@wisteria01 ~]\$ pjdel 15714
[INFO] PJM 0100 pjdel Accepted Job 15714

[t00XYZ@wisteria01 ~]\$ pjstat
Wisteria/BDEC-01 scheduled stop time: 2021/05/28(Fri) 09:00:00 (Remain: 4days 1:25:56)

JOB_ID	JOB_NAME	STATUS	PROJECT	RSCGROUP	START_DATE	ELAPSE	TOKEN	NODE	GPU
15713	f2_48	RUNNING	gt00	lecture-o	05/24 07:34:03	00:00:02	-	1	-

[t00XYZ@wisteria01 ~]\$ pjstat
Wisteria/BDEC-01 scheduled stop time: 2021/05/28(Fri) 09:00:00 (Remain: 4days 1:21:45)

No unfinished job found.

[t00XYZ@wisteria01 ~]\$ pjstat --rsc

SYSTEM: Odyssey

RSCGRP	STATUS	NODE
lecture-o	[ENABLE, START]	96
lecture0-o	[DISABLE, STOP]	2x12x16

[t00XYZ@wisteria01 ~]\$ pjstat --rsc -x

SYSTEM: Odyssey

RSCGRP	STATUS	MIN_NODE	MAX_NODE	MAX_ELAPSE	REMAIN_ELAPSE	MEM(GiB)	PROJECT
lecture-o	[ENABLE, START]	1	12	00:15:00	00:15:00	28	gt00
lecture0-o	[DISABLE, STOP]	1	12	00:15:00	--:--:--	28	gt00

[t00XYZ@wisteria01 ~]\$ pjstat --limit

SYSTEM: Odyssey

PROJECT	ACCEPT	RUN	BULK_ACCEPT	BULK_RUN	NODE
gt80	0/ 128	0/ 16	0/ 8	0/ 16	0/ 2304

SYSTEM: Aquarius

PROJECT	ACCEPT	RUN	BULK_ACCEPT	BULK_RUN	GPU
gt80	0/ 4	0/ 2	0/ 0	0/ 0	0/ 64

poi_gen (1/2): Main Part

Be carefule with “private” !!

```

!$omp parallel do private (icel, icN1, icN2, icN3, icN4, icN5, icN6, VOL0)      &
!$omp&          private (icou, k, coef, ii, jj, kk)
do icel= 1, ICELTOT
    icN1= NEIBcell(icel, 1)
    icN2= NEIBcell(icel, 2)
    icN3= NEIBcell(icel, 3)
    icN4= NEIBcell(icel, 4)
    icN5= NEIBcell(icel, 5)
    icN6= NEIBcell(icel, 6)

    VOL0= VOLCEL(icel)

    icou= 0
    if (icN5. ne. 0) then
        coef = RDZ * ZAREA
        D(icel)= D(icel) - coef
        icou= icou + 1
        k = icou + indexLU(icel-1)
        itemLU(k)= icN5
        AMAT(k)= coef
    endif
    (...)

    ii= XYZ(icel, 1)
    jj= XYZ(icel, 2)
    kk= XYZ(icel, 3)

    BFORCE(icel)= -dfloat(ii+jj+kk) * VOL0

enddo

```

poi_gen (2/2): Boundary Conditions

Be carefule with “private” !!

```
!$omp parallel do private (ib, icel, coef)
    do ib= 1, ZmaxCELtot
        icel= ZmaxCEL(ib)
        coef= 2. d0 * RDZ * ZAREA
        D(icel)= D(icel) - coef
    enddo
```

solve_PCG (1/5)

parallel computing by OpenMP

```
module solver_PCG
    contains
!C
!C*** solve_PCG
!C
    subroutine solve_PCG
&           ( N, NPLU, indexLU, itemLU, D, B, X, AMAT, EPS, ITR, IER)          &
use omp_lib
implicit REAL*8 (A-H, O-Z)

real(kind=8), dimension(N)      :: D
real(kind=8), dimension(N)      :: B
real(kind=8), dimension(N)      :: X
real(kind=8), dimension(NPLU)   :: AMAT

integer, dimension(0:N)        :: indexLU
integer, dimension(NPLU)       :: itemLU

real(kind=8), dimension(:, :), allocatable :: W

integer, parameter :: R= 1
integer, parameter :: Z= 2
integer, parameter :: Q= 2
integer, parameter :: P= 3
integer, parameter :: DD= 4
```

solve_PCG (2/5)

```

!$omp parallel do private(i)
do i= 1, N
    X(i) = 0. d0
    W(i, 2)= 0. ODO
    W(i, 3)= 0. ODO
    W(i, DD)= 1. d0/D(i)
enddo

!$omp parallel do private(i, VAL, k)
do i= 1, N
    VAL= D(i)*X(i)
    do k= indexLU(i-1)+1, indexLU(i)
        VAL= VAL + AMAT(k)*X(itemLU(k))
    enddo
    W(i, R)= B(i) - VAL
enddo

BNRM2= 0. ODO
!$omp parallel do private(i) reduction(+:BNRM2)
do i= 1, N
    BNRM2 = BNRM2 + B(i) **2
enddo

```

Compute $r^{(0)} = b - [A]x^{(0)}$

for $i = 1, 2, \dots$
 solve $[M] z^{(i-1)} = r^{(i-1)}$
 $\rho_{i-1} = r^{(i-1)} \cdot z^{(i-1)}$
 if $i = 1$
 $p^{(1)} = z^{(0)}$
 else
 $\beta_{i-1} = \rho_{i-1} / \rho_{i-2}$
 $p^{(i)} = z^{(i-1)} + \beta_{i-1} p^{(i-1)}$
 endif
 $q^{(i)} = [A]p^{(i)}$
 $\alpha_i = \rho_{i-1} / p^{(i)} q^{(i)}$
 $x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$
 $r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$
 check convergence $|r|$
end

solve_PCG (3/5)

```

ITR= N
Stime= omp_get_wtime()

do L= 1, ITR

!$omp parallel do private(i)
do i= 1, N
    W(i, Z)= W(i, R)*W(i, DD)
enddo

RHO= 0. d0
!$omp parallel do private(i) reduction(+:RHO)
do i= 1, N
    RHO= RHO + W(i, R)*W(i, Z)
enddo

if ( L.eq. 1 ) then
!$omp parallel do private(i)
    do i= 1, N
        W(i, P)= W(i, Z)
    enddo
else
    BETA= RHO / RH01
!$omp parallel do private(i)
    do i= 1, N
        W(i, P)= W(i, Z) + BETA*W(i, P)
    enddo
endif

```

Compute $r^{(0)} = b - [A]x^{(0)}$

for $i = 1, 2, \dots$

solve $[M]z^{(i-1)} = r^{(i-1)}$

$\rho_{i-1} = r^{(i-1)} \cdot z^{(i-1)}$

if $i=1$

$p^{(1)} = z^{(0)}$

else

$\beta_{i-1} = \rho_{i-1}/\rho_{i-2}$

$p^{(i)} = z^{(i-1)} + \beta_{i-1} p^{(i-1)}$

endif

$q^{(i)} = [A]p^{(i)}$

$\alpha_i = \rho_{i-1}/p^{(i)} q^{(i)}$

$x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$

$r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$

check convergence $|r|$

end

solve_PCG (4/5)

```

!$omp parallel do private(i, VAL, k)
do i= 1, N
    VAL= D(i)*W(i, P)
    do k= indexLU(i-1)+1, indexLU(i)
        VAL= VAL + AMAT(k)*W(itemLU(k), P)
    enddo
    W(i, Q)= VAL
enddo
C1= 0. d0
!$omp parallel do private(i) reduction(+:C1)
do i= 1, N
    C1= C1 + W(i, P)*W(i, Q)
enddo
ALPHA= RHO / C1

!$omp parallel do private(i)
do i= 1, N
    X(i) = X(i) + ALPHA * W(i, P)
    W(i, R)= W(i, R) - ALPHA * W(i, Q)
enddo

DNRM2= 0. d0
!$omp parallel do private(i) reduction(+:DNRM2)
do i= 1, N
    DNRM2= DNRM2 + W(i, R)**2
enddo

ERR = dsqrt(DNRM2/BNRM2)...

```

Compute $r^{(0)} = b - [A]x^{(0)}$

for $i = 1, 2, \dots$

solve $[M]z^{(i-1)} = r^{(i-1)}$

$\rho_{i-1} = r^{(i-1)} \cdot z^{(i-1)}$

if $i = 1$

$p^{(1)} = z^{(0)}$

else

$\beta_{i-1} = \rho_{i-1} / \rho_{i-2}$

$p^{(i)} = z^{(i-1)} + \beta_{i-1} p^{(i-1)}$

endif

$q^{(i)} = [A]p^{(i)}$

$\alpha_i = \rho_{i-1} / p^{(i)} q^{(i)}$

$x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$

$r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$

check convergence $|r|$

end

solve_PCG (5/5)

```

Stime = omp_get_wtime() ←
do L= 1, ITR
...
  if (ERR .lt. EPS) then
    IER = 0
    goto 900
  else
    RH01 = RHO
  endif

enddo
IER = 1
900 continue
Etime= omp_get_wtime() ←

write (*, '(i5,2(1pe16.6))') L, ERR
write (*, '(1pe16.6, a)') Etime-Stime, ' sec. (solver)'

ITR= L
deallocate (W)

return
end

```

Elapsed Time= Etime - Stime

f01,f02,f04,f06,f08,f12.sh

- /work/gt89/t89xxx/ompw/run/fYZ.sh
- Scheduling + Shell Script

```

#!/bin/sh
#PJM -N "fYZ"
#PJM -L rscgrp=lecture0-o
#PJM -L node=1
#PJM --omp thread=YZ
#PJM -L elapse=00:15:00
#PJM -g gt89
#PJM -j
#PJM -e err
#PJM -o fYZ.lst

module load fj
export OMP_NUM_THREADS=YZ           Thread # (--omp thread=YZ)
export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

numactl ./solf0
numactl -C 12-23 -m 4 ./solf0

```

#!/bin/sh
#PJM -N "fYZ" Job Name (not required)
#PJM -L rscgrp=lecture0-o Name of Queue (Resource Group)
#PJM -L node=1 Node # (=1)
#PJM --omp thread=YZ Thread # (1-48, ~12 for a while)
#PJM -L elapse=00:15:00 Elapsed Computation Time
#PJM -g gt89 Group Name (Wallet)
#PJM -j
#PJM -e err Standard Error
#PJM -o fYZ.lst Standard Output

module load fj
export OMP_NUM_THREADS=YZ Thread # (--omp thread=YZ)
export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

numactl ./solf0
numactl -C 12-23 -m 4 ./solf0

Time for PCG: Etime-Stime : Fortran

NX=NY=NZ=128

5 measurements, best result

Efficiency decreases as core# increases: decreasing of memory performance



Thread #	sec	Speed-up	Parallel Efficiency (%)
1	50.27	1.00	100.00
2	25.24	1.99	99.60
4	12.98	3.87	96.86
6	9.24	5.44	90.73
8	7.27	6.92	86.50
12	5.09	9.88	82.30

Parallel Efficiency(%) = 100 * (Speed-Up) / Thread#

f04.sh

```

#!/bin/sh
#PJM -N "f04"
#PJM -L rscgrp=lecture9-o
#PJM -L node=1
#PJM --omp thread=4
#PJM -L elapse=00:15:00
#PJM -g gt89
#PJM -j
#PJM -e err
#PJM -o f04.1st

module load fj
export OMP_NUM_THREADS=4
export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

numactl      ./solf0
numactl      ./solf0
numactl      ./solf0
numactl      ./solf0
numactl      ./solf0
numactl      -C    12-23  -m      4      ./solf0

```

f08.sh

```

#!/bin/sh
#PJM -N "f08"
#PJM -L rscgrp=lecture9-o
#PJM -L node=1
#PJM --omp thread=8
#PJM -L elapse=00:15:00
#PJM -g gt89
#PJM -j
#PJM -e err
#PJM -o f08.lst

module load fj
export OMP_NUM_THREADS=8
export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

numactl      ./solf0
numactl      ./solf0
numactl      ./solf0
numactl      ./solf0
numactl      ./solf0
numactl      -C    12-23  -m      4      ./solf0

