

Introduction to Parallel Programming for Multicore/Manycore Clusters

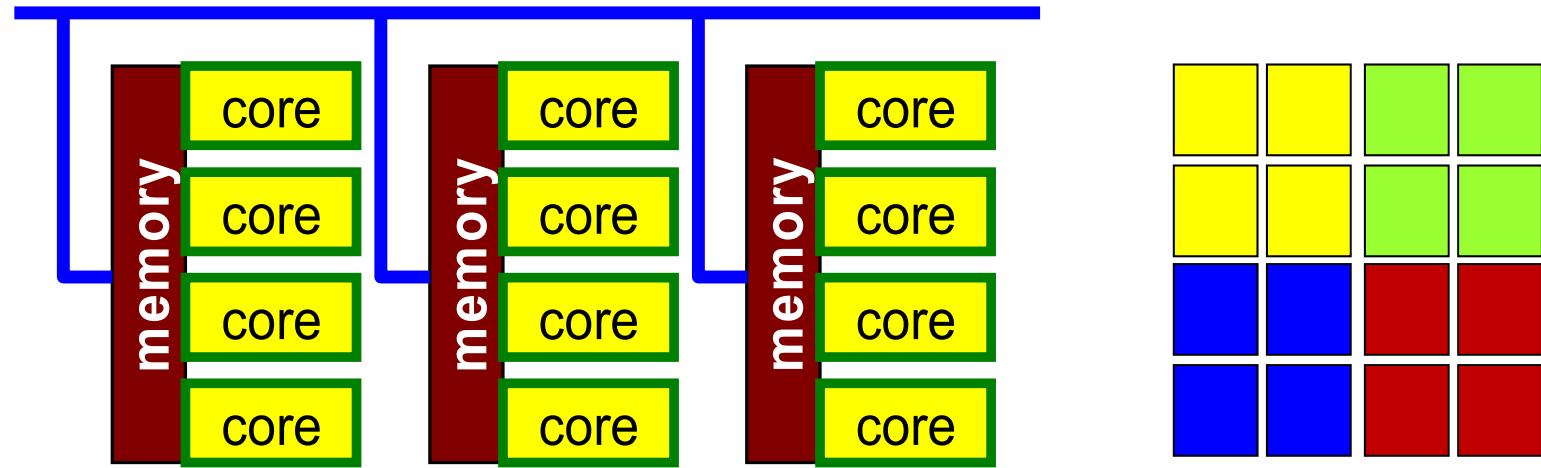
Part II-2: Parallel FVM using OpenMP

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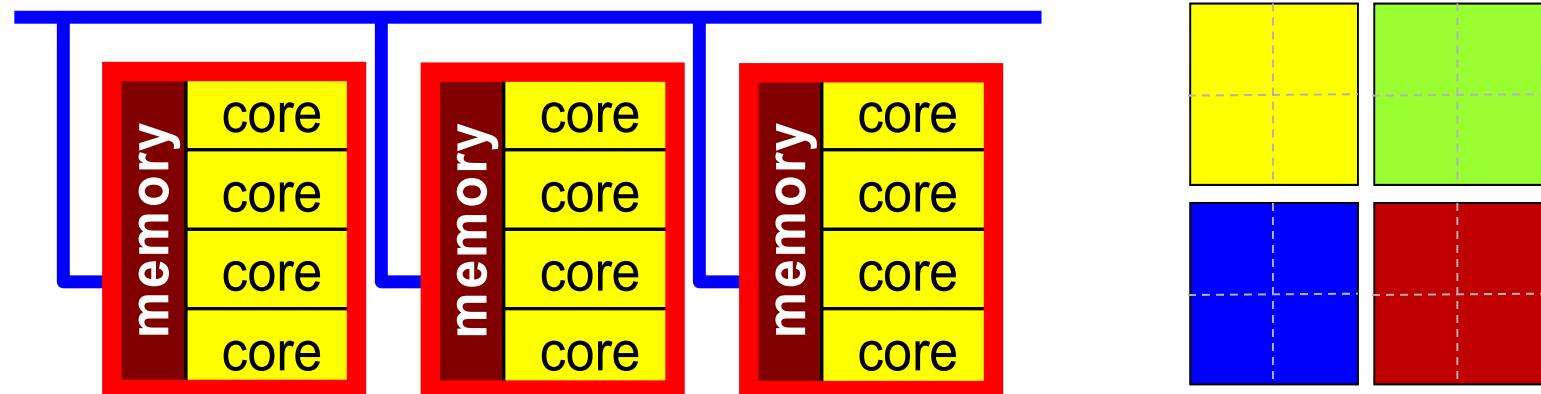
- OpenMP
- Parallel Version of the Code by OpenMP
- STREAM

