Report S2
Fortran

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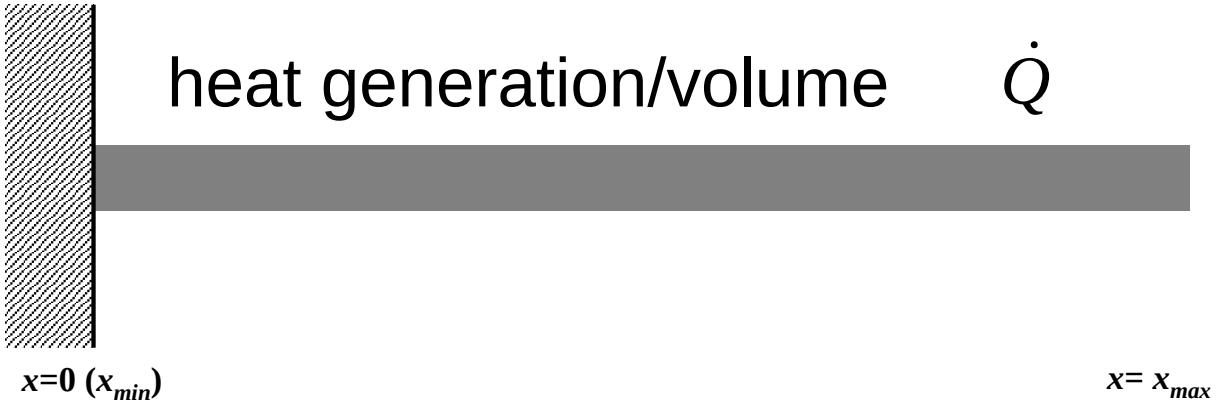
- Overview
- Distributed Local Data
- Program
- Results

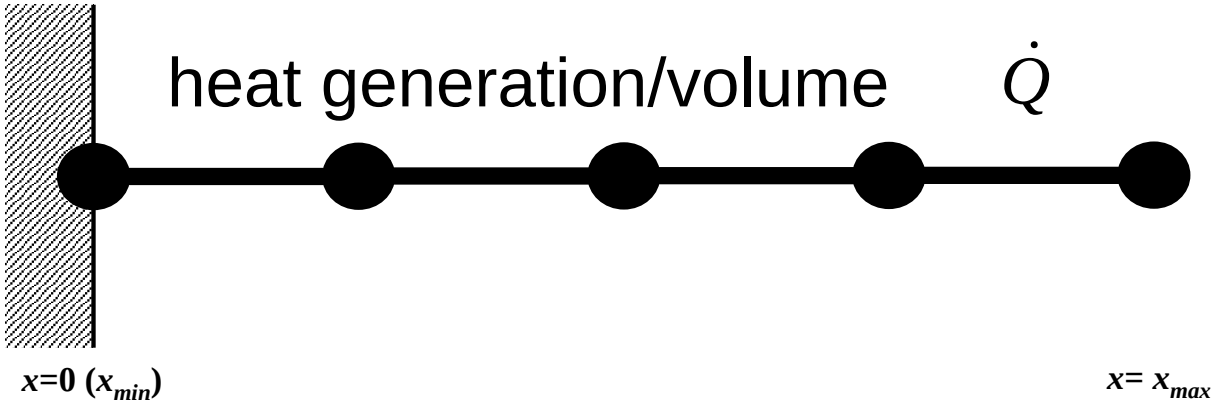
1D Steady State Heat Conduction



$$\frac{\partial}{\partial x}\left(\lambda \frac{\partial T}{\partial x}\right)+\dot{Q}=0$$

- Uniform: Sectional Area: A, Thermal Conductivity: λ
- Heat Generation Rate/Volume/Time [$QL^{-3}T^{-1}$] $\dot{Q}$
- Boundary Conditions
  - $x=0$: $T=0$ (Fixed Temperature)
  - $x=x_{max}$: $\dfrac{\partial T}{\partial x}=0$ (Insulated)

1D Steady State Heat Conduction



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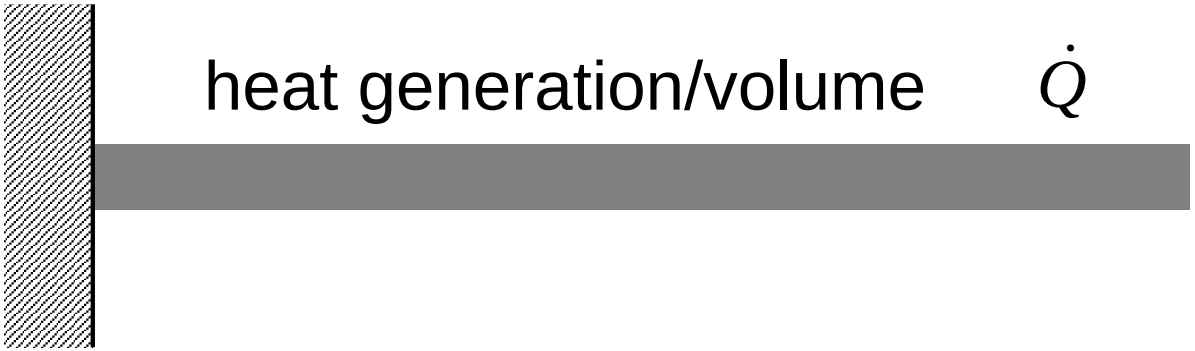
Analytical Solution