```

Multiple CMG's

Time for PCG: Etime-Stime : Fortran

NX=NY=NZ=128

5 measurements, best result



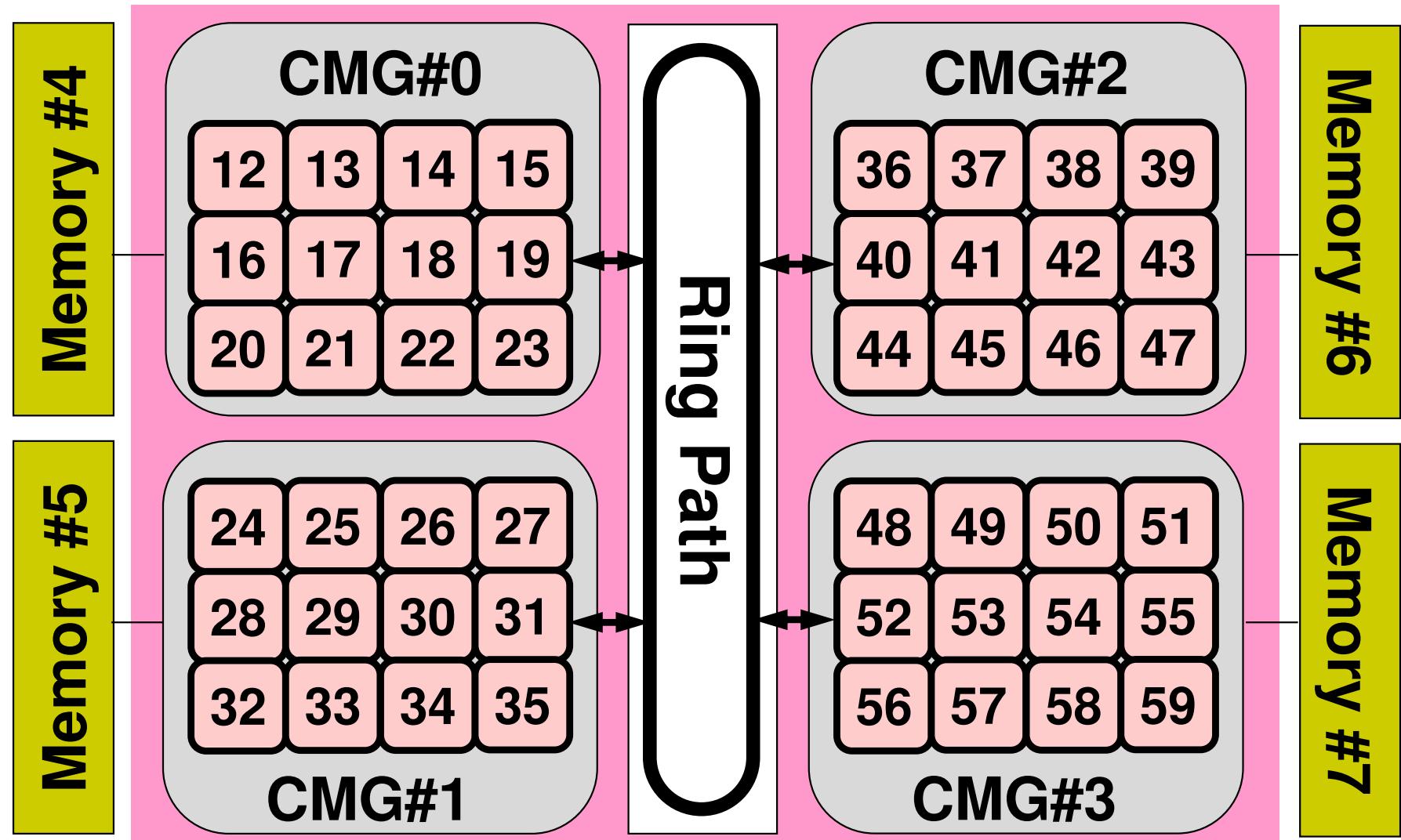
Granularity : 粒度
Problem Size/Thread

Thread #	sec	Speed-up	Parallel Efficiency (%)
12	5.09	12.00	100.00
24	2.79	21.88	91.18
36	1.99	30.75	85.41
48	1.70	35.97	74.95

$$\text{Parallel Efficiency}(\%) = 100 * (\text{Speed-Up}) / \text{Thread\#}$$

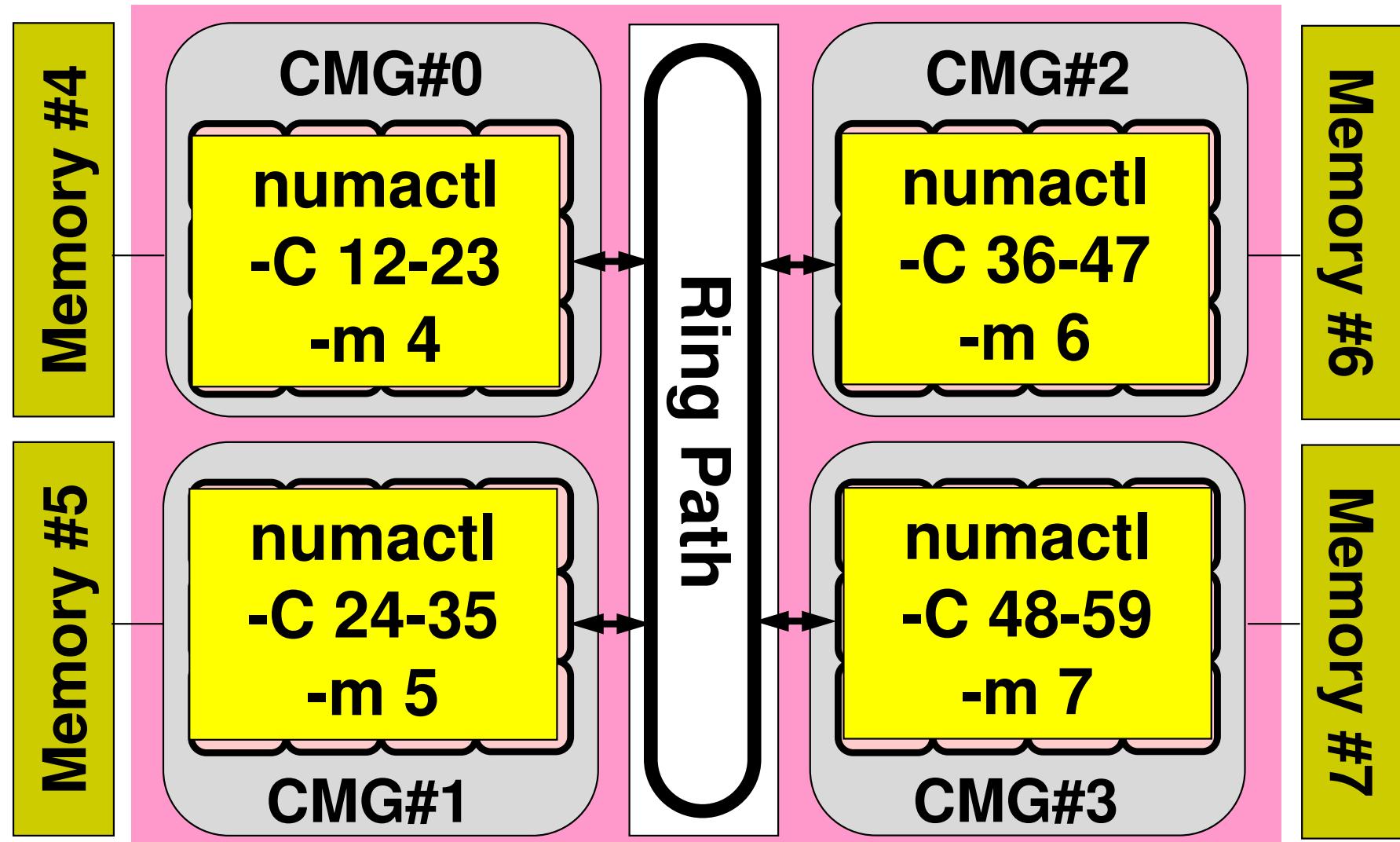
ID of CMGs, Cores, Memory's (1/2)

CMG:#0-#3, Core:#12-59, Memory:#4-#7



ID of CMGs, Cores, Memory's (2/2)

CMG:#0-#3, Core:#12-59, Memory:#4-#7



f0_12.sh

```

#!/bin/sh
#PJM -N "f0_12"
#PJM -L rscgrp=lecture9-o
#PJM -L node=1
#PJM --omp thread=12
#PJM -L elapse=00:15:00
#PJM -g gt89
#PJM -j
#PJM -e err
#PJM -o f0_12.lst

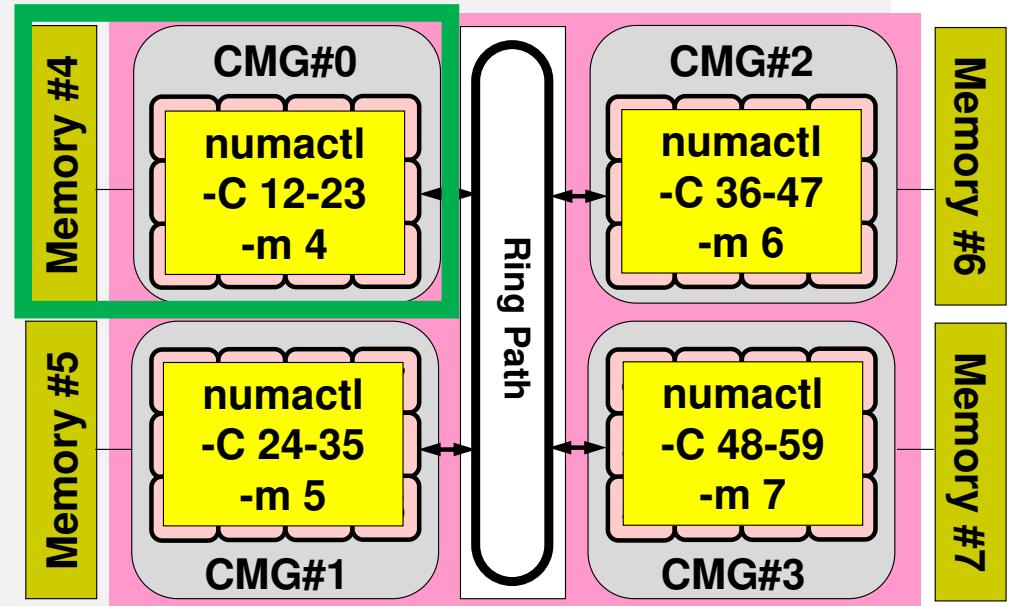
```

module load fj

export OMP_NUM_THREADS=12

export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

numactl	./solf0				
numactl	./solf0				
numactl	./solf0				
numactl	./solf0				
numactl	./solf0				
numactl	-C 12-23 -m 4	./solf0			
numactl	-C 12-23 -m 4	./solf0			
numactl	-C 12-23 -m 4	./solf0			
numactl	-C 12-23 -m 4	./solf0			
numactl	-C 12-23 -m 4	./solf0			



f0_24.sh

```
#!/bin/sh
#PJM -N "f0_24"
#PJM -L rscgrp=lectur9-o
#PJM -L node=1
#PJM --omp thread=24
#PJM -L elapse=00:15:00
#PJM -g gt89
#PJM -j
#PJM -e err
#PJM -o f0_24.lst
```

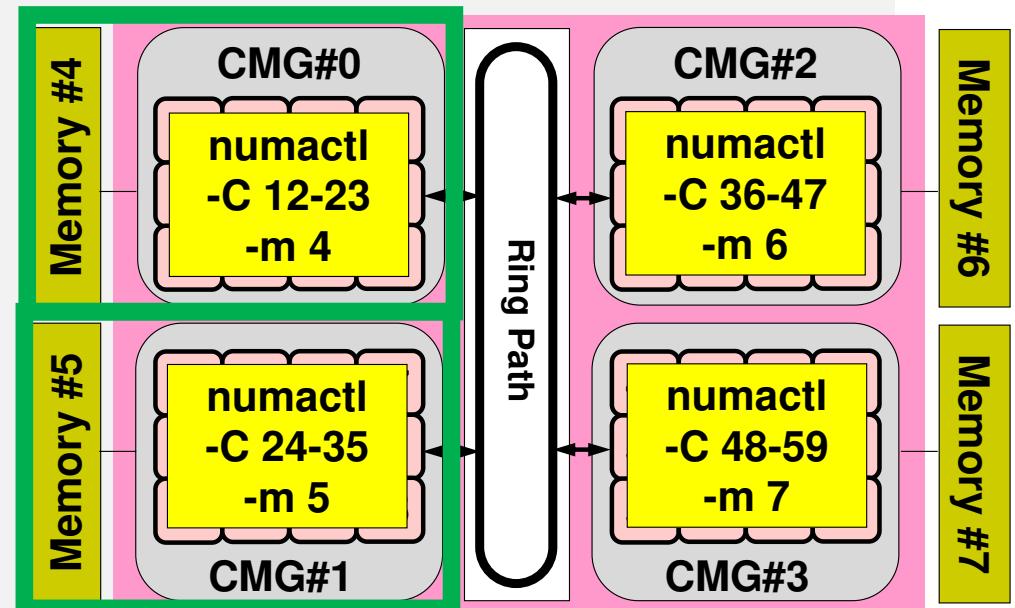
module load fj

export OMP_NUM_THREADS=24

export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

```
numactl ./solf0
numactl ./solf0
numactl ./solf0
numactl ./solf0
numactl ./solf0
```

numactl	-C	12-35	-m	4-5	./solf0
numactl	-C	12-35	-m	4-5	./solf0
numactl	-C	12-35	-m	4-5	./solf0
numactl	-C	12-35	-m	4-5	./solf0
numactl	-C	12-35	-m	4-5	./solf0



f0_36.sh

```
#!/bin/sh
#PJM -N "f0_36"
#PJM -L rscgrp=lecture9-o
#PJM -L node=1
#PJM --omp thread=36
#PJM -L elapse=00:15:00
#PJM -g gt89
#PJM -j
#PJM -e err
#PJM -o f0_36.lst
```