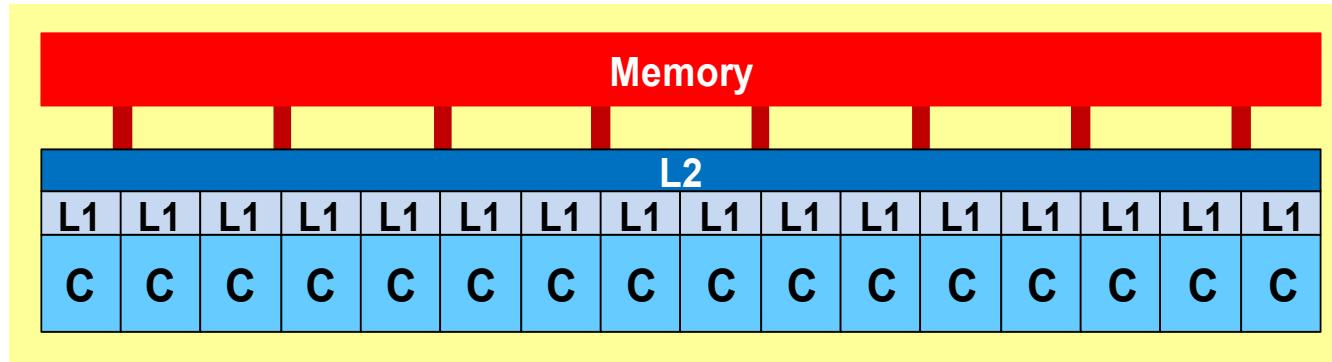
Flat MPI vs. Hybrid

Flat-MPI: Each Core -> Independent

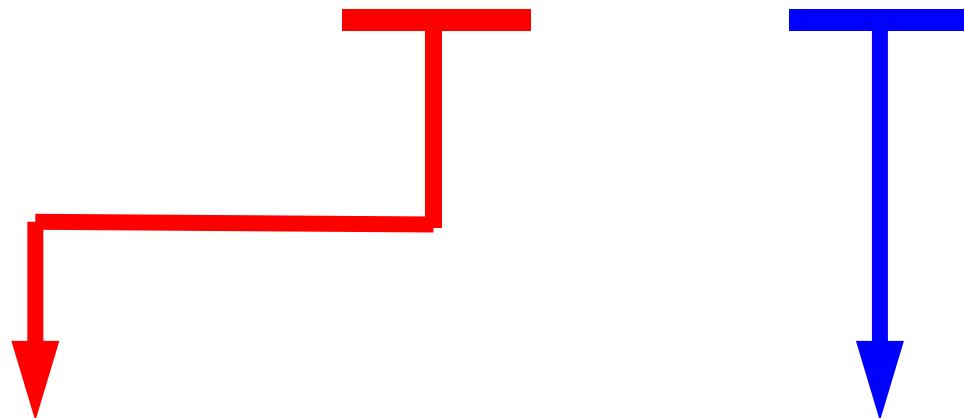


Hybrid: Hierarchical Structure





HB M x N

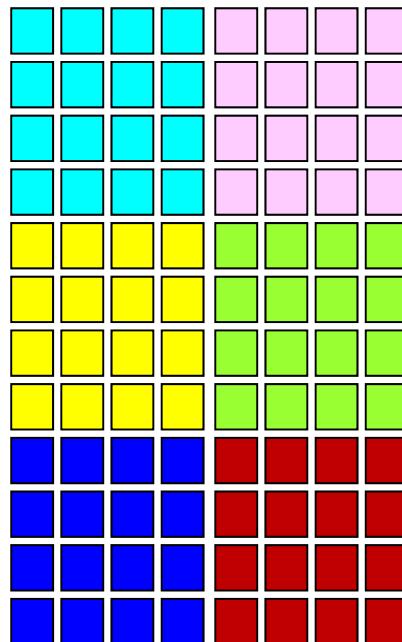


Number of OpenMP threads
per a single MPI process

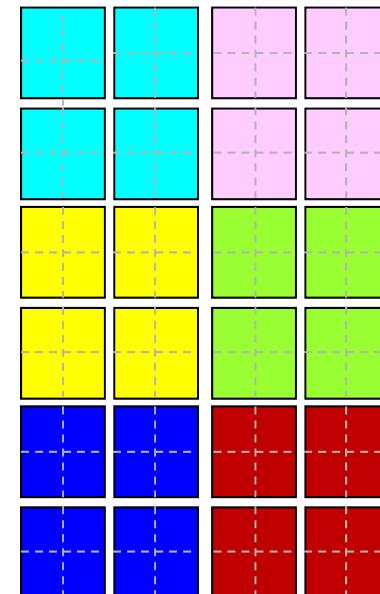
Number of MPI process
per a single node

Size of data for each MPI process varies according to HB MxN

example: 6 nodes, 96 cores

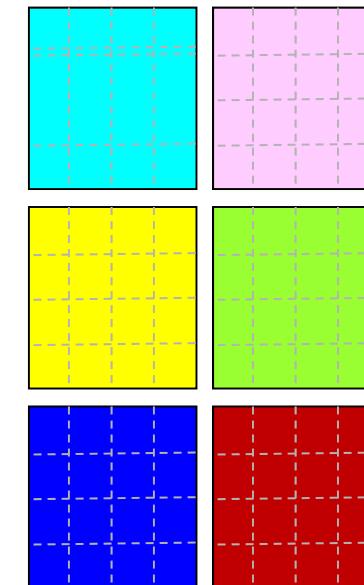


128	192	64
8	12	1
pcube		



HB 4x4

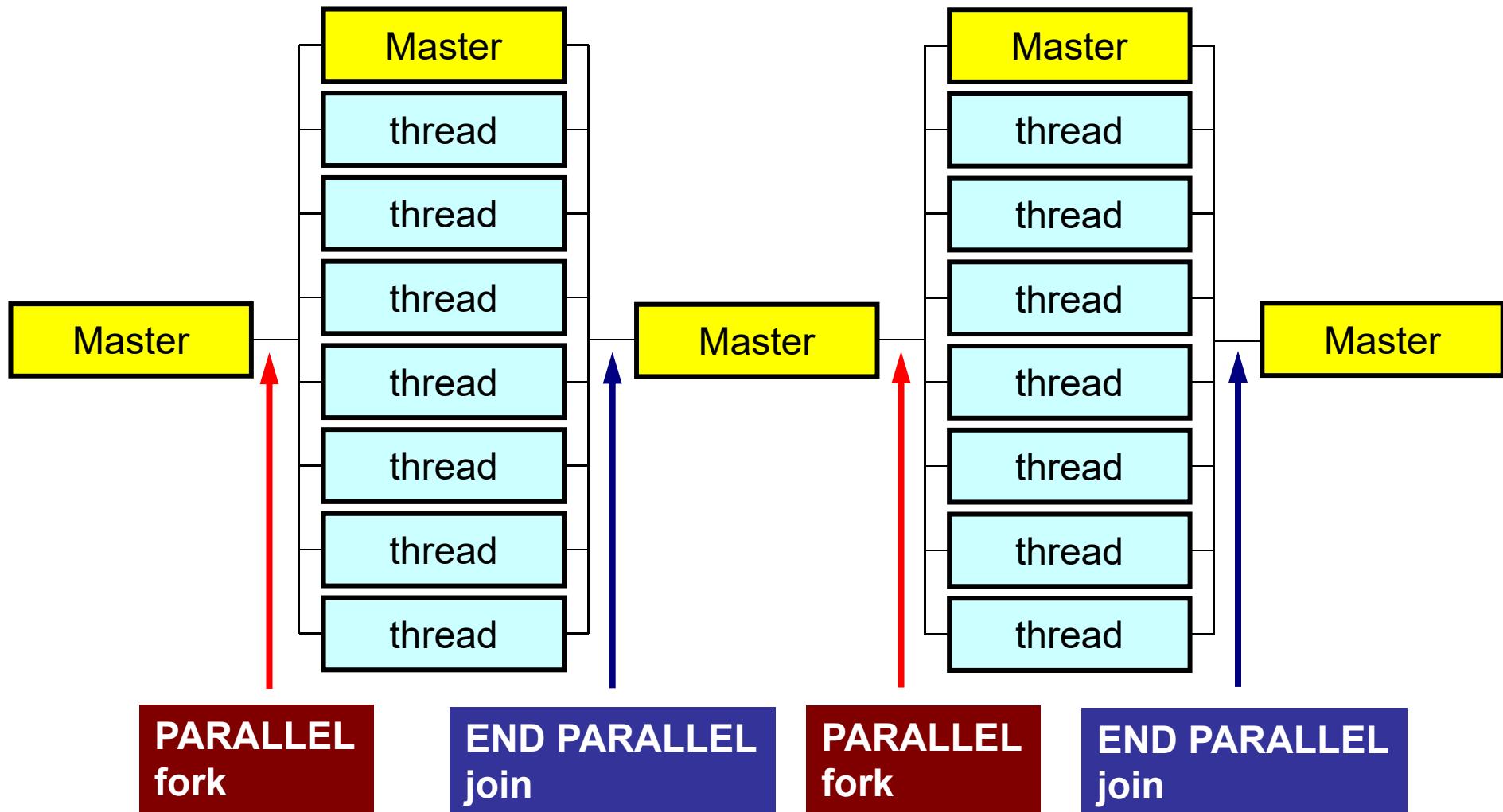
128	192	64
4	6	1
pcube		



HB 16x1

128	192	64
2	3	1
pcube		

Fork-Join Parallel Execution Model



Features of OpenMP

- Directives
 - Loops right after the directives are parallelized.
 - If the compiler does not support OpenMP, directives are considered as just comments.

OpenMP/Directives Array Operations

Simple Substitution

```
#pragma omp parallel for private (i)
for (i=0; i<N; i++) {
    X[i] = 0.0;
    W[0][i] = 0.0;
    W[1][i] = 0.0;
    W[2][i] = 0.0;
}
```

Dot Products

```
RHO = 0.0;
#pragma omp parallel for private (i)
reduction (+:RHO)
for (i=0; i<N; i++) {
    RHO += W[R][i] * W[Z][i];
}
```

DAXPY

```
#pragma omp parallel for private (i)
for (i=0; i<N; i++) {
    Y[i] = Y[i] + alpha*X[i];
}
```

OpenMP/Direceives Matrix/Vector Products

```
#pragma omp parallel for private (i, VAL, j)
for (i=0; i<N; i++) {
    VAL = D[i] * W[P][i];
    for (j=indexLU[i]; j<indexLU[i+1]; j++) {
        VAL += AMAT[j] * W[P][itemLU[j]-1];
    }
    W[Q][i] = VAL;
}
```

Features of OpenMP

- Directives
 - Loops right after the directives are parallelized.
 - If the compiler does not support OpenMP, directives are considered as just comments.
- Nothing happen without explicit directives
 - Different from “automatic parallelization/vectorization”
 - Something wrong may happen by un-proper way of usage
 - Data configuration, ordering etc. are done under users’ responsibility
- “Threads” are created according to the number of cores on the node
 - Thread: “Process” in MPI