$$\frac{\partial}{\partial x}\left(\lambda\frac{\partial T}{\partial x}\right)+\dot{Q}=0$$

$$T=0 @ x=0$$

$$\frac{\partial T}{\partial x}=0 @ x=x_{\max}$$

$$\lambda T''=-\dot{Q}$$

$$\lambda T'=-\dot{Q}x+C_1 \Rightarrow C_1=\dot{Q}x_{\max}, \quad T'=0 @ x=x_{\max}$$

$$\lambda T=-\frac{1}{2}\dot{Q}x^2+C_1x+C_2 \Rightarrow C_2=0, \quad T=0 @ x=0$$

$$\therefore \quad T=-\frac{1}{2\lambda}\dot{Q}x^2+\frac{\dot{Q}x_{\max}}{\lambda}x$$

Report S2 (1/2)

- Parallelize 1D code (1d.f) using MPI
- Read entire element number, and decompose into sub-domains in your program

- Validate the results
  - Answer of Original Code = Answer of Parallel Code
  - Explain why number of iterations does not change, as number of MPI processes changes.
- Measure parallel performance

Report S2 (2/2)

- Deadline: January 28th (Thu), 2021, 17:00@ITC-LMS
- Problem
  - Apply “Generalized Communication Table”
  - Read entire elem. #, decompose into sub-domains in your program
  - Evaluate parallel performance
    - You need huge number of elements, to get excellent performance.
    - Fix number of iterations (e.g. 100), if computations cannot be completed.
- Report
  - Cover Page: Name, ID, and Problem ID (S2) must be written.
  - Less than eight pages including figures and tables (A4).
    - Strategy, Structure of the Program, Remarks
  - Source list of the program (if you have bugs)
  - Output list (as small as possible)

Copy and Compile

Fortran

```
>$ cd /work/gt39/t39XXX/pFEM
>$ cp /work/gt00/z30088/pFEM/F/s2r-f.tar .
>$ tar xvf s2r-f.tar
```

C

```
>$ cd /work/gt39/t39XXX/pFEM
>$ cp /work/gt00/z30088/pFEM/C/s2r-c.tar .
>$ tar xvf s2r-c.tar
```

Confirm/Compile

```
>$ cd mpi/S2-ref
>$ mpiifort -align array64byte -O3 -axCORE-AVX512 1d.f -o 1da
>$ mpiifort -O3 1d.f -o 1db

>$ mpiicc -align -O3 -axCORE-AVX512 1d.c -o 1da
>$ mpiicc -O3 1d.c -o 1db
```

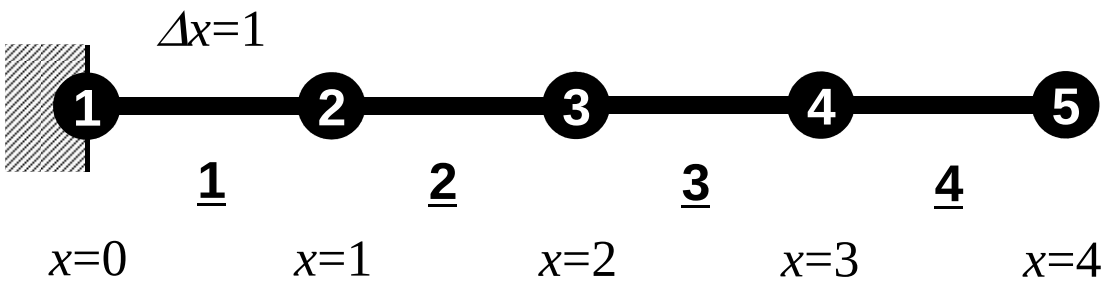
<$O-S2r> = <$O-TOP>/mpi/S2-ref

Control File: input.dat

Control Data **`input.dat`**

```
1000000
1.0  1.0  1.0  1.0
100
1.e-8
```

`NE` (Number of Elements)
Δx (Length of Each Elem.: `L`) , `Q`, `A`, λ
Number of MAX. Iterations for CG Solver
Convergence Criteria for CG Solver