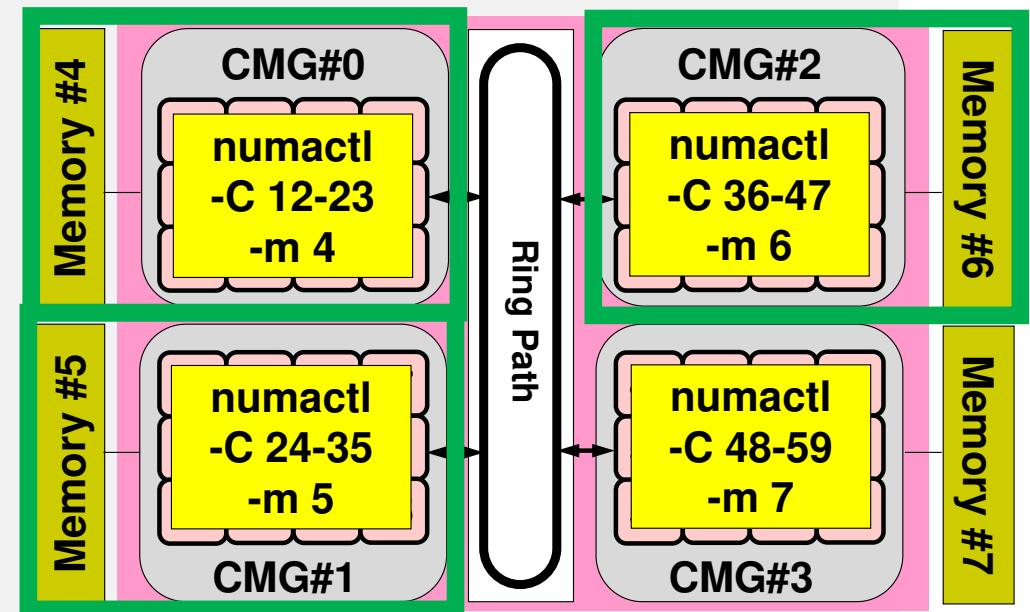
module load fj

export OMP_NUM_THREADS=36

export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

```
numactl      ./solf0
numactl      ./solf0
numactl      ./solf0
numactl      ./solf0
numactl      ./solf0
```

numactl	-C 12-47 -m 4-6	./solf0
numactl	-C 12-47 -m 4-6	./solf0
numactl	-C 12-47 -m 4-6	./solf0
numactl	-C 12-47 -m 4-6	./solf0
numactl	-C 12-47 -m 4-6	./solf0



f0_48.sh

```
#!/bin/sh
#PJM -N "f0_48"
#PJM -L rscgrp=lecture9-o
#PJM -L node=1
#PJM --omp thread=48
#PJM -L elapse=00:15:00
#PJM -g gt89
#PJM -j
#PJM -e err
#PJM -o f0_48.lst
```

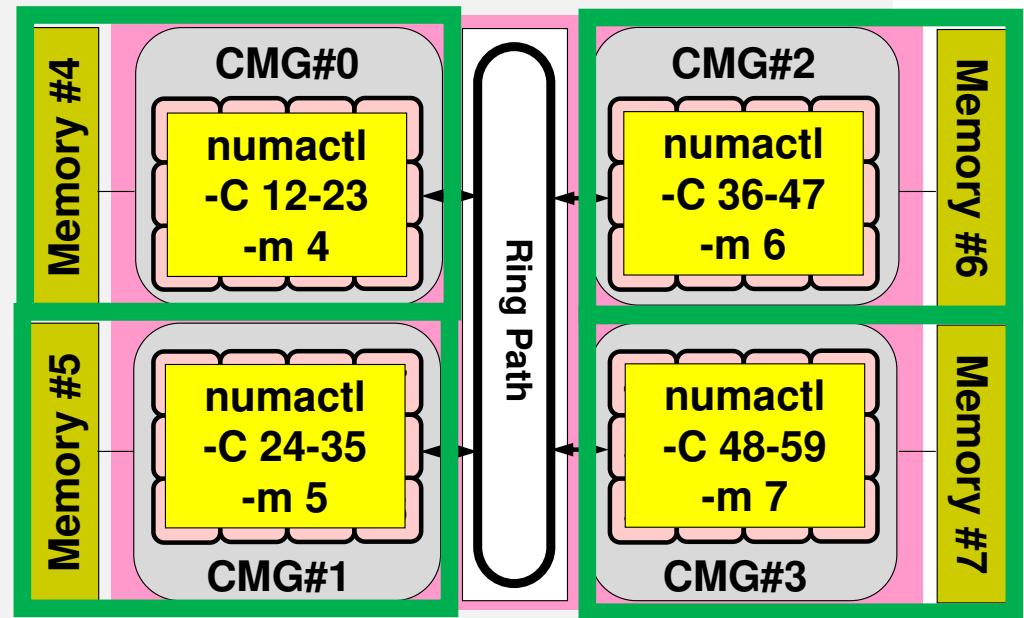
module load fj

export OMP_NUM_THREADS=48

export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

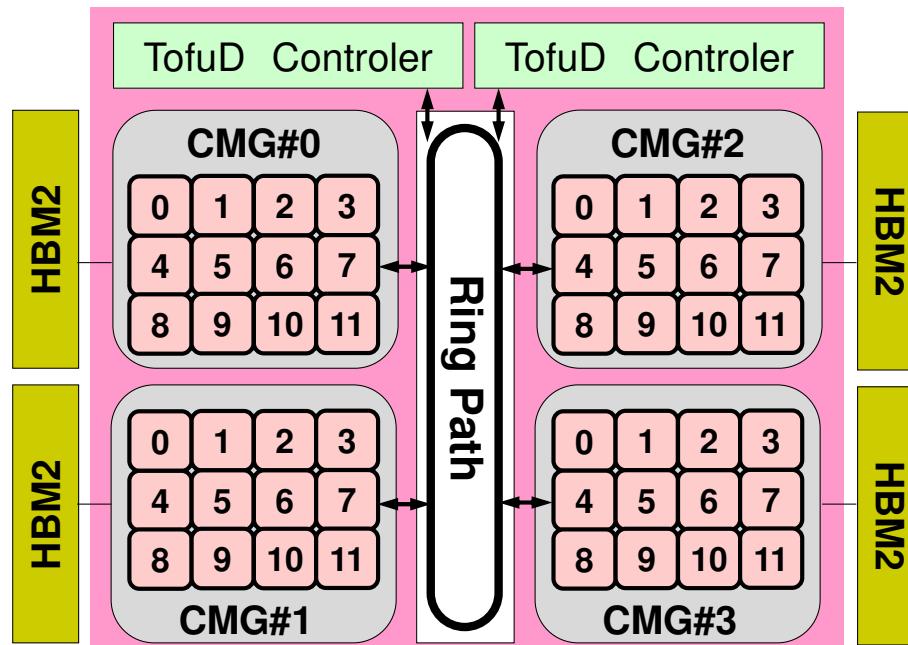
numactl	./solf0

numactl	-C 12-59 -m 4-7 ./solf0
numactl	-C 12-59 -m 4-7 ./solf0
numactl	-C 12-59 -m 4-7 ./solf0
numactl	-C 12-59 -m 4-7 ./solf0
numactl	-C 12-59 -m 4-7 ./solf0



- OpenMP
- Login to Wisteria/BDEC-01
- Parallel Code by OpenMP (0): up to 12 cores
- **Parallel Code by OpenMP (1): First Touch**
- Parallel Code by OpenMP (2): +ELL
- Parallel Code by OpenMP (3): reduced omp-parallel
- Parallel Code by OpenMP (4): Further Optimization (Fortran only)

A64FX Processor on Odyssey



Name	A64FX
Processor # (Core #)	1 (48+ 2or4 Assistant Cores)
Frequency	2.2 GHz
Peak Performance	3.3792 TFLOPS
Memory Size	32 GiB
Memory Bandwidth	1,024 GB/s
L1 Cache	64 KiB/core (Inst/Data)
L2 Cache	8 MiB/CMG

- 4 CMG's (Core Memory Group), 12 cores/CMG
 - 48 Cores/Node (Processor)
 - $2.2\text{GHz} \times 32\text{DP} \times 48 = 3379.2 \text{ GFLOPS} = 3.3792 \text{ TFLOPS}$
- **NUMA Architecture (Non-Uniform Memory Access)**
 - Each core of a CMG can access to the memory on other CMG's
 - Utilization of the local memory is more efficient

First Touch Data Placement

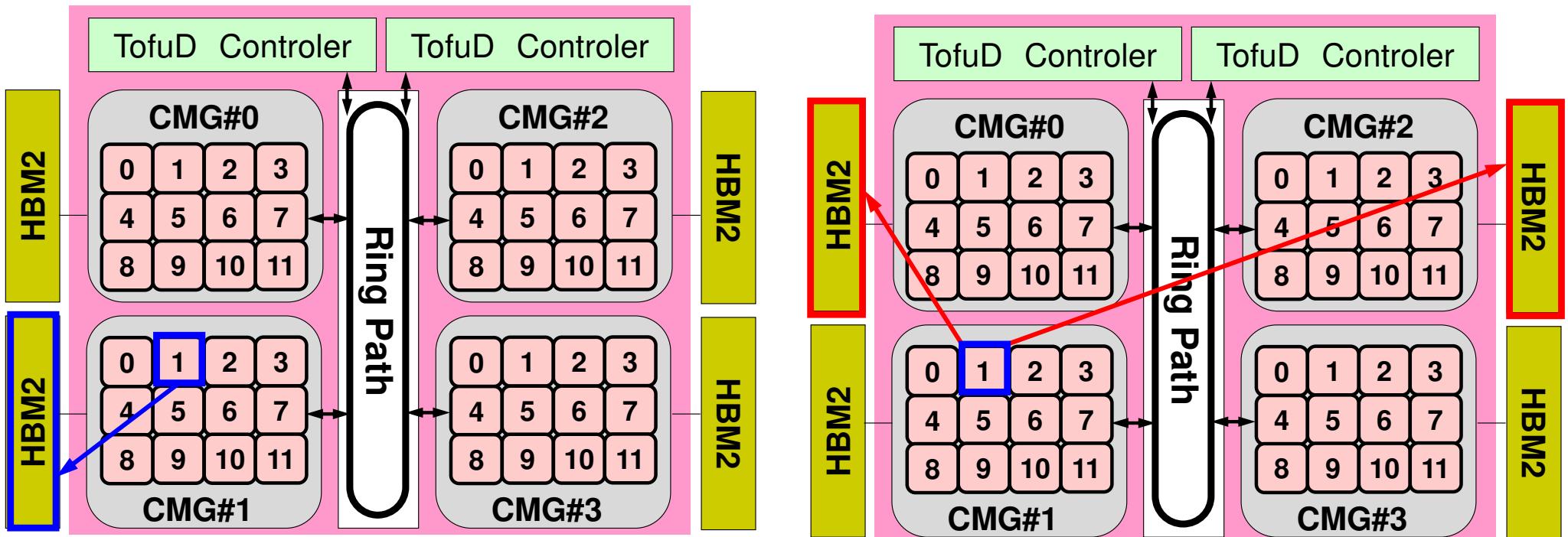
“Patterns for Parallel Programming” Mattson, T.G. et al.

- To reduce memory traffic in the system, it is important to keep the data close to the PEs that will work with the data (e.g. NUMA control).
- On NUMA computers, this corresponds to making sure the pages of memory are allocated and “owned” by the PEs that will be working with the data contained in the page.
 - ✓ Page/Memory Page/Virtual Page: A fixed-length continuous block of virtual memory, smallest unit of data for memory management in a virtual memory OS
- The most common NUMA page-placement algorithm is the “first touch” algorithm, in which the PE first referencing a region of memory will have the page holding that memory assigned to it.
- A very common technique in OpenMP program for optimization is to initialize data in parallel using the same loop schedule as will be used later in the computations.