 - Generally, “# threads = # cores”: Xeon Phi supports 4 threads per core (Hyper Multithreading)

Features of OpenMP (cont.)

- “for” loops with “#pragma omp parallel for”
- Global (Shared) Variables, Private Variables
 - Default: Global (Shared)
 - Dot Products: reduction

W[:, :], R, Z
global (shared)

```
RHO = 0.0;  
#pragma omp parallel for private (i) reduction (+:RHO)  
for(i=0; i<N; i++) {  
    RHO += W[R][i] * W[Z][i];  
}
```

FORTRAN & C

```
use omp_lib

```
 !$omp parallel do shared(n, x, y) private(i)
 do i= 1, n
 x(i)= x(i) + y(i)
 enddo
 !$omp end parallel do
```

```
#include <omp.h>
```
{
    #pragma omp parallel for default(none) shared(n, x, y) private(i)
    for (i=0; i<n; i++)
        x[i] += y[i];
}
```

OpenMP Directives (Fortran)

```
sentinel directive_name [clause[,] clause...]
```

- NO distinctions between upper and lower cases.
- sentinel
 - Fortran: !\$OMP, C\$OMP, *\$OMP
 - !\$OMP only for free format
 - Continuation Lines (Same rule as that of Fortran compiler is applied)
 - Example for !\$OMP PARALLEL DO SHARED(A,B,C)

```
!$OMP PARALLEL DO  
!$OMP+SHARED (A,B,C)
```

```
!$OMP PARALLEL DO &  
!$OMP SHARED (A,B,C)
```

OpenMP Directives (C)

```
#pragma omp directive_name [clause[,] clause...]
```

- “\” for continuation lines
- Only lower case (except names of variables)

```
#pragma omp parallel for shared (a,b,c)
```

PARALLEL DO

```
!$OMP PARALLEL DO[clause[,] clause] ... ]  
  (do_loop)  
 !$OMP END PARALLEL DO
```

```
#pragma parallel for [clause[,] clause] ... ]  
  (for_loop)
```

- Parallelize DO/for Loops
- Examples of “clause”
 - PRIVATE(list)
 - SHARED(list)
 - DEFAULT(PRIVATE|SHARED|NONE)
 - REDUCTION({operation|intrinsic}: list)

REDUCTION

```
REDUCTION ({operator|instinsic}: list)
```

```
reduction ({operator|instinsic}: list)
```

- Similar to “MPI_Reduce”
- Operator
 - +, *, -, .AND., .OR., .EQV., .NEQV.
- Intrinsic
 - MAX, MIN, IAND, IOR, IEQR

Example-1: A Simple Loop

```
#pragma omp parallel for
for(i=0; i<N; i++){
    B[i]= (A[i] + B[i]) * 0.50;
}
```

- Default status of loop variables (“i” in this case) is private. Therefore, explicit declaration is not needed.