Element ID
Node ID (Global)

g16.sh: 8-nodes, 256-cores, 16x2x8

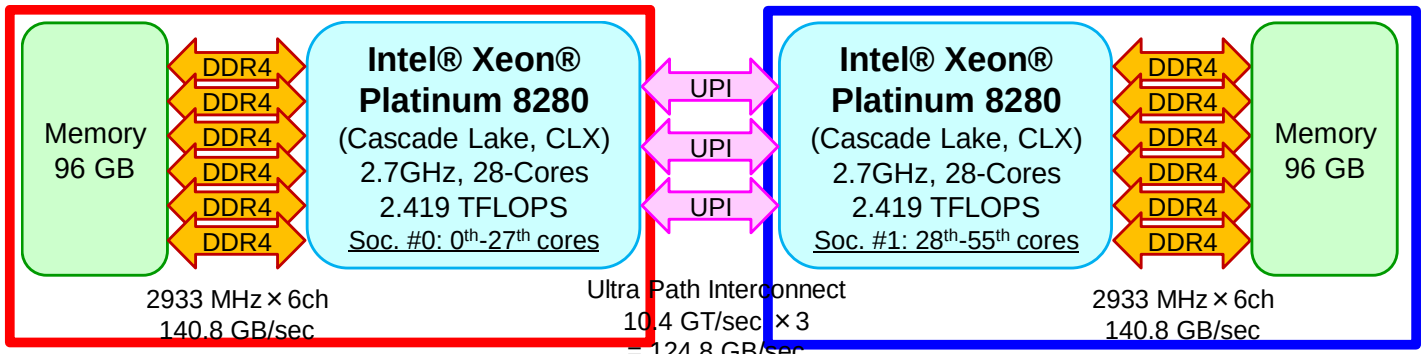
```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture2
#PJM -L node=8
#PJM --mpi proc=256
#PJM -L elapse=00:15:00
#PJM -g gt62
#PJM -j
#PJM -e err
#PJM -o test.lst

mpiexec.hydra -n ${PJM_MPI_PROC} ./1da
mpiexec.hydra -n ${PJM_MPI_PROC} ./1db
export I_MPI_PIN_PROCESSOR_LIST=0-15,28-43
mpiexec.hydra -n ${PJM_MPI_PROC} ./1da
mpiexec.hydra -n ${PJM_MPI_PROC} ./1db
```

256/8= 32 cores/node

1da: -O3 + AVX512
1db: -O3 Only

32-cores are randomly selected from 56-cores on the node

32-cores on each socket are assigned. A little bit more stable



Socket #0: 0th-27th Cores

Socket #1: 28th-55th Cores

g24.sh: 8-nodes, 384-cores, 24x2x8

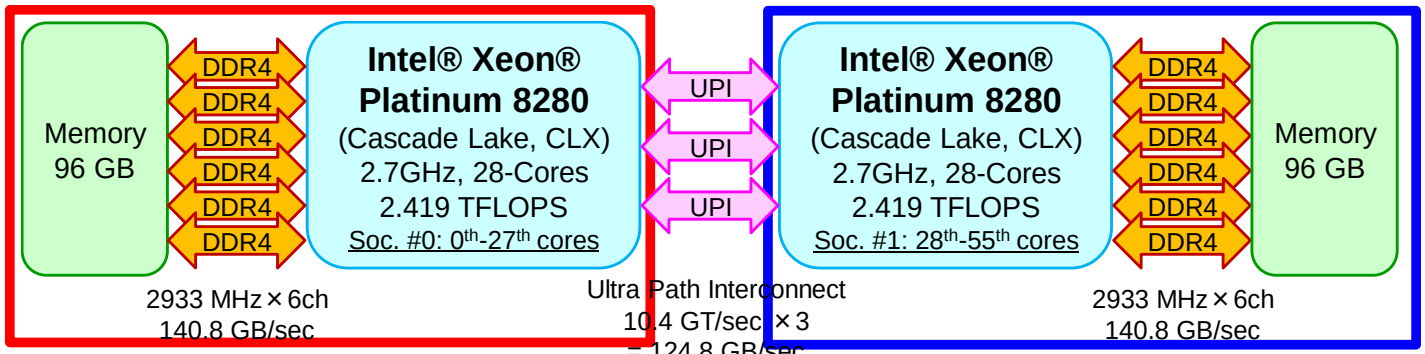
```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture2
#PJM -L node=8
#PJM --mpi proc=384
#PJM -L elapse=00:15:00
#PJM -g gt62
#PJM -j
#PJM -e err
#PJM -o test.lst

mpiexec.hydra -n ${PJM_MPI_PROC} ./1da
mpiexec.hydra -n ${PJM_MPI_PROC} ./1db
export I_MPI_PIN_PROCESSOR_LIST=0-23,28-51
mpiexec.hydra -n ${PJM_MPI_PROC} ./1da
mpiexec.hydra -n ${PJM_MPI_PROC} ./1db
```