Summary: First Touch Data Placement

- On NUMA architecture (Non-Uniform Memory Access), “pages of memory” are not allocated when variables and arrays are declared/allocated in the program.
- “Pages” are allocated at the local memory of the “socket” for the “core/thread” that first touches the variables and/or arrays.
- If the pages are not on the local memory of the socket for each thread, performance of the program is very bad.
- A very common technique in OpenMP program for optimization is to initialize data in parallel using the same loop schedule as will be used later in the computations.
- You have to consider this if you use multiple CMG's of the Odyssey system for a single OpenMP program
 - If you don't care, all pages are created at the local memory of CMG#0
 - Not needed for a single CMG case

Local/Remote Memory



Local Memory

Remote Memory

Program by “First-Touch”: src-f1 Original Program in src-f0

```
>$ cd /work/gt89/t89XYZ/ompw
>$ cd run

<modify "INPUT.DAT", "f1_XY.sh"> (XY:12,24,36,48)

>$ pbsub f1_XY.sh
```

```
[XYZ@wisteria01 run]$ cd ../src-f0
[XYZ@wisteria01 src-f0]$ diff poi_gen.f ../src-f1/poi_gen.f
25,29c25,31
```

```
<      PHI    = 0.d0
<      BFORCE= 0.d0
<      D     = 0.d0
<
<      INLU= 0
```

src-f0

```
> !$omp parallel do private (icel)
>       do icel= 1, ICELTOT
>           PHI  (icel)= 0.d0
>           BFORCE(icel)= 0.d0
>           D    (icel)= 0.d0
>           INLU (icel)= 0
>       enddo
```

src-f1

71, 72c73, 75

```
<     indexLU= 0
<
<     do icel= 1, ICELTOT
<         indexLU(icel)= INLU(icel)
<     enddo
```

src-f0

```
>     indexLU(0)= 0
>
> !$omp parallel do private (icel)
>     do icel= 1, ICELTOT
>         indexLU(icel)= INLU(icel)
>     enddo
```

src-f1

85, 86c88, 94

```
<     itemLU= 0
<     AMAT= 0.d0
```

src-f0

```
> !$omp parallel do private (icel,k)
>     do icel= 1, ICELTOT
>         do k= indexLU(icel-1)+1, indexLU(icel)
>             itemLU(k)= 0
>             AMAT(k)= 0.d0
>         enddo
>     enddo
```

src-f1

f1_48.sh

```
#!/bin/sh
#PJM -N "f1_48"
#PJM -L rscgrp=lecture9-o
#PJM -L node=1
#PJM --omp thread=48
#PJM -L elapse=00:15:00
#PJM -g
#PJM -j
#PJM -e
#PJM -o f1_48.lst
```

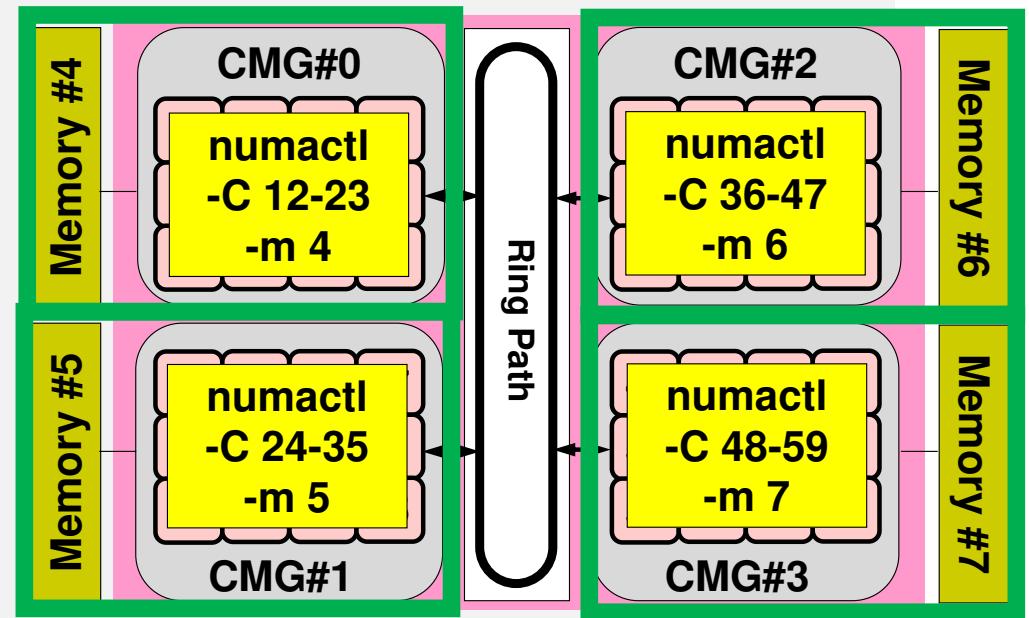
module load fj

export OMP_NUM_THREADS=48

export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

numactl	./solf1

numactl	-C 12-59 -m 4-7 ./solf1
numactl	-C 12-59 -m 4-7 ./solf1
numactl	-C 12-59 -m 4-7 ./solf1
numactl	-C 12-59 -m 4-7 ./solf1
numactl	-C 12-59 -m 4-7 ./solf1



Multiple CMG's

Time for PCG: Etime-Stime : Fortran

$\text{NX}=\text{NY}=\text{NZ}=128$, Best Case for 5 Measurements
based on src-f0 with 12-threads

	Thread #	sec	Speed-up	Parallel Efficiency (%)
src-f0	12	5.09	12.00	100.00
	24	2.79	21.88	91.18
	36	1.99	30.75	85.41
	48	1.70	35.97	74.95

	12	5.27	-	-
src-f1	24	2.70	22.61	94.20
	36	1.87	32.65	90.69
	48	1.52	40.09	83.51

Multiple CMG's

Time for PCG: Etime-Stime : C(clang)

NX=NY=NZ=128, Best Case for 5 Measurements
based on src-c0 with 12-threads

	Thread #	sec	Speed-up	Parallel Efficiency (%)
src-c0	12	5.03	12.00	100.00
	24	2.69	22.42	93.41
	36	1.91	31.57	87.69
	48	1.57	38.51	80.23

	12	5.22	-	-
	24	2.70	22.39	93.29
	36	1.85	32.66	90.71
	48	1.17	51.70	107.7

Cache is well-utilized,
because the problem size is small

Multiple CMG's

Time for PCG: Etime-Stime

$NX=NY=NZ=128$, Best Case for 5 Measurements

	Language	12	24	36	48
Original	Fortran	5.09	2.79	1.99	1.70
	C (clang)	5.03	2.69	1.91	1.57
	C (trad)	7.75	4.19	2.90	2.36

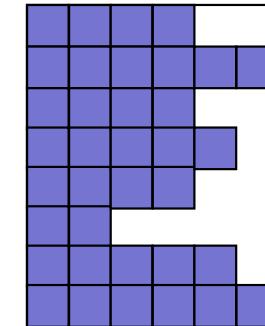
First-Touch	Fortran	5.27	2.70	1.87	1.52
	C (clang)	5.22	2.70	1.85	1.17
	C (trad)	7.85	4.05	2.79	1.72 (112.6 %)

- OpenMP
- Login to Wisteria/BDEC-01
- Parallel Code by OpenMP (0): up to 12 cores
- Parallel Code by OpenMP (1): First Touch
- **Parallel Code by OpenMP (2): +ELL**
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Storage of Sparse Matrices

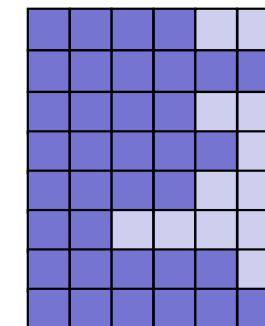
CRS (Compressed Row Storage)

```
do i= 1, N
  W(i, Q) = D(i)*W(i, P)
  do k= indexLU(i-1)+1, indexLU(i)
    W(i, Q) = W(i, Q) + AMAT(k)*W(itemLU(k), P)
  enddo
enddo
```



ELL (ELLPACK/ITPACK)

```
do i= 1, N
  W(i, Q) = D(i)*W(i, P)
  do j= 1, 6
    W(i, Q) = W(i, Q) + AMAT(j, i)*W(itemLU(j, i), P)
  enddo
enddo
```

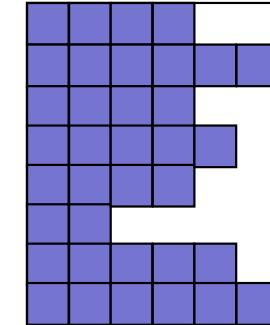


- CRS: Compressed Row Storage
 - Only non-zero off-diag's: saving memory, low performance
- ELL : ELLPACK/ITPACK
 - Fixed # of non-zero off-diag's, 0 padding needed
 - More expensive, better memory performance due to prefetch

Storage of Sparse Matrices

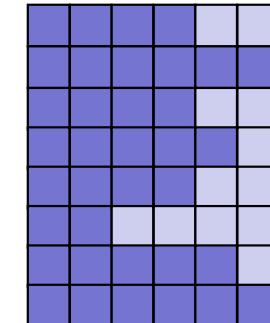
CRS (Compressed Row Storage)

```
do i= 1, N
  W(i, Q) = D(i)*W(i, P)
  do k= indexLU(i-1)+1, indexLU(i)
    W(i, Q) = W(i, Q) + AMAT(k)*W(itemLU(k), P)
  enddo
enddo
```



ELL (ELLPACK/ITPACK)

```
do i= 1, N
  W(i, Q) = D(i)*W(i, P)
  do j= 1, 6
    W(i, Q) = W(i, Q) + AMAT(j, i)*W(itemLU(j, i), P)
  enddo
enddo
```



ELL (ELLPACK/ITPACK) This is a little bit slow

```
do i= 1, N
  W(i, Q) = D(i)*W(i, P)
  do j= 1, 6
    W(i, Q) = W(i, Q) + AMAT(6*(i-1)+k)*W(itemLU(6*(i-1)+j), P)
  enddo
enddo
```

poi_gen (1/2): Private

```

!C
!C-- INTERIOR & NEUMANN BOUNDARY CELLS

!$omp parallel do private (icel, icN1, icN2, icN3, icN4, icN5, icN6, VOL0)      &
!$omp&                      private (icou, k, coef, i i, jj, kk)

do icel= 1, ICELTOT
  icN1= NEIBcell(icel, 1)
  icN2= NEIBcell(icel, 2)
  icN3= NEIBcell(icel, 3)
  icN4= NEIBcell(icel, 4)
  icN5= NEIBcell(icel, 5)
  icN6= NEIBcell(icel, 6)

  VOL0= VOLCEL(icel)

  icou= 0
  if (icN5.ne.0) then
    coef= RDZ * ZAREA
    D(icel)= D(icel) - coef
    icou      = icou + 1
    itemLU (icou, icel)= icN5
    AMAT(icou, icel)= coef
  endif

  if (icN3.ne.0) then
    coef= RDY * YAREA
    D(icel)= D(icel) - coef
    icou      = icou + 1
    itemLU (icou, icel)= icN3
    AMAT(icou, icel)= coef
  endif
  (...)

enddo

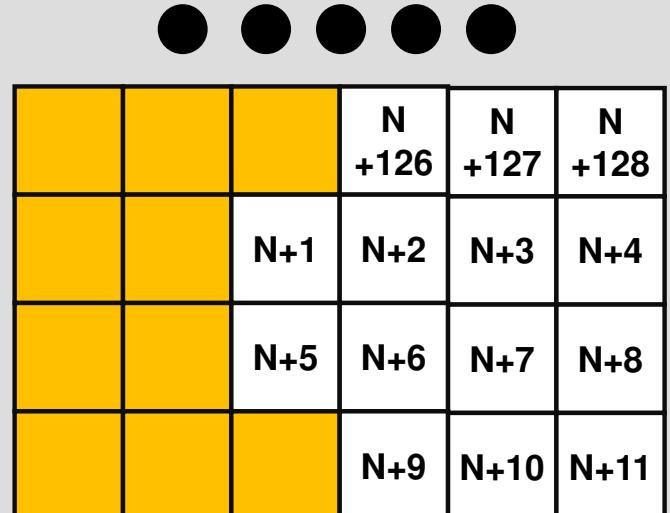
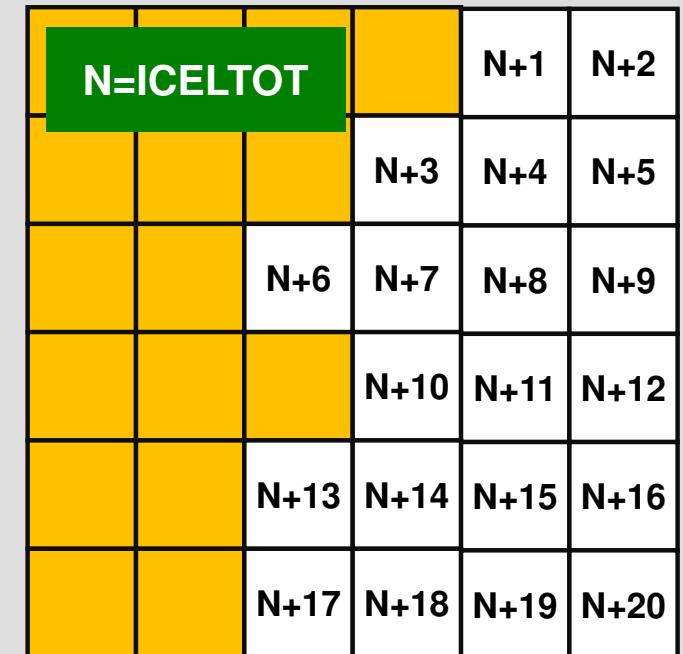
```

poi_gen (2/2): Padding

N2= 128 (pcg. f)

```
!C
!C-- PADDING
!C

      icou= 0
      do icel= 1, ICELTOT
        do k = 1, 6
          if (itemLU(k, icel). eq. 0) then
            icou= icou + 1
            itemLU(k, icel)= icou + ICELTOT
            if (icou. eq. N2) icou= 0
          endif
        enddo
      enddo
```



solve_PCG (1/2)

```
!C
!C +-----+
!C | INIT |
!C +-----+
!C====

      allocate (W(N+N2, 4))

 !$omp parallel do private(i)
   do i= 1, N
     X(i) = 0.d0
     W(i, 2)= 0.0D0
     W(i, 3)= 0.0D0
     W(i, DD)= 1. d0/D(i)
   enddo

 !$omp parallel do private(i)
   do i= N+1, N+N2
     X(i) = 0.d0
     W(i, 2)= 0.0D0
     W(i, 3)= 0.0D0
     W(i, DD)= 1. d0/D(i)
   enddo
```

solve_PCG (2/2)

```
!C
!C +-----+
!C | {q} = [A] {p} |
!C +-----+
!C===
!
!$omp parallel do private(i, VAL, k)
    do i= 1, N
        VAL= D(i)*W(i, P)
        do k= 1, 6
            VAL= VAL + AMAT(k, i)*W(itemLU(k, i), P)
        enddo
        W(i, Q)= VAL
    enddo
!C===
!
```