Example-1: REDUCTION

```
#pragma omp parallel default(private) reduction(+:A,B)
for(i=0; i<N; i++){
    err= work(Alocal, Blocl);
    A= A + Alocal;
    B= B + Blocl;
}
```

Functions in OpenMP

functions	description
<code>int omp_get_num_threads (void)</code>	Thread #
<code>int omp_get_thread_num (void)</code>	Thread ID
<code>double omp_get_wtime (void)</code>	Timer
<code>void omp_set_num_threads (int num_threads)</code> <code>call omp_set_num_threads (num_threads)</code>	Specifying Thread #

OpenMP for Dot Products

```
VAL= 0.0;  
for(i=0; i<N; i++) {  
    VAL= VAL + W[R][i] * W[Z][i];  
}
```

OpenMP for Dot Products

```
VAL= 0.0;  
for(i=0; i<N; i++) {  
    VAL= VAL + W[R][i] * W[Z][i];  
}
```

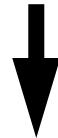


```
VAL= 0.0;  
#pragma omp parallel for private (i) reduction(+:VAL)  
for(i=0; i<N; i++) {  
    VAL= VAL + W[R][i] * W[Z][i];  
}
```

Directives are just inserted.

OpenMP for Dot Products

```
VAL= 0.0;
for(i=0; i<N; i++) {
    VAL= VAL + W[R][i] * W[Z][i];
}
```



```
VAL= 0.0;
#pragma omp parallel for private (i) reduction(+:VAL)
for(i=0; i<N; i++) {
    VAL= VAL + W[R][i] * W[Z][i];
}
```

Directives are just inserted.



```
VAL= 0.0;
#pragma omp parallel for private (i, ip)
reduction(+:VAL)
for(ip=0; ip<PEsmpTOT; ip++) {
    for (i=INDEX[ip]; i<INDEX[ip+1]; i++) {
        VAL= VAL + W[R][i] * W[Z][i];
    }
}
```

Multiple Loop
PEsmpTOT: Number of threads
Additional array **INDEX[:]** is needed.

Efficiency is not necessarily good, but users can specify thread for each component of data.

OpenMP for Dot Products

```

VAL= 0.0;
#pragma omp parallel for private (i, ip)
reduction(+:VAL)
    for(ip=0; ip<PEsmpTOT; ip++) {
        for (i=INDEX[ip]; i<INDEX[ip+1]; i++) {
            VAL= VAL + W[R][i] * W[Z][i];
        }
    }
}

```

Multiple Loop

PEsmpTOT: Number of threads

Additional array **INDEX[:]** is needed.

Efficiency is not necessarily good, but users can specify thread for each component of data.

e.g.: N=100, PEsmpTOT=4

```

INDEX[0]= 0
INDEX[1]= 25
INDEX[2]= 50
INDEX[3]= 75
INDEX[4]= 100

```

Matrix-Vector Multiply

```
for (i=0; i<N; i++) {  
    VAL = D[i] * W[P][i];  
    for (j=indexLU[i]; j<indexLU[i+1]; j++) {  
        VAL += AMAT[j] * W[P][itemLU[j]-1];  
    }  
    W[Q][i] = VAL;  
}
```

Matrix-Vector Multiply

```
#pragma omp parallel for private(ip, i, VAL, j)
for (ip=0; ip<PEsmpTOT; ip++) {
    for(i=SMPindexG[ip]; i<SMPindexG[ip+1]; i++) {
        VAL = D[i] * W[P][i];
        for (j=indexLU[i]; j<indexLU[i+1]; j++) {
            VAL += AMAT[j] * W[P][itemLU[j]-1];
        }
        W[Q][i] = VAL;
    }
}
```

Matrix-Vector Multiply: Other Approach

This is rather better for GPU and (very) many-core architectures: simpler structure of loops

```
#pragma omp parallel for private(i, VAL, j)
for (i=0; i<N; i++) {
    VAL = D[i] * W[P][i];
    for (j=indexLU[i]; j<indexLU[i+1]; j++) {
        VAL += AMAT[j] * W[P][itemLU[j]-1];
    }
    W[Q][i] = VAL;
}
```

omp parallel (do)

- Each “omp parallel-omp end parallel” pair starts & stops threads: fork-join
- If you have many loops, these operations on threads could be overhead
- omp parallel + omp do/omp for

```
!$omp parallel ...  
  
 !$omp do  
   do i= 1, N  
 ...  
 !$omp do  
   do i= 1, N  
 ...  
 !$omp end parallel
```

必須

```
#pragma omp parallel ...  
  
#pragma omp for {  
 ...  
#pragma omp for {
```

- OpenMP
- **Parallel Version of the Code by OpenMP**
- STREAM

Target for Parallelization

- FVM code
- Preconditioned CG solver: PCG
 - Diagonal Scaling, Point Jacobi
- NO sample code.
- Please develop the parallel code by yourself

Preconditioned Conjugate Gradient Method (PCG)

```

Compute  $r^{(0)} = b - [A]x^{(0)}$ 
for i= 1, 2, ...
  solve  $[M]z^{(i-1)} = r^{(i-1)}$ 
   $\rho_{i-1} = r^{(i-1)} \cdot z^{(i-1)}$ 
  if i=1
     $p^{(1)} = z^{(0)}$ 
  else
     $\beta_{i-1} = \rho_{i-1}/\rho_{i-2}$ 
     $p^{(i)} = p^{(i-1)} + \beta_{i-1} z^{(i-1)}$ 
  endif
   $q^{(i)} = [A]p^{(i)}$ 
   $\alpha_i = \rho_{i-1}/p^{(i)}q^{(i)}$ 
   $x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$ 
   $r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$ 
  check convergence |r|
end

```

Solving the following equation:

$$\{z\} = [M]^{-1}\{r\}$$

“Approximate Inverse Matrix”

$$[M]^{-1} \approx [A]^{-1}, \quad [M] \approx [A]$$

Ultimate Preconditioning:
Inverse Matrix

$$[M]^{-1} = [A]^{-1}, \quad [M] = [A]$$

Diagonal Scaling: Simple but weak

$$[M]^{-1} = [D]^{-1}, \quad [M] = [D]$$

Diagonal Scaling, Point-Jacobi

$$[M] = \begin{bmatrix} D_1 & 0 & \dots & 0 & 0 \\ 0 & D_2 & & 0 & 0 \\ \dots & & \dots & & \dots \\ 0 & 0 & & D_{N-1} & 0 \\ 0 & 0 & \dots & 0 & D_N \end{bmatrix}$$

- **solve $[M]z^{(i-1)} = r^{(i-1)}$** is very easy.
- Provides fast convergence for simple problems.