384/8= 48 cores/node

1da: -O3 + AVX512
1db: -O3 Only

48-cores are randomly selected from 56-cores on the node

24-cores on each socket are assigned. A little bit more stable



Socket #0: 0th-27th Cores

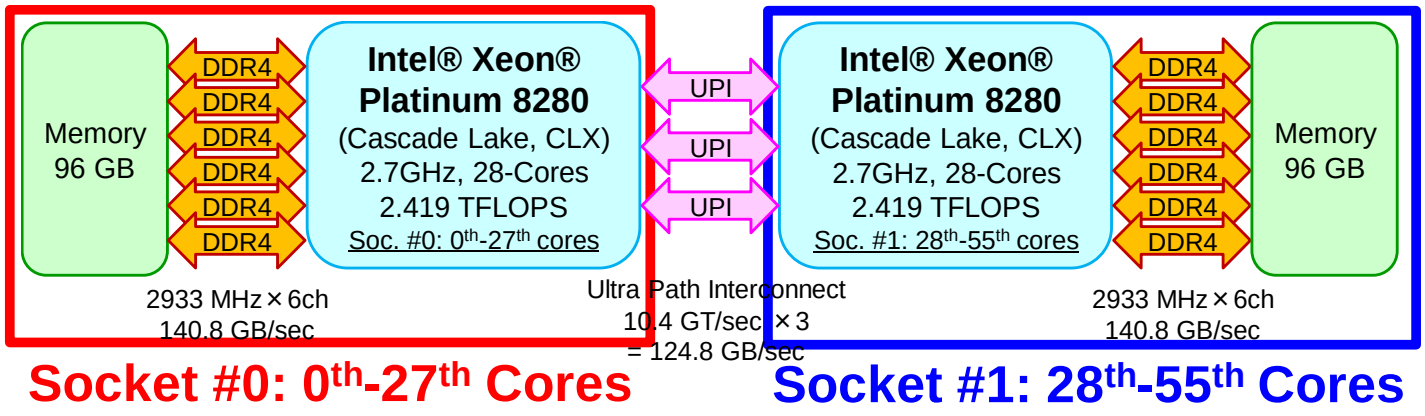
Socket #1: 28th-55th Cores

g28.sh: 8-nodes, 448-cores, 28x2x8

```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture2
#PJM -L node=8
#PJM --mpi proc=448
#PJM -L elapse=00:15:00
#PJM -g gt62
#PJM -j
#PJM -e err
#PJM -o test.lst

mpiexec.hydra -n ${PJM_MPI_PROC} ./1da
mpiexec.hydra -n ${PJM_MPI_PROC} ./1db
export I_MPI_PIN_PROCESSOR_LIST=0-55
mpiexec.hydra -n ${PJM_MPI_PROC} ./1da
mpiexec.hydra -n ${PJM_MPI_PROC} ./1db
```