Program by “ELL”: src-c2

```
>$ cd /work/gt89/t89XYZ/ompw  
>$ cd run  
  
<modify "INPUT.DAT", "f2_48.sh">  
  
>$ pbsub f2_48.sh
```

f2_48.sh

```
#!/bin/sh
#PJM -N "f2_48"
#PJM -L rscgrp=lecture9-o
#PJM -L node=1
#PJM --omp thread=48
#PJM -L elapse=00:15:00
#PJM -g
#PJM -j
#PJM -e
#PJM -o f2_48.lst
```

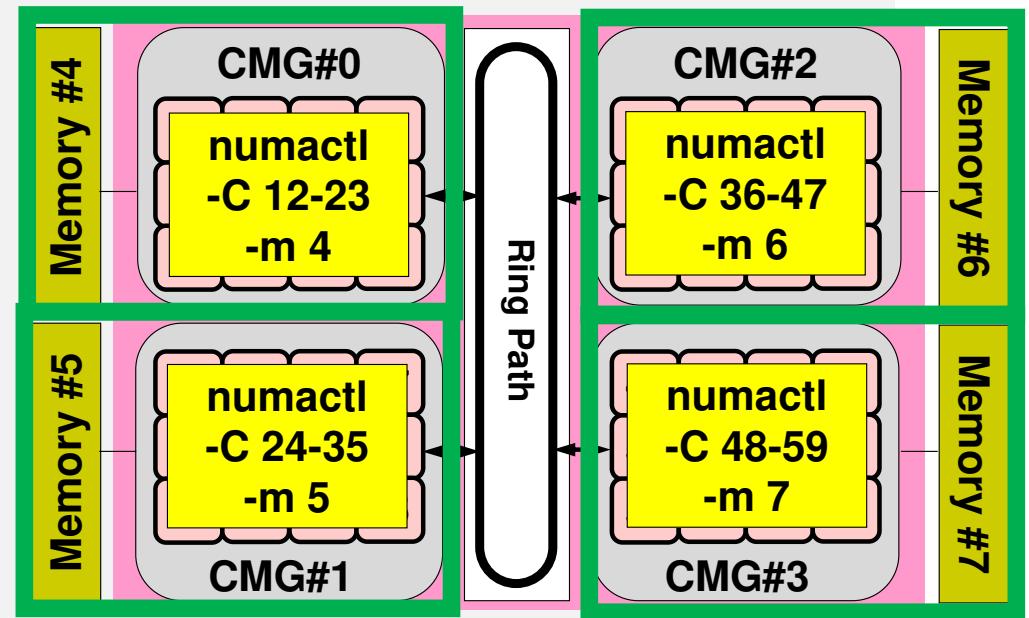
module load fj

export OMP_NUM_THREADS=48

export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

numactl	./solf2

numactl	-C 12-59 -m 4-7 ./solf2
numactl	-C 12-59 -m 4-7 ./solf2
numactl	-C 12-59 -m 4-7 ./solf2
numactl	-C 12-59 -m 4-7 ./solf2
numactl	-C 12-59 -m 4-7 ./solf2



Results (1/2): Fortran

```
[t00XYZ@wisteria01 run]$ cat f12.1st
```

```
1      8.958216E+00
101    8.313496E+00
201    2.090443E+00
301    3.811029E-01
401    3.769653E-02
501    9.429978E-04
601    4.940783E-05
701    1.888611E-06
801    2.243179E-08
826    9.818026E-09
5.275418E+00 sec. (solver)
```

```
##ANSWER    2097152    1.459831E+04
```

```
1      8.958216E+00
101    8.313496E+00
201    2.090443E+00
301    3.811029E-01
401    3.769653E-02
501    9.429978E-04
601    4.940783E-05
701    1.888611E-06
801    2.243179E-08
826    9.818026E-09
5.270398E+00 sec. (solver)
```

```
##ANSWER    2097152    1.459831E+04
```

Results (1/2): Best for 5 Measurements

```
[XYZ@wisteria01 run]$ grep "(sol" f1_48.1st
```

```
1.480524E+00 sec. (solver)  
1.501454E+00 sec. (solver)  
1.441297E+00 sec. (solver) numactl ./solf1  
1.483405E+00 sec. (solver)  
1.481864E+00 sec. (solver)
```

```
1.475129E+00 sec. (solver)  
1.483695E+00 sec. (solver)  
1.485036E+00 sec. (solver)  
1.502549E+00 sec. (solver)  
1.487192E+00 sec. (solver)
```

```
[XYZ@wisteria01 run]$ grep "(sol" f2_48.1st
```

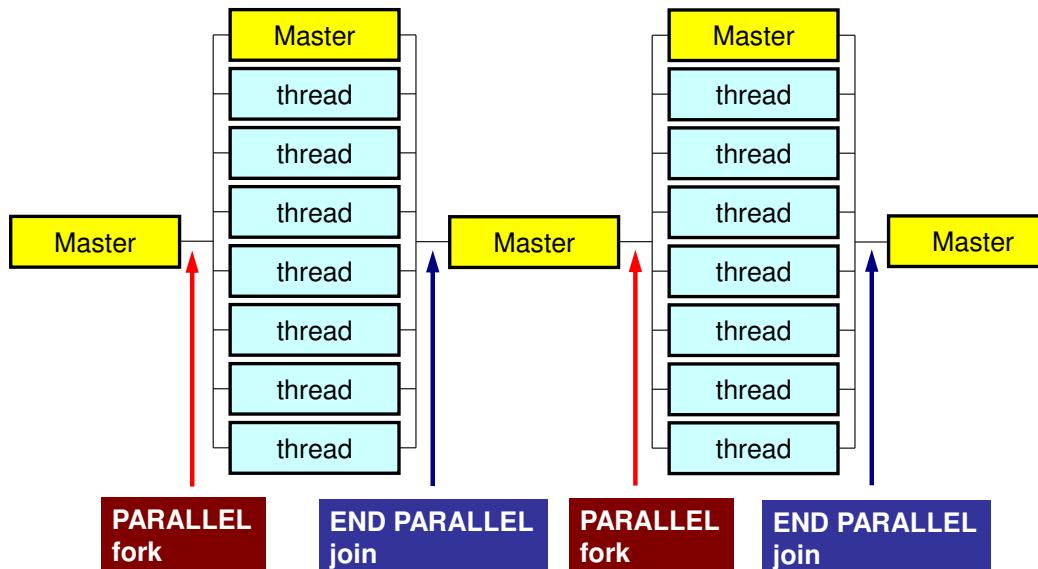
```
7.713702E-01 sec. (solver)  
7.568300E-01 sec. (solver)  
7.328739E-01 sec. (solver) numactl ./solf2  
7.826090E-01 sec. (solver)  
7.884219E-01 sec. (solver)
```

```
7.546160E-01 sec. (solver)  
7.937970E-01 sec. (solver)  
7.403760E-01 sec. (solver)  
7.745121E-01 sec. (solver)  
7.862871E-01 sec. (solver)
```

- OpenMP
- Login to Wisteria/BDEC-01
- Parallel Code by OpenMP (0): up to 12 cores
- Parallel Code by OpenMP (1): First Touch
- Parallel Code by OpenMP (2): +ELL
- **Parallel Code by OpenMP (3): reduced omp-parallel**
- **Parallel Code by OpenMP (4): Further Optimization (Fortran only)**

omp parallel (do)

- “omp parallel-omp end parallel” = “fork-join”
- If you have many loops, these “fork-join’s” cause overheads
- omp parallel + omp do/omp for



```
#pragma omp parallel ...
```

```
#pragma omp for {
```

...

```
#pragma omp for {
```

```
!$omp parallel ...
```

```
!$omp do  
do i= 1, N
```

...

```
!$omp do  
do i= 1, N
```

...

```
!$omp end parallel required
```

!\$omp parallel do: Fork-Join

```

!$omp parallel do private(i, VAL, k)
do i= 1, N
    VAL= D(i)*W(i, P)
    do k= indexLU(i-1)+1, indexLU(i)
        VAL= VAL + AMAT(k)*W(itemLU(k), P)
    enddo
    W(i, Q)= VAL
enddo

C1= 0. d0
!$omp parallel do private(i) reduction(+:C1)
do i= 1, N
    C1= C1 + W(i, P)*W(i, Q)
enddo

ALPHA= RHO / C1

!$omp parallel do private(i)
do i= 1, N
    X(i) = X(i) + ALPHA * W(i, P)
    W(i, R)= W(i, R) - ALPHA * W(i, Q)
enddo

DNRM2= 0. d0
!$omp parallel do private(i) reduction(+:DNRM2)
do i= 1, N
    DNRM2= DNRM2 + W(i, R)**2
enddo

ERR = dsqrt(DNRM2/BNRM2)...

```

Compute $r^{(0)} = b - [A]x^{(0)}$

for $i = 1, 2, \dots$

solve $[M]z^{(i-1)} = r^{(i-1)}$

$\rho_{i-1} = r^{(i-1)} \cdot z^{(i-1)}$

if $i = 1$

$p^{(1)} = z^{(0)}$

else

$\beta_{i-1} = \rho_{i-1}/\rho_{i-2}$

$p^{(i)} = z^{(i-1)} + \beta_{i-1} p^{(i-1)}$

endif

$q^{(i)} = [A]p^{(i)}$

$\alpha_i = \rho_{i-1}/p^{(i)} q^{(i)}$

$x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$

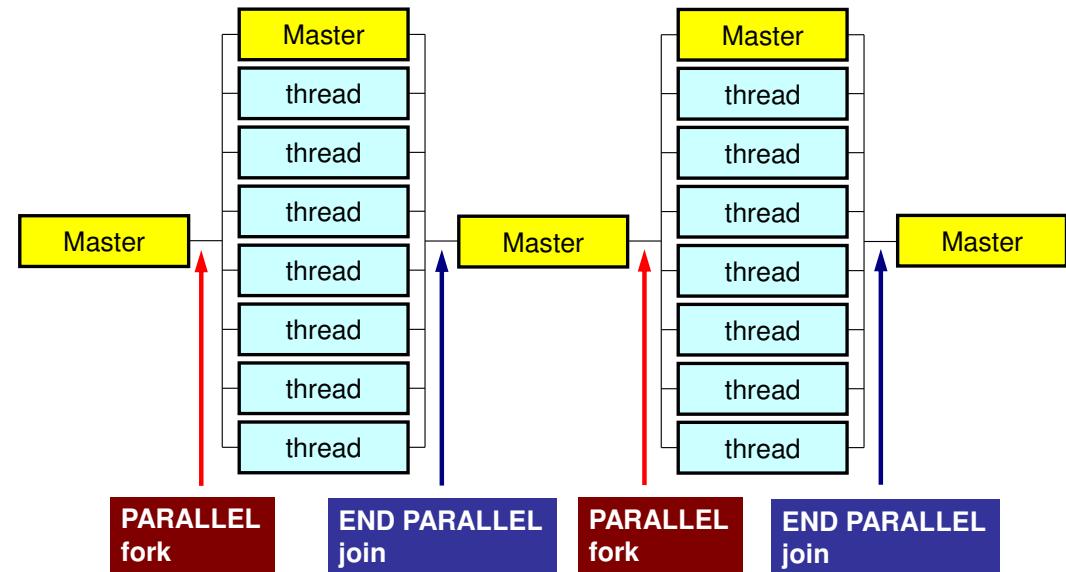
$r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$

check convergence $|r|$

end

Strategy for Further Optimization

- src-c3, src-f3
 - Only 1 omp-parallel in each iteration
- src-f4 (Fortran Only)
 - Only 1 !\$omp parallel during PCG
 - NO !\$omp do
 - Operations for reduction in dot-product's are NOT parallelized



Further Optimization

src-f3

```
do L= 1, ITR
!$omp parallel private (i, k, VAL)
(...)
!$omp end parallel
enddo
```