Files on Oakleaf-FX (1/2)

```
>$ cd  
  
>$ cp /home/z30088/omp/omp2-c.tar .  
>$ cp /home/z30088/omp/omp2-f.tar .
```

```
>$ tar xvf omp2-c.tar  
>$ tar xvf omp2-f.tar
```

```
>$ cd multicore
```

Confirm Directories:

omp2 stream

<\$O-omp2>, <\$O-stream>

Files on Oakleaf-FX (2/2)

```
>$ cd <$O-omp2>
>$ cd src

>$ make

>$ cd ../run
>$ pbsub go.sh
```

<\$O-omp2>/src/Makefile

NOT for parallel computing

```

CC      = fccpx          : Compiler
OPTFLAGS = -Kfast        : Optimization
TARGET   = ../run/sol    : Exec File

.SUFFIXES:
.SUFFIXES: .o .c

.c.o:
    $(CC) -c $(CFLAGS) $(OPTFLAGS) $< -o $@

OBJS = input.o pointer_init.o \
       boundary_cell.o cell_metrics.o \
       rcm.o ...

HEADERS = \
    struct.h struct_ext.h \
    pcg.h ...

all: $(TARGET)
$(TARGET): $(OBJS)
    $(CC) $(CFLAGS) $(OPTFLAGS) -o $@ $(OBJS) $(LDFLAGS)

$(OBJS): $(HEADERS)

clean:
    rm -f *.o $(TARGET) *.log *~ *.lst

```

Running Job

- Batch Jobs
 - Only batch jobs are allowed.
 - Interactive executions of jobs are not allowed.
- How to run
 - writing job script
 - submitting job
 - checking job status
 - checking results
- Utilization of computational resources
 - 1-node (16 cores) is occupied by each job.
 - Your node is not shared by other jobs.

Job Script

- <\$O-omp>/run/go.sh
- Scheduling + Shell Script

```
#!/bin/sh
#PJM -L "node=1"
#PJM -L "elapse=00:10:00"
#PJM -L "rscgrp=lecture7"
#PJM -g "gt17"
#PJM -j
#PJM -o "test.lst"

./sol
```

Number of Nodes
Computation Time
Name of “QUEUE”
Group Name (Wallet)
Standard Output
Execs

Available QUEUE's

- Following 2 queues are available.
- 1 Tofu (12 nodes) can be used
 - **lecture**
 - 12 nodes (192 cores), 15 min., **valid until March 30 08:30**
 - Shared by all “educational” users
 - **lecture7**
 - 12 nodes (192 cores), 15 min., active during class time (09:00-17:00)
 - **More jobs (compared to lecture) can be processed up on availability.**
 - **Just during until Feb.23 17:00**
- **Please use “lecture” after Feb.24 !!**

Submitting & Checking Jobs

- Submitting Jobs `pjsub SCRIPT NAME`
- Checking status of jobs `pjstat`
- Deleting/aborting `pjdel JOB ID`
- Checking status of queues `pjstat --rsc`
- Detailed info. of queues `pjstat --rsc -x`
- Number of running jobs `pjstat --rsc -b`
- Limitation of submission `pjstat --limit`

```
[z30088@oakleaf-fx-6 S2-ref]$ pjstat
```

```
Oakleaf-FX scheduled stop time: 2012/09/28(Fri) 09:00:00 (Remain: 31days 20:01:46)
```

JOB_ID	JOB_NAME	STATUS	PROJECT	RSCGROUP	START_DATE	ELAPSE	TOKEN	NODE:COORD
334730	go.sh	RUNNING	gt61	lecture	08/27 12:58:08	00:00:05		0.0 1

solve_PCG (1/3)

```

for(i=0; i<N; i++) {
    W[DD][i] = 1.e0/D[i];
}

...
for(L=0; L<(*ITR); L++) {

/******
 * {z} = [Minv] {r} *
*****/
    for(i=0; i<N; i++) {
        W[Z][i] = W[R][i]*W[DD][i];
    }

/******
 * RHO = {r} {z} *
*****/
    RHO = 0.0;
    for(i=0; i<N; i++) {
        RHO += W[R][i] * W[Z][i];
    }
}

```

Compute $r^{(0)} = b - [A]x^{(0)}$

for $i = 1, 2, \dots$

solve $[M]z^{(i-1)} = r^{(i-1)}$

$\rho_{i-1} = r^{(i-1)} \cdot z^{(i-1)}$

if $i = 1$

$p^{(1)} = z^{(0)}$

else

$\beta_{i-1} = \rho_{i-1}/\rho_{i-2}$

$p^{(i)} = z^{(i-1)} + \beta_{i-1} p^{(i)}$

endif

$q^{(i)} = [A]p^{(i)}$

$\alpha_i = \rho_{i-1}/p^{(i)} q^{(i)}$

$x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$

$r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$

check convergence $|r|$

end

solve_PCG (2/3)

```

/*****
 * {p} = {z} if ITER=0
 * BETA = RHO / RH01 otherwise
 *****/
if(L == 0) {
    for(i=0; i<N; i++) {
        W[P][i] = W[Z][i];
    } else {
        BETA = RHO / RH01;
        for(i=0; i<N; i++) {
            W[P][i] = W[Z][i] + BETA * W[P][i];
        }
    }
}

/*****
 * {q} = [A]{p}
 *****/
for(i=0; i<N; i++) {
    VAL = D[i] * W[P][i];
    for(j=indexLU[i]; j<indexLU[i+1]; j++) {
        VAL += AMAT[j] * W[P][itemLU[j]-1];
    }
    W[Q][i] = VAL;
}