448/8= 56 cores/node

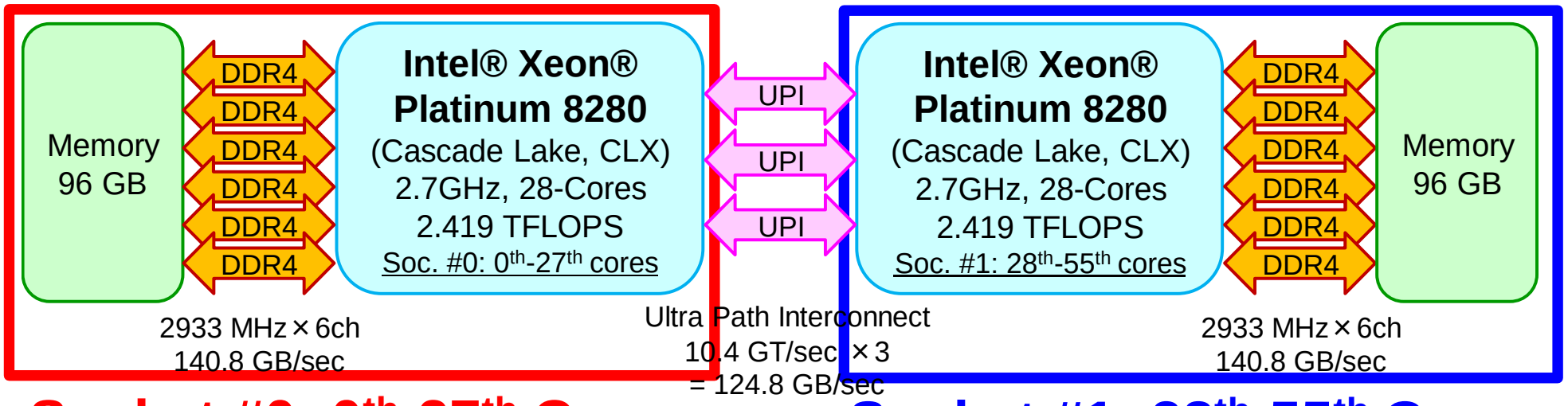
1da: -O3 + AVX512
1db: -O3 Only



Socket #0: 0th-27th Cores

Socket #1: 28th-55th Cores

NUMA Architecture



- Each Node of Oakbridge-CX (OBCX)
  - 2 Sockets (CPU's) of Intel CLX
  - Each socket has 28 cores
- Each core of a socket can access to the memory on the other socket : NUMA (Non-Uniform Memory Access)
  - `numactl -l` : local memory to be used
  - Sometimes (not always), more stable with this option

Procedures for Parallel FEM

- Reading control file, entire element number etc.
- Creating “distributed local data” in the program
- Assembling local and global matrices for linear solvers
- Solving linear equations by CG

- Not so different from those of original code

- Overview
- **Distributed Local Data**
- Program
- Results

Finite Element Procedures

- Initialization
  - Control Data
  - Node, Connectivity of Elements (N: Node#, NE: Elem#)
  - Initialization of Arrays (Global/Element Matrices)
  - Element-Global Matrix Mapping (Index, Item)
- Generation of Matrix
  - Element-by-Element Operations (do icel= 1, NE)
    - Element matrices
    - Accumulation to global matrix
  - Boundary Conditions
- Linear Solver
  - Conjugate Gradient Method

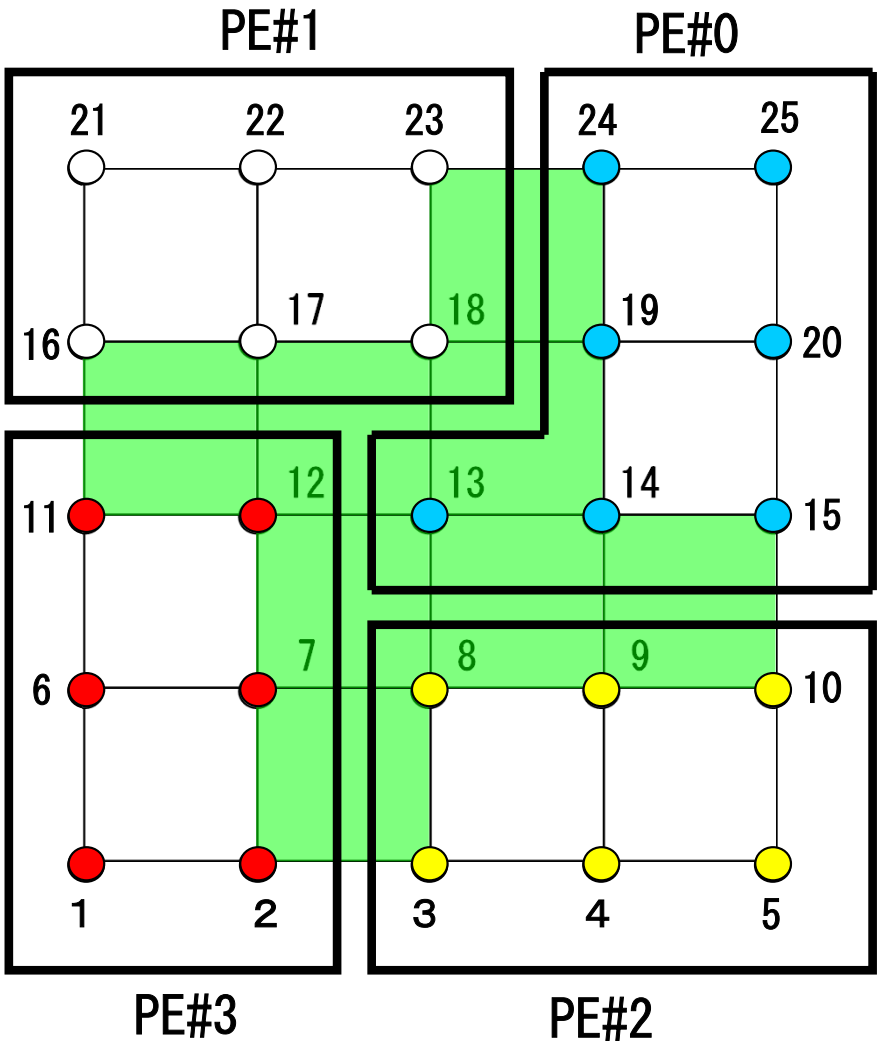
Distributed Local Data Structure for Parallel FEM