900 continue

ITR= L

deallocate (W)

return
end

src-f4

```
!$omp parallel private (...)
```

```
do L= 1, ITR
(...)
enddo
```

900 continue

```
ITR= L
!$omp end parallel
```

deallocate (W)

return
end

src_f3 (1/2)

```

ITR= N
Stime= omp_get_wtime()

do L= 1, ITR
!$omp parallel private(i, k, VAL)
!$omp do
  do i= 1, N
    W(i, Z)= W(i, R)*W(i, DD)
  enddo

  RH0= 0. d0
!$omp do reduction(+:RH0)
  do i= 1, N
    RH0= RH0 + W(i, R)*W(i, Z)
  enddo

  if ( L.eq. 1 ) then
!$omp do
    do i= 1, N
      W(i, P)= W(i, Z)
    enddo
  else
    BETA= RH0 / RH01
!$omp do
    do i= 1, N
      W(i, P)= W(i, Z) + BETA*W(i, P)
    enddo
  endif

```

Compute $r^{(0)} = b - [A]x^{(0)}$

for $i = 1, 2, \dots$

solve $[M]z^{(i-1)} = r^{(i-1)}$

$\rho_{i-1} = r^{(i-1)} \cdot z^{(i-1)}$

if $i=1$

$p^{(1)} = z^{(0)}$

else

$\beta_{i-1} = \rho_{i-1}/\rho_{i-2}$

$p^{(i)} = z^{(i-1)} + \beta_{i-1} p^{(i-1)}$

endif

$q^{(i)} = [A]p^{(i)}$

$\alpha_i = \rho_{i-1}/p^{(i)} q^{(i)}$

$x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$

$r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$

check convergence $|r|$

end

src_f3 (2/2)

```

!$omp do
do i= 1, N
    VAL= D(i)*W(i, P)
    do k= indexLU(i-1)+1, indexLU(i)
        VAL= VAL + AMAT(k)*W(itemLU(k), P)
    enddo
    W(i, Q)= VAL
enddo

C1= 0. d0
!$omp do reduction(+:C1)
do i= 1, N
    C1= C1 + W(i, P)*W(i, Q)
enddo

ALPHA= RHO / C1

!$omp do
do i= 1, N
    X(i)= X(i) + ALPHA * W(i, P)
    W(i, R)= W(i, R) - ALPHA * W(i, Q)
enddo

DNRM2= 0. d0
!$omp do reduction(+:DNRM2)
do i= 1, N
    DNRM2= DNRM2 + W(i, R)**2
enddo

!$omp end parallel
ERR = dsqrt(DNRM2/BNRM2)...

```

Compute $r^{(0)} = b - [A]x^{(0)}$

for $i = 1, 2, \dots$

solve $[M]z^{(i-1)} = r^{(i-1)}$

$\rho_{i-1} = r^{(i-1)} \cdot z^{(i-1)}$

if $i=1$

$p^{(1)} = z^{(0)}$

else

$\beta_{i-1} = \rho_{i-1}/\rho_{i-2}$

$p^{(i)} = z^{(i-1)} + \beta_{i-1} p^{(i-1)}$

endif

$q^{(i)} = [A]p^{(i)}$

$\alpha_i = \rho_{i-1}/p^{(i)} q^{(i)}$

$x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$

$r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$

check convergence $|r|$

end

Program for “src-f3”

```
>$ cd /work/gt89/t89XYZ/ompw  
>$ cd run  
  
<modify "INPUT.DAT", "f3_48.sh">  
  
>$ pbsub f3_48.sh
```

f3_48.sh

```
#!/bin/sh
#PJM -N "f3_48"
#PJM -L rscgrp=lecture9-o
#PJM -L node=1
#PJM --omp thread=48
#PJM -L elapse=00:15:00
#PJM -g
#PJM -j
#PJM -e
#PJM -o f3_48.lst
```

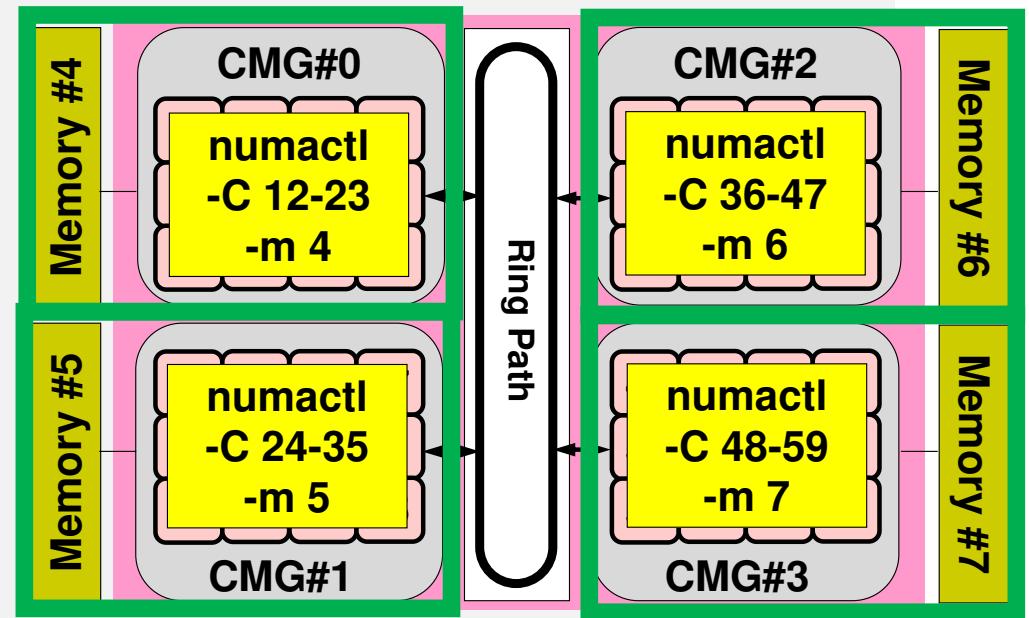
module load fj

export OMP_NUM_THREADS=48

export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

numactl	./solf3

numactl	-C 12-59 -m 4-7 ./solf3
numactl	-C 12-59 -m 4-7 ./solf3
numactl	-C 12-59 -m 4-7 ./solf3
numactl	-C 12-59 -m 4-7 ./solf3
numactl	-C 12-59 -m 4-7 ./solf3



C Language : trad

```
>$ cd /work/gt89/t89XYZ/ompw
>$ cd run

<modify "INPUT.DAT", "c48org.sh">

>$ pbsub c48org.sh
```

c48org.sh

```
#!/bin/sh
#PJM -N "cx48"
#PJM -L rscgrp=lecture9-o
#PJM -L node=1
#PJM --omp thread=48
#PJM -L elapse=00:15:00
#PJM -g gt89
#PJM -j
#PJM -e err
#PJM -o c48org_160.lst

module load fj
export OMP_NUM_THREADS=48
export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

numactl ./solc0org
numactl ./solc0org
numactl ./solc0org
numactl ./solc0org
numactl ./solc0org
...
...
```

Time for PCG, N=128³, 48 Threads

N= 2,097,152

	Fortran	C (clang)	C (trad)
src0	1.671	1.564	2.354
src1 (First Touch)	1.480	1.122	1.720
src2 (+ ELL)	0.747	0.809	1.127
src3 (+ reduced “omp-parallel”)	0.707	0.834	0.854

Time for PCG, N=160³, 48 Threads

N= 4,096,000

	Fortran	C (clang)	C (trad)
src0	3.610	3.484	4.067
src1 (First Touch)	2.993	2.228	3.425
src2 (+ ELL)	1.534	1.690	2.340
src3 (+ reduced “omp-parallel”)	1.556	1.693	1.742

Time for PCG, N=200³, 48 Threads

N= 8,000,000

	Fortran	C (clang)	C (trad)
src0	7.666	8.321	9.397
src1 (First Touch)	6.952	5.102	8.008
src2 (+ ELL)	3.421	3.910	5.381
src3 (+ reduced “omp-parallel”)	3.440	3.920	3.824

Time for PCG, N=256³, 48 Threads

N= 16,777,216

	Fortran	C (clang)	C (trad)
src0	34.308	24.772	25.547
src1 (First Touch)	32.202	22.172	23.814
src2 (+ ELL)	8.916	10.761	14.566
src3 (+ reduced “omp-parallel”)	8.915	10.764	10.415

src_f4 (1/5)

parallel computing by OpenMP

```

module solver_PCG
    contains
!C
!C*** solve_PCG
!C
    subroutine solve_PCG
&        ( N, NPLU, indexLU, itemLU, D, B, X, AMAT, EPS, ITR, IER) &
use omp_lib
implicit REAL*8 (A-H, O-Z)

real(kind=8), dimension(N)      :: D
real(kind=8), dimension(N)      :: B
real(kind=8), dimension(N)      :: X
real(kind=8), dimension(NPLU)   :: AMAT

integer, dimension(0:N)         :: indexLU
integer, dimension(NPLU)        :: itemLU

real(kind=8), dimension(:, :), allocatable :: W
integer(kind=4), dimension(: ), allocatable :: SMPindex

integer, parameter :: R= 1
integer, parameter :: Z= 2
integer, parameter :: Q= 2
integer, parameter :: P= 3
integer, parameter :: DD= 4

real(kind=8), dimension(:), allocatable :: W_RHO, W_C1, W_DNRM2

```

```

allocate (W(N+N2, 4))

!$omp parallel do private(i)
do i= 1, N
  X(i) = 0. d0
  W(i, 2)= 0. ODO
  W(i, 3)= 0. ODO
  W(i, DD)= 1. d0/D(i)
enddo

!$omp parallel do private(i)
do i= N+1, N+N2
  X(i) = 0. d0
  W(i, 2)= 0. ODO
  W(i, 3)= 0. ODO
  W(i, DD)= 1. d0/D(i)
enddo

```

```

!$omp parallel
PEsmpTOT= omp_get_num_threads()
!$omp end parallel

```

```

allocate (SMPindex(0:PEsmpTOT))
SMPindex(0)= 0
m = N/PEsmpTOT
nr= N - PEsmpTOT*m
do ip= 1, PEsmpTOT
  SMPindex(ip)= m
  if (ip.le.nr) SMPindex(ip)= m+1
enddo
do ip= 1, PEsmpTOT
  SMPindex(ip)= SMPindex(ip)+SMPindex(ip-1)
enddo

```

```
allocate (W_RHO(PEsmpTOT), W_C1(PEsmpTOT), W_DNRM2(PEsmpTOT))
```

src_f4 (2/5)

PEsmpTOT: Total Number of Threads

SMPindex(0:PEsmpTOT) : Element# for each thread

内積用

src_f4 (3/5)

```
!$omp parallel do private(i, VAL, k)
do i= 1, N
    VAL= D(i)*X(i)
    do k= 1, 6
        VAL= VAL + AMAT(k, i)*X(itemLU(k, i))
    enddo
    W(i, R)= B(i) - VAL
enddo

BNRM2= 0.0D0
!$omp parallel do private(i) reduction(+:BNRM2)
do i= 1, N
    BNRM2 = BNRM2 + B(i) **2
enddo
```

```

ITR= N
Stime= omp_get_wtime()

 !$omp parallel private(L, ip, ip1, ip2, i, k, VAL)
 !$omp& private(RHO, BETA, RH01, C1, ALPHA, DNRM2)
 do L= 1, ITR

     ip = omp_get_thread_num() + 1
     ip1= SMPIndex(ip-1)+1
     ip2= SMPIndex(ip)

 !$omp simd
 do i= ip1, ip2
     W(i, Z)= W(i, R)*W(i, DD)
 enddo

     W_RHO(ip)= 0.0d0
 !$omp simd
 do i= ip1, ip2
     W_RHO(ip)= W_RHO(ip) + W(i, R)*W(i, Z)
 enddo
 !$omp barrier
 RHO= 0.0d0
 !$omp simd
 do i = 1, PEsmptot
     RHO= RHO + W_RHO(i)
 enddo

 if ( L.eq. 1 ) then
 !$omp simd
 do i= ip1, ip2
     W(i, P)= W(i, Z)
 enddo
 else
     BETA= RHO / RH01
 !$omp simd
 do i= ip1, ip2
     W(i, P)= W(i, Z) + BETA*W(i, P)
 enddo
 endif
 !$omp barrier

```

src_f4 (4/5)

Compute $r^{(0)} = b - [A]x^{(0)}$

for $i = 1, 2, \dots$

solve $[M]z^{(i-1)} = r^{(i-1)}$

$\rho_{i-1} = r^{(i-1)} \cdot z^{(i-1)}$

if $i=1$

$p^{(1)} = z^{(0)}$

else

$\beta_{i-1} = \rho_{i-1}/\rho_{i-2}$

$p^{(i)} = z^{(i-1)} + \beta_{i-1} p^{(i-1)}$

endif

$q^{(i)} = [A]p^{(i)}$

$\alpha_i = \rho_{i-1}/p^{(i)} q^{(i)}$

$x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$

$r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$

check convergence $|r|$

end

```

do i= ip1, ip2
  VAL= D(i)*W(i, P)
  do k= 1, 6
    VAL= VAL + AMAT(k, i)*W(itemLU(k, i), P)
  enddo
  W(i, Q)= VAL
enddo

W_C1(ip)= 0.0d0
 !$omp simd
 do i= ip1, ip2
   W_C1(ip)= W_C1(ip) + W(i, P)*W(i, Q)
 enddo
 !$omp barrier
 C1= 0. d0
 !$omp simd
 do i = 1, PEsmptOT
   C1= C1 + W_C1(i)
 enddo
 ALPHA= RHO / C1

 !$omp simd
 do i= ip1, ip2
   X(i) = X(i) + ALPHA * W(i, P)
   W(i, R)= W(i, R) - ALPHA * W(i, Q)
 enddo

 W_DNRM2(ip)= 0.0d0
 !$omp simd
 do i= ip1, ip2
   W_DNRM2(ip)= W_DNRM2(ip) + W(i, R)**2
 enddo
 !$omp barrier
 DNRM2= 0. d0
 !$omp simd
 do i = 1, PEsmptOT
   DNRM2= DNRM2 + W_DNRM2(i)
 enddo
 ERR = dsqrt(DNRM2/BNRM2)...