```

Compute $r^{(0)} = b - [A]x^{(0)}$

for $i = 1, 2, \dots$

solve $[M]z^{(i-1)} = r^{(i-1)}$

$\rho_{i-1} = r^{(i-1)} \cdot z^{(i-1)}$

if $i = 1$

$p^{(1)} = z^{(0)}$

else

$\beta_{i-1} = \rho_{i-1} / \rho_{i-2}$

$p^{(i)} = z^{(i-1)} + \beta_{i-1} p^{(i)}$

endif

$q^{(i)} = [A]p^{(i)}$

$\alpha_i = \rho_{i-1} / p^{(i)} q^{(i)}$

$x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$

$r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$

check convergence $|r|$

end

```

/*****************/
 * ALPHA = RHO / {p} {q} *
/*****************/
C1 = 0.0;
for(i=0; i<N; i++) {
    C1 += W[P][i] * W[Q][i];
}
ALPHA = RHO / C1;

/*****************/
 * {x} = {x} + ALPHA * {p} *
 * {r} = {r} - ALPHA * {q} *
/*****************/
for(i=0; i<N; i++) {
    X[i] += ALPHA * W[P][i];
    W[R][i] -= ALPHA * W[Q][i];
}

DNRM2 = 0.0;
for(i=0; i<N; i++) {
    DNRM2 += W[R][i]*W[R][i];
}

ERR = sqrt(DNRM2/BNRM2);
if((L+1)%100 ==1) {
    fprintf(stderr, "%5d%16.6e\n", L+1, ERR);
}
if(ERR < EPS) {
    *IER = 0; goto N900;
} else {
    RH01 = RHO;
}
*IER = 1;

```

$r = b - Ax$
 $DNRM2 = |r|^2$
 $BNRM2 = |b|^2$
 $ERR = |r| / |b|$

solve_PCG (3/3)

```

Compute r^(0) = b - [A]x^(0)
for i= 1, 2, ...
    solve [M]z^(i-1) = r^(i-1)
    rho_i-1 = r^(i-1) z^(i-1)
    if i=1
        p^(1) = z^(0)
    else
        beta_i-1 = rho_i-1 / rho_i-2
        p^(i) = z^(i-1) + beta_i-1 p^(i)
    endif
    q^(i) = [A]p^(i)
    alpha_i = rho_i-1 / p^(i) q^(i)
    x^(i) = x^(i-1) + alpha_i p^(i)
    r^(i) = r^(i-1) - alpha_i q^(i)
    check convergence |r|
end

```

Parallelization by OpenMP

- Focusing on “solver_PCG.c” (solve_PCG)
- Just insert OpenMP directives

```
>$ cd <$O-omp2>
>$ cd ex
(modify files)

>$ make
>$ cd ../run
>$ pbsub g.sh
```

```
#!/bin/sh
#PJM -L "node=1"
#PJM -L "elapse=00:10:00"
#PJM -L "rscgrp=lecture7"
#PJM -g "gt17"
#PJM -j
#PJM -o "100_08_002.lst"
```

```
export OMP_NUM_THREADS=8          1-16
./sol0
```

<\$O-omp2>/ex/Makefile

parallel computing by OpenMP

```

CC      = fccpx          : Compiler
OPTFLAGS = -Kfast,openmp : Optimization + OpenMP
TARGET  = ../run/sol0   : Exec File

.SUFFIXES:
.SUFFIXES: .o .c

.c.o:
    $(CC) -c $(CFLAGS) $(OPTFLAGS) $< -o $@

OBJS = input.o pointer_init.o \
        boundary_cell.o cell_metrics.o \
        rcm.o ...

HEADERS = \
    struct.h struct_ext.h \
    pcg.h ...

all: $(TARGET)
$(TARGET): $(OBJS)
    $(CC) $(CFLAGS) $(OPTFLAGS) -o $@ $(OBJS) $(LDFLAGS)

$(OBJS): $(HEADERS)

clean:
    rm -f *.o $(TARGET) *.log *~ *.lst

```

solve_PCG (1/5)

parallel computing by OpenMP

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <errno.h>
#include <math.h>
#include <omp.h>

#include "solver_PCG.h"

extern int
solve_PCG (int N, int NPLU, int *indexLU, int *itemLU,
           double *D, double *B, double *X, double *AMAT,
           double EPS, int *ITR, int *IER, int *N2)
{
    double **W;
    double VAL, BNRM2, WVAL, SW, RHO, BETA, RH01, C1, DNRM2, ALPHA, ERR;
    double Stime, Etime;
    int i, j, ic, ip, L, ip1, N3;
    int R = 0;
    int Z = 1;
    int Q = 1;
    int P = 2;
    int DD = 3;
```

solve_PCG (2/5)

```

#pragma omp parallel for private (i)
    for(i=0; i<N; i++) {
        X[i] = 0.0;
        W[1][i] = 0.0;
        W[2][i] = 0.0;
        W[3][i] = 0.0;
    }

#pragma omp parallel for private (i,VAL,j)
    for(i=0; i<N; i++) {
        VAL = D[i] * X[i];
        for(j=indexLU[i]; j<indexLU[i+1]; j++) {
            VAL += AMAT[j] * X[itemLU[j]-1];
        }
        W[R][i] = B[i] - VAL;
    }

BNRM2 = 0.0;
#pragma omp parallel for private (i) reduction
(+:BNRM2)
    for(i=0; i<N; i++) {
        BNRM2 += B[i]*B[i];
    }

#pragma omp parallel for private (i)
    for(i=0; i<N; i++) {
        W[DD][i]= 1.e0/D[i];
    }

```

Compute $r^{(0)} = b - [A]x^{(0)}$
for $i = 1, 2, \dots$
 solve $[M]z^{(i-1)} = r^{(i-1)}$
 $\rho_{i-1} = r^{(i-1)} \cdot z^{(i-1)}$
if $i=1$
 $p^{(1)} = z^{(0)}$
else
 $\beta_{i-1} = \rho_{i-1}/\rho_{i-2}$
 $p^{(i)} = z^{(i-1)} + \beta_{i-1} p^{(i)}$
endif
 $q^{(i)} = [A]p^{(i)}$
 $\alpha_i = \rho_{i-1}/p^{(i)}q^{(i)}$
 $x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$
 $r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$
 check convergence $|r|$
end