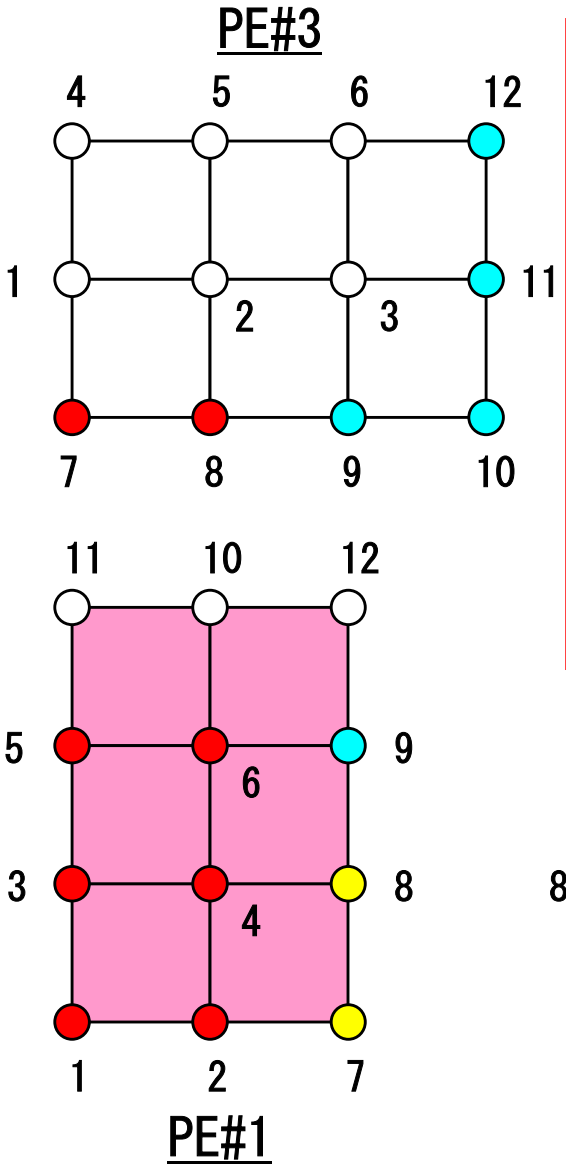
- Node-based partitioning
- Local data includes:
  - Nodes originally assigned to the domain/PE/partition
  - Elements which include above nodes
  - Nodes which are included above elements, and originally NOT-assigned to the domain/PE/partition
- 3 categories for nodes
  - Internal nodes Nodes originally assigned to the domain/PE/partition
  - External nodes Nodes originally NOT-assigned to the domain/PE/partition
  - Boundary nodes External nodes of other domains/PE's/partitions
- Communication tables

- Global info. is not needed except relationship between domains
  - Property of FEM: local element-by-element operations