```

src_f4 (5/5)

Compute $r^{(0)} = b - [A]x^{(0)}$

for $i = 1, 2, \dots$

solve $[M]z^{(i-1)} = r^{(i-1)}$

$\rho_{i-1} = r^{(i-1)} \cdot z^{(i-1)}$

if $i = 1$

$p^{(1)} = z^{(0)}$

else

$\beta_{i-1} = \rho_{i-1} / \rho_{i-2}$

$p^{(i)} = z^{(i-1)} + \beta_{i-1} p^{(i-1)}$

endif

$q^{(i)} = [A]p^{(i)}$

$\alpha_i = \rho_{i-1} / p^{(i)} q^{(i)}$

$x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$

$r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$

check convergence $|r|$

end

Program for “src4”

```
>$ cd /work/gt89/t89XYZ/ompw  
>$ cd run  
  
<modify "INPUT.DAT", "f4_48.sh">  
  
>$ pbsub f4_48.sh
```

f4_48.sh

```
#!/bin/sh
#PJM -N "f4_48"
#PJM -L rscgrp=lecture9-o
#PJM -L node=1
#PJM --omp thread=48
#PJM -L elapse=00:15:00
#PJM -g
#PJM -j
#PJM -e
#PJM -o f4_48.lst
```

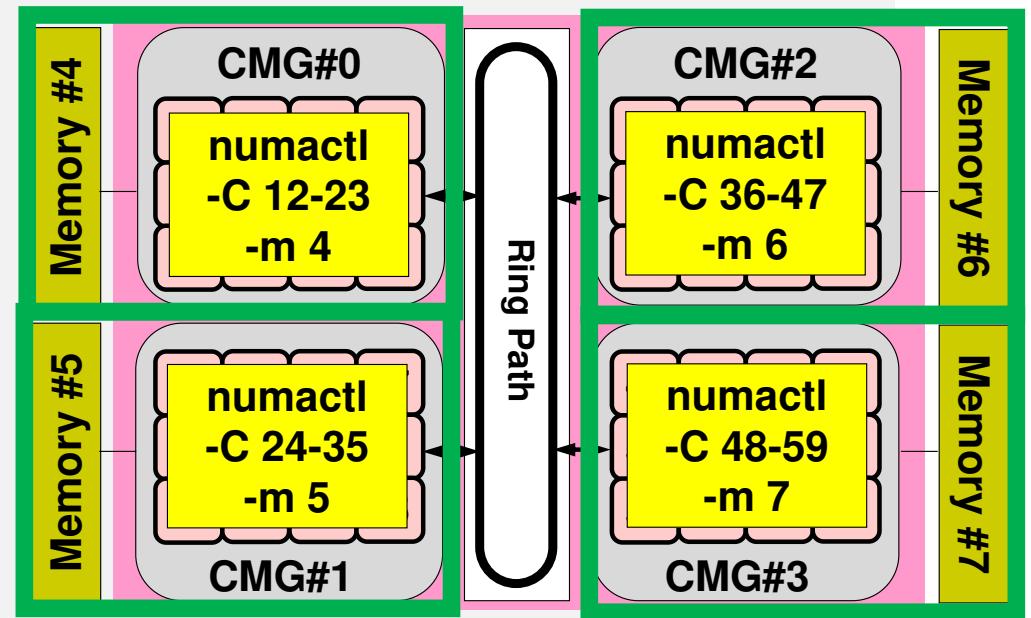
module load fj

export OMP_NUM_THREADS=48

export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

numactl	./solf4

numactl	-C 12-59 -m 4-7 ./solf4
numactl	-C 12-59 -m 4-7 ./solf4
numactl	-C 12-59 -m 4-7 ./solf4
numactl	-C 12-59 -m 4-7 ./solf4
numactl	-C 12-59 -m 4-7 ./solf4



Time for PCG : N=128³, 48 Threads

N= 2,097,152

	Fortran	C (clang)	C (trad)
src0	1.671	1.564	2.354
src1 (First Touch)	1.480	1.122	1.720
src2 (+ ELL)	0.747	0.809	1.127
src3 (+ reduced “omp-parallel”)	0.707	0.834	0.854
src3b (+ clang loop unroll_count)	-	0.764	-
src4 (+ No OMP-DO/Reduction)	0.676	-	-

src_c3b (5/5)

src-c3

```
#pragma omp for
for(i=0; i<N; i++) {
    VAL = D[i] * W[P][i];
    for(j=0; j<6; j++) {
        VAL += AMAT[6*i+j]*W[P][itemLU[6*i+j]];
    }
    W[Q][i] = VAL;
}
```

src-c3b: clang only

```
#pragma omp for
#pragma clang loop unroll_count(8)

for(i=0; i<N; i++) {
    VAL = D[i] * W[P][i];
    for(j=0; j<6; j++) {
        VAL += AMAT[6*i+j]*W[P][itemLU[6*i+j]];
    }
    W[Q][i] = VAL;
}
```

Compute $r^{(0)} = b - [A]x^{(0)}$

```
for i= 1, 2, ...
    solve [M] z(i-1) = r(i-1)
    ρi-1 = r(i-1) z(i-1)
    if i=1
        p(1) = z(0)
    else
        βi-1 = ρi-1/ρi-2
        p(i) = z(i-1) + βi-1 p(i-1)
    endif
    q(i) = [A]p(i)
    αi = ρi-1/p(i)q(i)
    x(i) = x(i-1) + αip(i)
    r(i) = r(i-1) - αiq(i)
    check convergence |r|
end
```

Time for PCG : N=160³, 48 Threads

N= 4,096,000

	Fortran	C (clang)	C (trad)
src0	3.610	3.484	4.067
src1 (First Touch)	2.993	2.228	3.425
src2 (+ ELL)	1.534	1.690	2.340
src3 (+ reduced “omp-parallel”)	1.556	1.693	1.742
src3b (+ clang loop unroll_count)	-	1.586	-
src4 (+ No OMP-DO/Reduction)	1.435	-	-

Time for PCG : N=200³, 48 Threads

N= 8,000,000

	Fortran	C (clang)	C (trad)
src0	7.666	8.321	9.397
src1 (First Touch)	6.952	5.102	8.008
src2 (+ ELL)	3.421	3.910	5.381
src3 (+ reduced “omp-parallel”)	3.440	3.920	3.824
src3b (+ clang loop unroll_count)	-	3.624	-
src4 (+ No OMP-DO/Reduction)	3.276	-	-

Time for PCG : N=256³, 48 Threads

N= 16,777,216

	Fortran	C (clang)	C (trad)
src0	34.308	24.772	25.547
src1 (First Touch)	32.202	22.172	23.814
src2 (+ ELL)	8.916	10.761	14.566
src3 (+ reduced “omp-parallel”)	8.915	10.764	10.415
src3b (+ clang loop unroll_count)	-	10.003	-
src4 (+ No OMP-DO/Reduction)	8.620	-	-

Exercises

- Problem size (NX, NY, NZ)
- Thread # (OMP_NUM_THREADS: 1-48)
- Various Types of Implementation
- Profiling

Performance Evaluation: Profiler

- Specify the measuring unit in the program (multiple settings are possible)
- 17 runs needed for each case
- Performance of Computation & Memory, Power Consumption
- Excel Macro File
 - https://www.dropbox.com/s/kat9ny5aoxp7cqm/cpu_pa_report.xls?dl=0

Profiling (1/4)

Specifying the measuring unit in the program, No additional options for compiling: solver_PCG.c/f

```
#include "fj_tool/fapp.h"
```

```
fapp_start ("CG", 1, 0); ←
Stime = omp_get_wtime();
```

```
for (L=0; L<(*ITR); L++) {
```

```
...
if (ERR < EPS) {
    *IER = 0;
    goto N900;
} else {
    RH01 = RHO;
}
```

```
}
```

```
*IER = 1;
```

```
N900:
```

```
Etime = omp_get_wtime();
fapp_stop ("CG", 1, 0); ←
```

```
return 0;
```

```
}
```

```
call fapp_start ("CG", 1, 0) ←
Stime = omp_get_wtime()
```

```
do L= 1, ITR
```

```
...
if (ERR .lt. EPS) then
    IER = 0
    goto 900
else
    RH01 = RHO
endif
```

```
enddo
```

```
IER = 1
```

```
900 continue
```

```
Etime= omp_get_wtime()
call fapp_stop ("CG", 1, 0) ←
```

```
return
end
```

Profiling (2/4)

Running on Wisteria/BDEC-01 (Odyssey)

```
>$ cd /work/gt89/t89XYZ/ompw
>$ cd run

<modify "fapp.sh", "data.sh">

>$ pbsub fapp.sh

(after finishing)

>$ pjsbu data.sh

>$ ls pa*.csv
    pa1.csv ... pa17.csv
```

Profiling (3/4)

Each Directory (e.g. repo01 etc.) should be empty before running the code

fapp.sh

```
#!/bin/sh
#PJM -N "fapp"
#PJM -L rscgrp=lecture9-o
#PJM -L node=1
#PJM -omp thread=48
#PJM -L elapse=00:15:00
#PJM -g gt89
#PJM -j
#PJM -e err
#PJM -o test.lst

module load fj
export OMP_NUM_THREADS=48
export XOS_MMM_L_PAGING_POLICY=demand:demand:demand

fapp -C -d ./repo01 -Hevent=pa1 ./solf2
fapp -C -d ./repo02 -Hevent=pa2 ./solf2
fapp -C -d ./repo03 -Hevent=pa3 ./solf2
fapp -C -d ./repo04 -Hevent=pa4 ./solf2
fapp -C -d ./repo05 -Hevent=pa5 ./solf2
fapp -C -d ./repo06 -Hevent=pa6 ./solf2
fapp -C -d ./repo07 -Hevent=pa7 ./solf2
fapp -C -d ./repo08 -Hevent=pa8 ./solf2
fapp -C -d ./repo09 -Hevent=pa9 ./solf2
fapp -C -d ./repo10 -Hevent=pa10 ./solf2
fapp -C -d ./repo11 -Hevent=pa11 ./solf2
fapp -C -d ./repo12 -Hevent=pa12 ./solf2
fapp -C -d ./repo13 -Hevent=pa13 ./solf2
fapp -C -d ./repo14 -Hevent=pa14 ./solf2
fapp -C -d ./repo15 -Hevent=pa15 ./solf2
fapp -C -d ./repo16 -Hevent=pa16 ./solf2
fapp -C -d ./repo17 -Hevent=pa17 ./solf2
```

data.sh

```
#!/bin/sh
#PJM -N "data"
#PJM -L rscgrp=lecture9-o
#PJM -L node=1
#PJM --mpi proc=1
#PJM -L elapse=00:15:00
#PJM -g gt89
#PJM -j
#PJM -e err
#PJM -o data.lst

module load fj
module load fjmpi

fapp -A -d ./repo01 -Icpupa,mpi -tcsv -o pa1.csv
fapp -A -d ./repo02 -Icpupa,mpi -tcsv -o pa2.csv
fapp -A -d ./repo03 -Icpupa,mpi -tcsv -o pa3.csv
fapp -A -d ./repo04 -Icpupa,mpi -tcsv -o pa4.csv
fapp -A -d ./repo05 -Icpupa,mpi -tcsv -o pa5.csv
fapp -A -d ./repo06 -Icpupa,mpi -tcsv -o pa6.csv
fapp -A -d ./repo07 -Icpupa,mpi -tcsv -o pa7.csv
fapp -A -d ./repo08 -Icpupa,mpi -tcsv -o pa8.csv
fapp -A -d ./repo09 -Icpupa,mpi -tcsv -o pa9.csv
fapp -A -d ./repo10 -Icpupa,mpi -tcsv -o pa10.csv
fapp -A -d ./repo11 -Icpupa,mpi -tcsv -o pa11.csv
fapp -A -d ./repo12 -Icpupa,mpi -tcsv -o pa12.csv
fapp -A -d ./repo13 -Icpupa,mpi -tcsv -o pa13.csv
fapp -A -d ./repo14 -Icpupa,mpi -tcsv -o pa14.csv
fapp -A -d ./repo15 -Icpupa,mpi -tcsv -o pa15.csv
fapp -A -d ./repo16 -Icpupa,mpi -tcsv -o pa16.csv
fapp -A -d ./repo17 -Icpupa,mpi -tcsv -o pa17.csv
```