```

*ITR = N;
Stime = omp_get_wtime();
for (L=0; L<(*ITR); L++) {
#pragma omp parallel for private (i)
    for(i=0; i<N; i++) {
        W[Z][i] = W[R][i]*W[DD][i];
    }

RHO = 0.0;
#pragma omp parallel for private (i)
reduction(+:RHO)
    for(i=0; i<N; i++) {
        RHO += W[R][i] * W[Z][i];
    }

if(L == 0) {
#pragma omp parallel for private (i)
    for(i=0; i<N; i++) {
        W[P][i] = W[Z][i];
    }
} else {
    BETA = RHO / RH01;
#pragma omp parallel for private (i)
    for(i=0; i<N; i++) {
        W[P][i] = W[Z][i] + BETA * W[P][i];
    }
}

```

solve_PCG (3/5)

Compute $r^{(0)} = b - [A]x^{(0)}$
for $i = 1, 2, \dots$
solve $[M]z^{(i-1)} = r^{(i-1)}$
 $\rho_{i-1} = r^{(i-1)} \cdot z^{(i-1)}$
if $i = 1$
 $p^{(1)} = z^{(0)}$
else
 $\beta_{i-1} = \rho_{i-1} / \rho_{i-2}$
 $p^{(i)} = z^{(i-1)} + \beta_{i-1} p^{(i)}$
endif
 $q^{(i)} = [A]p^{(i)}$
 $\alpha_i = \rho_{i-1} / p^{(i)} q^{(i)}$
 $x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$
 $r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$
 check convergence $|r|$
end

solve PCG (4/5)

```

#pragma omp parallel for private (i,VAL,j)
for(i=0; i<N; i++) {
    VAL = D[i] * W[P][i];
    for(j=indexLU[i]; j<indexLU[i+1]; j++) {
        VAL += AMAT[j] * W[P][itemLU[j]-1];
    }
    W[Q][i] = VAL;
}

C1 = 0.0;
#pragma omp parallel for private (i) reduction(+:C1)
for(i=0; i<N; i++) {
    C1 += W[P][i] * W[Q][i];
}
ALPHA = RHO / C1;

#pragma omp parallel for private (i)
for(i=0; i<N; i++) {
    X[i]    += ALPHA * W[P][i];
    W[R][i] -= ALPHA * W[Q][i];
}

DNRM2 = 0.0;
#pragma omp parallel for private (i) reduction(+:DNRM2)
for(i=0; i<N; i++) {
    DNRM2 += W[R][i]*W[R][i];
}
ERR = sqrt(DNRM2/BNRM2);
...

```

Compute $r^{(0)} = b - [A]x^{(0)}$

for $i = 1, 2, \dots$

solve $[M]z^{(i-1)} = r^{(i-1)}$

$\rho_{i-1} = r^{(i-1)} \cdot z^{(i-1)}$

if $i=1$

$p^{(1)} = z^{(0)}$

else

$\beta_{i-1} = \rho_{i-1}/\rho_{i-2}$

$p^{(i)} = z^{(i-1)} + \beta_{i-1} p^{(i)}$

endif

$q^{(i)} = [A]p^{(i)}$

$\alpha_i = \rho_{i-1}/p^{(i)}q^{(i)}$

$x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$

$r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$

check convergence $|r|$

end

solve_PCG (5/5)

```
Stime = omp_get_wtime(); ←  
  
for (L=0; L<(*ITR); L++) {  
    ...  
    if (ERR < EPS) {  
        *IER = 0;  
        goto N900;  
    } else {  
        RH01 = RHO;  
    }  
}  
*IER = 1;  
N900:  
    Etime = omp_get_wtime(); ←  
  
    fprintf(stderr, "%5d%16.6e\n", L+1, ERR);  
    fprintf(stderr, "%16.6e sec. (solver)\n", Etime - Stime);  
  
    *ITR = L;  
    free(W);  
    return 0;  
}
```

$$\text{Elapsed Time} = \text{Etime} - \text{Stime}$$

Etime-Stime

$N_X=N_Y=N_Z=100$, 10^6 DOF

```
#!/bin/sh
#PJM -L "node=1"
#PJM -L "elapse=00:10:00"
#PJM -L "rscgrp=lecture7"
#PJM -g "gt17"
#PJM -j
#PJM -o "100_08_002.lst"

export OMP_NUM_THREADS=M      M=1,4 8,16
./sol
```

M	sec.	Speed-Up
1	33.7	1.00
4	8.73	3.86
8	4.95	6.80
16	3.30	10.20

Exercises

- Develop your own program by inserting OpenMP directives !
- Effect of problem size (NX, NY, NZ)
- Effect of Thread # (OMP_NUM_THREADS: 1-16)

- OpenMP
- Login to FX10
- Parallel Version of the Code by OpenMP
- **STREAM**

Why less than 16x ?

- Memory Contention
- Performance of memory per each thread decreases if number of threads on each node increases
- Sparse Matrix Solver: Memory-Bound
 - Effect of this decreasing is more significant
- Problem size is not so larger

Sparse/Dense Matrices

```

for (i=0; i<N; i++) {
    Y[i] = Diag[i] * X[i];
    for (k=Index[i]; k<Index[i+1]; k++) {
        Y[i] += AMat[k]*X[Item[k]];
    }
}

```