Node-based Partitioning

internal nodes - elements - external nodes

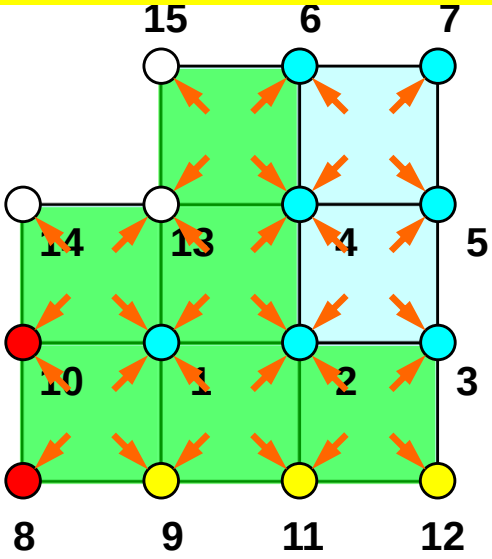


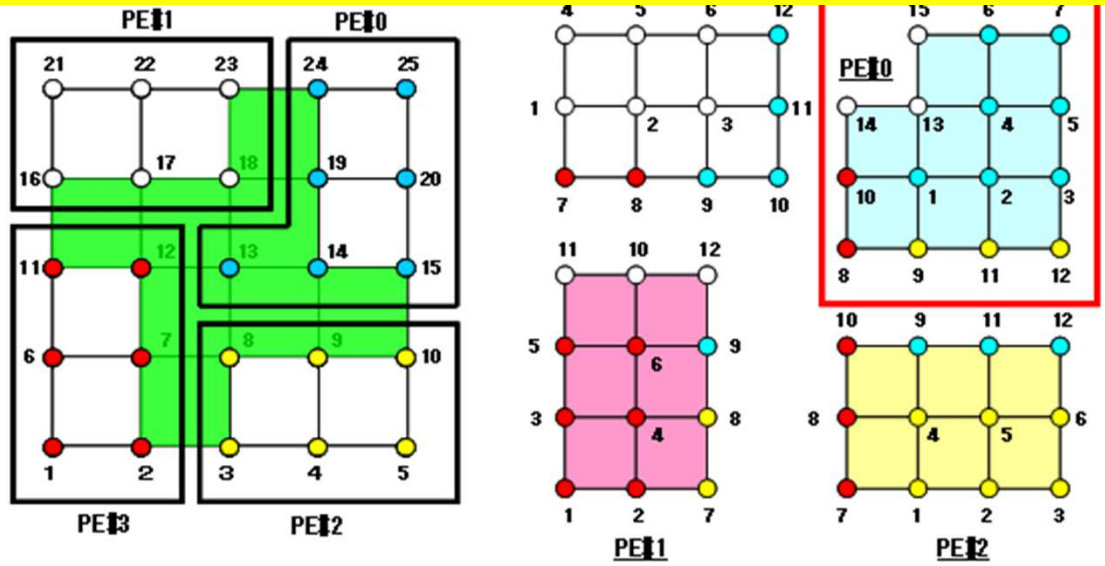


Node-based Partitioning

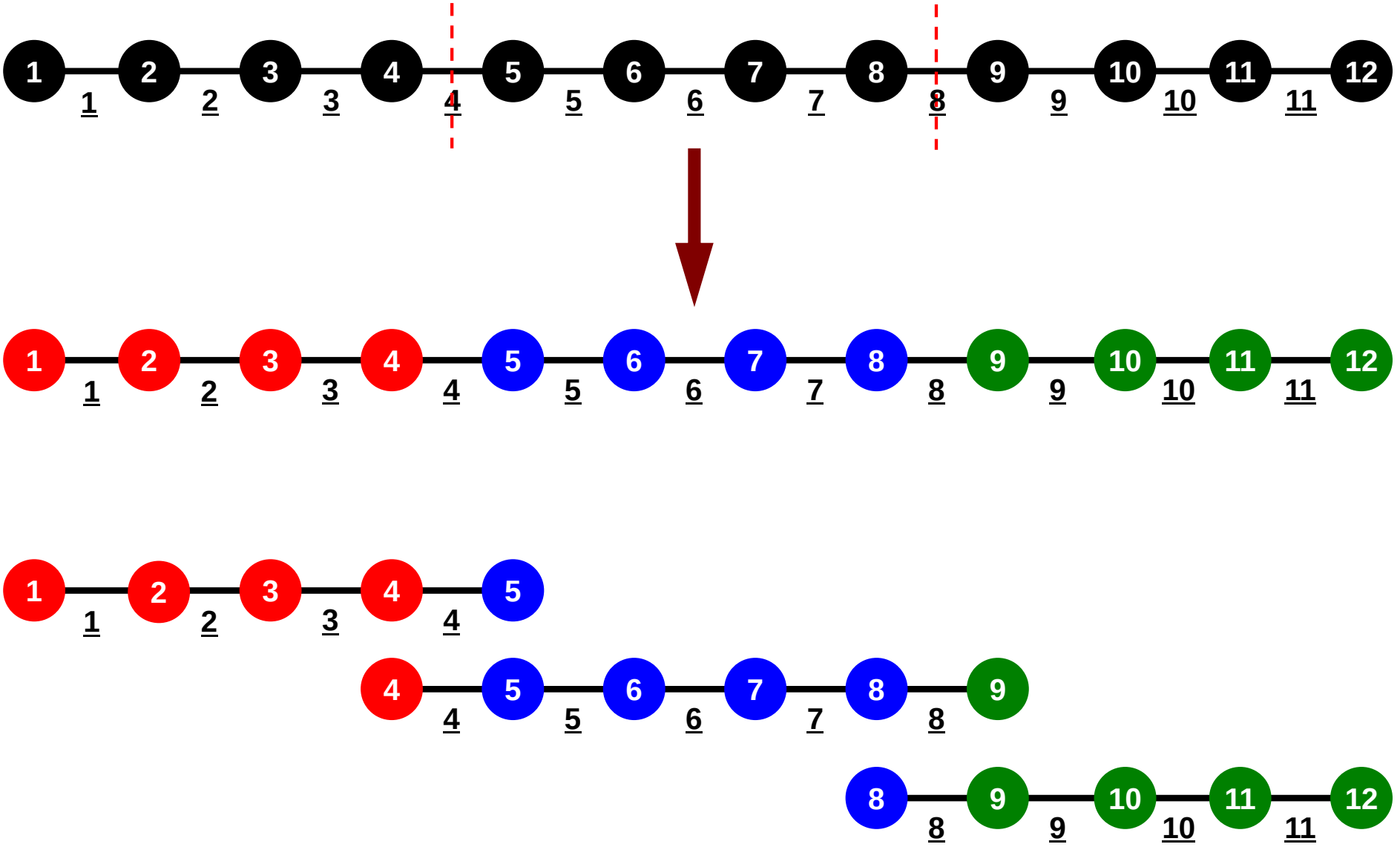
internal nodes - elements - external nodes

- Partitioned nodes themselves (Internal Nodes)　内点
- Elements which include Internal Nodes　内点を含む要素
- External Nodes included in the Elements　外点
  in overlapped region among partitions.
- Info of External Nodes are required for completely local element–based operations on each processor.