Profiling (4/4): Operations on PC

- Copying all “pa*.csv”’s to your PC

```
>$ scp t00xyz@wisteria.cc.u-tokyo.ac.jp:/work/gt00/t00xyz/ompw/run/pa*.csv .
```

- All “pa*.csv”’s and Macro for Excel should be in the same directory
- “Double Click” the Excel Macro
 - Just follow instructions
 - Please select “CG”, not “All”

Time for PCG: N=160³, 48 Threads

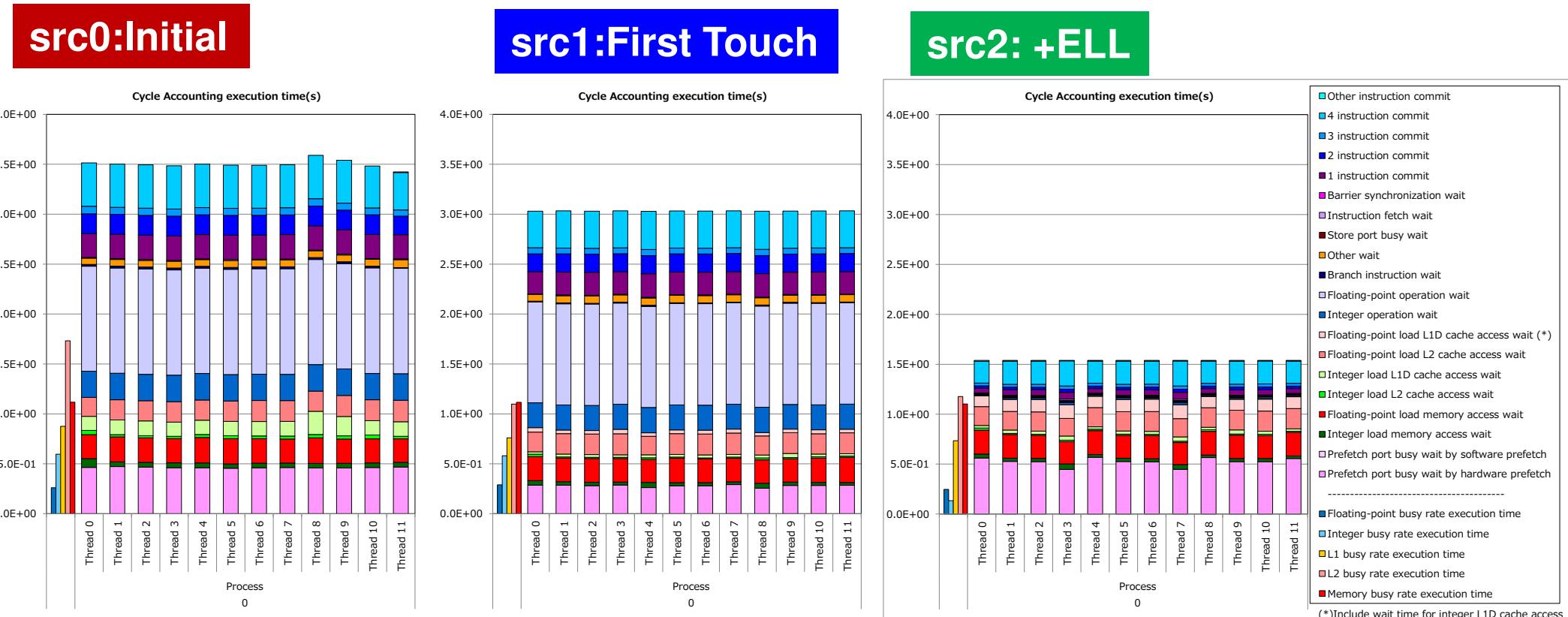
N= 4,096,000, for Each Node, Fortran

Optimization – Busy Memory – More Watt

	Time (sec.)	Peak Perf. Ratio (%)	SIMD Ratio (%)	Memory Throughput (%)	Instruction		Power (W)	
					Effective	Load/ Store	Core L2 Memory	Node
solf0	3.69	1.59	20.0	30.3	3.39×10^{11}	8.27×10^{10}	81.6 10.9 20.2	112.
solf1 (First Touch)	2.98	1.97	28.8	37.5	2.35×10^{11}	5.33×10^{10}	92.1 10.8 33.3	136.
solf2 (+ ELL)	1.58	3.73	49.7	70.0	1.19×10^{11}	4.17×10^{10}	104. 15.0 51.7	170.
solf3 (+ reduced “omp- parallel”)	1.58	3.72	48.4	69.8	1.22×10^{11}	4.10×10^{10}	101. 14.9 51.0	167.
solf4 (+ further optimization)	1.45	4.06	55.1	75.8	1.10×10^{11}	3.78×10^{10}	102. 15.8 56.5	174.

PCG, N=160³, 48 Threads

Fortran, Each CMG



Prefetch Port Busy Wait (H/W)
Integer Load L1D Cache Access Wait
Floating-Point Operation Wait
Instruction Commit

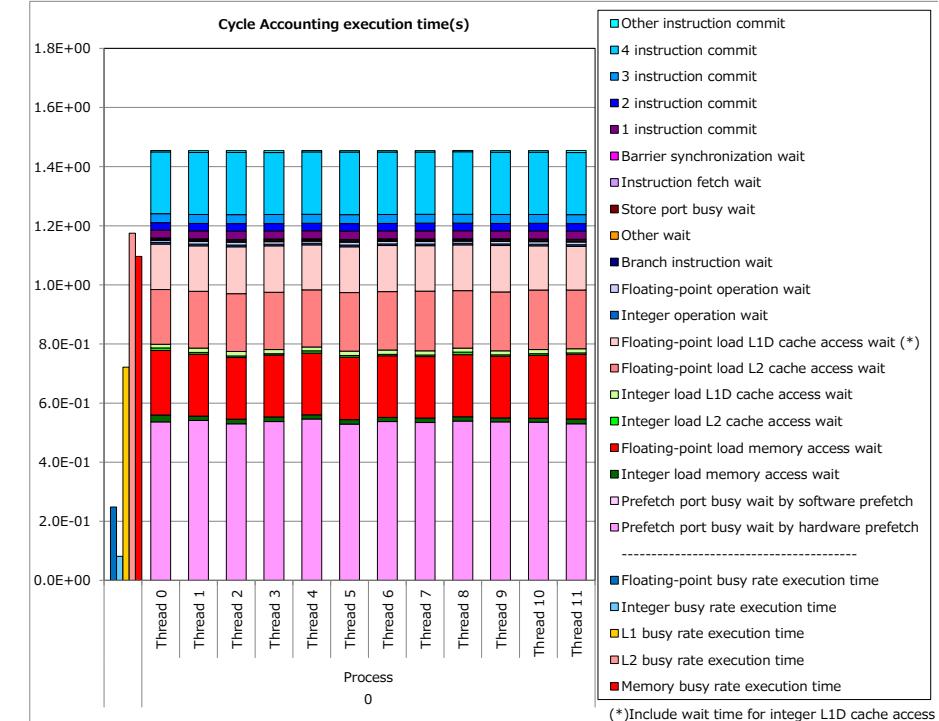
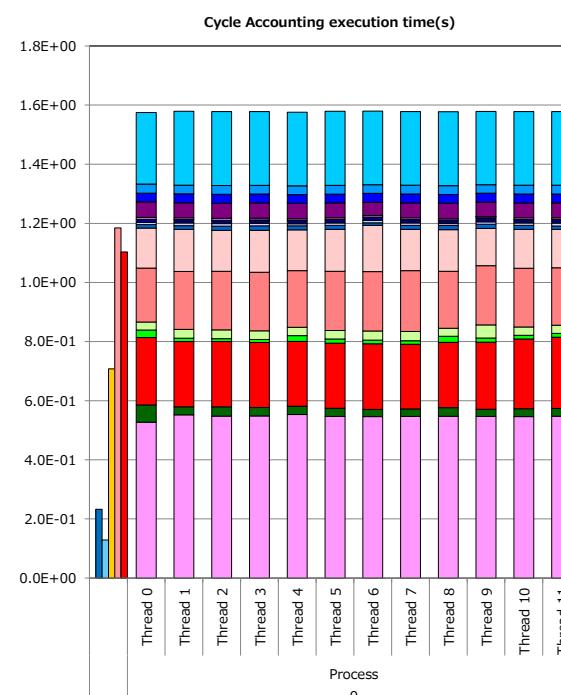
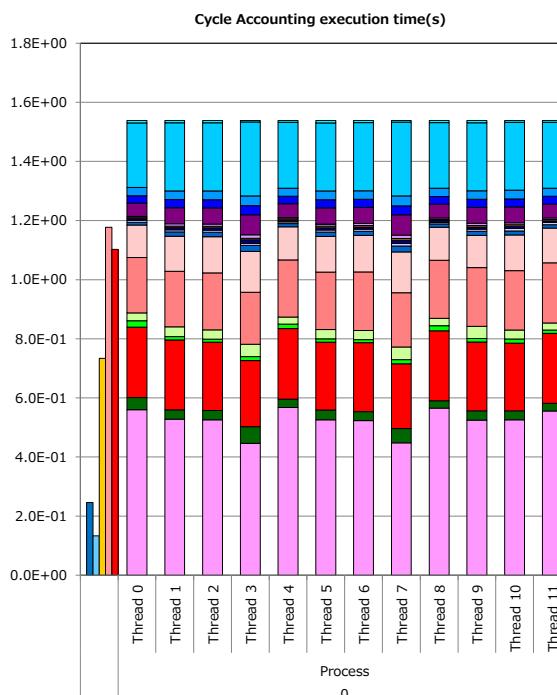
PCG, N=160³, 48 Threads

Fortran, Each CMG

src2:+ELL

src3

src4



- █ Prefetch Port Busy Wait (H/W)
- █ Floating-Point Load Memory Access Wait
- █ Floating-Point Load L2D Cache Access Wait
- █ Floating-Point Load L1D Cache Access Wait
- █ Instruction Commit

█ Other instruction commit
█ 4 instruction commit
█ 3 instruction commit
█ 2 instruction commit
█ 1 instruction commit
█ Barrier synchronization wait
█ Instruction fetch wait
█ Store port busy wait
█ Other wait
█ Branch instruction wait
█ Floating-point operation wait
█ Integer operation wait
█ Floating-point load L1D cache access wait (*)
█ Floating-point load L2D cache access wait
█ Integer load L1D cache access wait
█ Integer load L2D cache access wait
█ Floating-point load memory access wait
█ Integer load memory access wait
█ Prefetch port busy wait by software prefetch
█ Prefetch port busy wait by hardware prefetch
█ Floating-point busy rate execution time
█ Integer busy rate execution time
█ L1 busy rate execution time
█ L2 busy rate execution time
█ Memory busy rate execution time

(*)Include wait time for integer L1D cache access

Time for PCG: N=160³, 48 Threads

N= 4,096,000, for Each Node, C (clang)

Optimization – Busy Memory – More Watt

	Time (sec.)	Peak Perf. Ratio (%)	SIMD Ratio (%)	Memory Throughput (%)	Instruction		Power (W)	
					Effective	Load/ Store	Core L2 Memory	Node
solc0	3.53	0.90	2.95	31.7	4.83×10^{11}	1.65×10^{11}	103.	134.
solc1 (First Touch)	2.42	1.34	4.82	46.2	2.97×10^{11}	1.10×10^{11}	107.	153.
solc2 (+ ELL)	1.73	1.88	8.21	63.8	1.74×10^{11}	8.92×10^{10}	104.	163.
solc3 (+ reduced “omp- parallel”)	1.75	1.86	8.51	62.9	1.68×10^{11}	8.88×10^{10}	108.	167.
solc3b (+ clang loop unroll_count)	1.62	2.00	10.4	67.8	1.37×10^{11}	8.22×10^{10}	109.	168.

Time for PCG: N=160³, 48 Threads

N= 4,096,000, for Each Node, C (trad)

Optimization – Busy Memory – More Watt

	Time (sec.)	Peak Perf. Ratio (%)	SIMD Ratio (%)	Memory Throughput (%)	Instruction		Power (W)	
					Effective	Load/ Store	Core L2 Memory	Node
solc0org	4.14	2.66	27.3	27.1	3.78×10^{11}	7.40×10^{10}	91.3 11.3 18.4	121.
solc1org (First Touch)	3.52	3.13	37.4	31.8	2.75×10^{11}	6.14×10^{10}	90. 9.86 28.4	128.
solc2org (+ ELL)	2.34	4.70	58.1	47.1	1.63×10^{11}	4.77×10^{10}	95.1 11.7 37.6	144.
solc3org (+ reduced “omp- parallel”)	1.70	3.46	47.1	64.8	1.17×10^{11}	4.14×10^{10}	96.7 13.9 48.7	159.