```

for (j=0; j<N; i++) {
    Y[j]= 0.0;
    for (i=0; i<N; i++) {
        Y[j]+= A[j][i]*X[i];
    }
}

```

- “X” in RHS
 - Dense: continuous on memory, easy to utilize cache
 - Sparse: continuity is not assured, difficult to utilize cache
 - more “memory-bound”

GeoFEM Benchmark

ICCG in FEM for Solid Mechanics

	SR11K/J2	SR16K/M1	T2K	FX10	京
Core #/Node	16	32	16	16	8
Peak Performance (GFLOPS)	147.2	980.5	147.2	236.5	128.0
STREAM Triad (GB/s)	101.0	264.2	20.0	64.7	43.3
B/F	0.686	0.269	0.136	0.274	0.338
GeoFEM (GFLOPS)	19.0	72.7	4.69	16.0	11.0
% to Peak	12.9	7.41	3.18	6.77	8.59
LLC/core (MB)	18.0	4.00	2.00	0.75	0.75

Sparse Linear Solver: Memory-Bound

STREAM benchmark

<http://www.cs.virginia.edu/stream/>

- Benchmarks for Memory Bandwidth
 - Copy: $c(i) = a(i)$
 - Scale: $c(i) = s \cdot b(i)$
 - Add: $c(i) = a(i) + b(i)$
 - Triad: $c(i) = a(i) + s \cdot b(i)$

Double precision appears to have 16 digits of accuracy
Assuming 8 bytes per DOUBLE PRECISION word

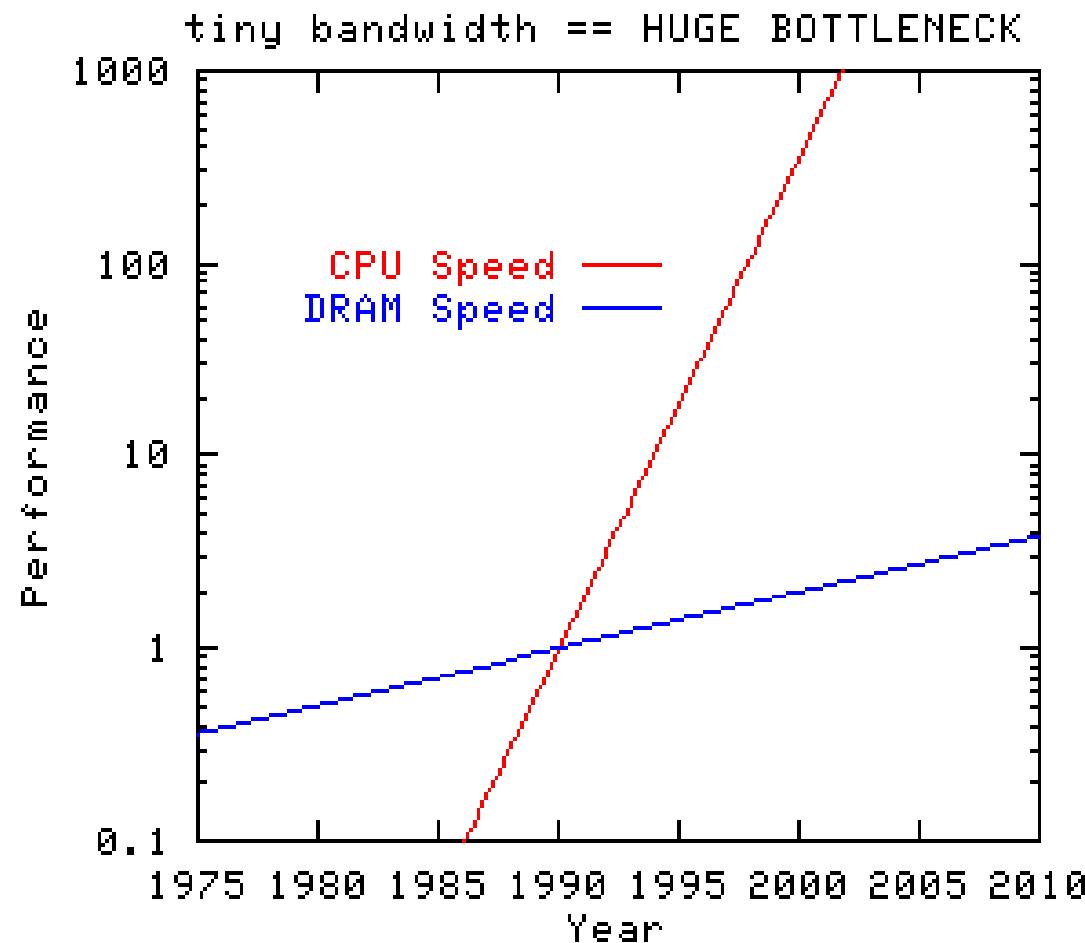
Number of processors = 16
 Array size = 2000000
 Offset = 0
 The total memory requirement is 732.4 MB
 (45.8MB/task)

You are running each test 10 times

--
 The *best* time for each test is used
 EXCLUDING the first and last iterations

Function	Rate (MB/s)	Avg time	Min time	Max time
Copy:	18334.1898	0.0280	0.0279	0.0280
Scale:	18035.1690	0.0284	0.0284	0.0285
Add:	18649.4455	0.0412	0.0412	0.0413
Triad:	19603.8455	0.0394	0.0392	0.0398

Gap between performance of CPU and Memory



OpenMP version of STREAM

```
>$ cd <$O-stream>
>$ pjsub go.sh
```

- <http://www.cs.virginia.edu/stream/>
- C, Fortran, MPI, OpenMP etc.

go.sh

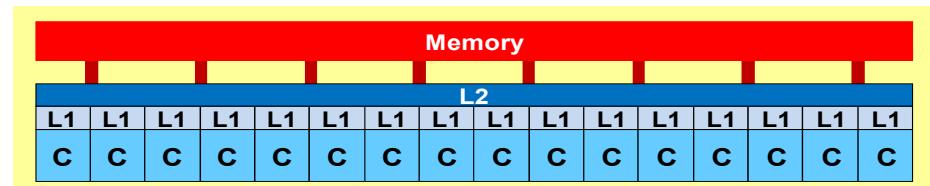
```
#!/bin/sh
#PJM -L "rscgrp=lecture7"
#PJM -L "node=1"
#PJM -L "elapse=10:00"
#PJM -j

export PATH=...
export LD_LIBRARY_PATH=...
export PARALLEL=16
export OMP_NUM_THREADS=16          Number of threads (1-16)

./stream.out > 16-01.lst 2>&1      Name of output file
```

Results of Triad

`<$O-stream>/stream/*.lst`
 Peak is 85.3 GB/sec., 75%



Thread #	MB/sec.	Speed-up
1	8606.14	1.00
2	16918.81	1.97
4	34170.72	3.97
8	59505.92	6.91
16	64714.32	7.52

Exercises

- Running the code
- Try various number of threads (1-16)
- MPI-version and Single PE version are available
 - Fortran, C
 - Web-site of STREAM