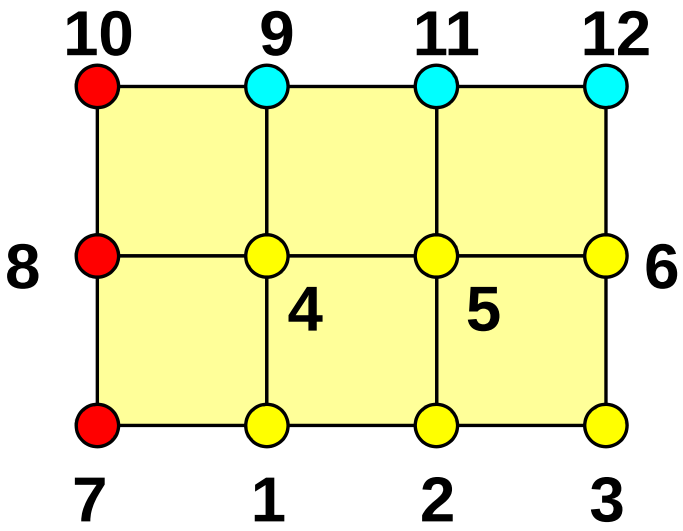




1D FEM: 12 nodes/11 elem's/3 domains



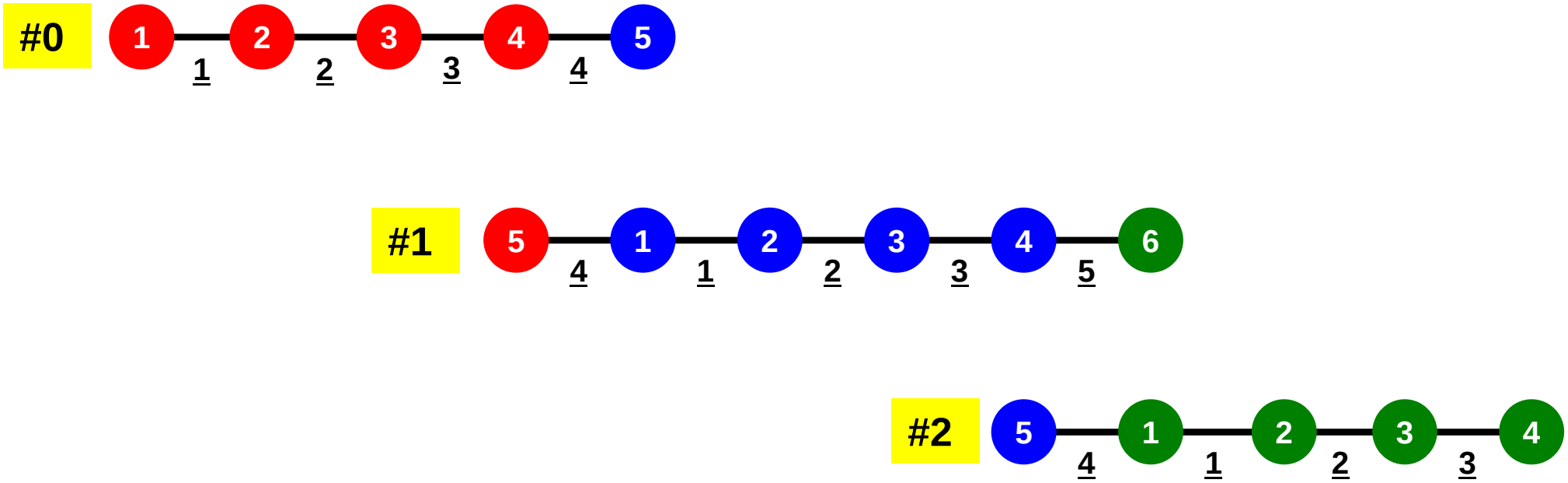
Description of Distributed Local Data



- Internal/External Points
  - Numbering: Starting from internal pts, then external pts after that
- Neighbors
  - Shares overlapped meshes
  - Number and ID of neighbors
- External Points
  - From where, how many, and which external points are received/imported ?
- Boundary Points
  - To where, how many and which boundary points are sent/exported ?

1D FEM: 12 nodes/11 elem’s/3 domains

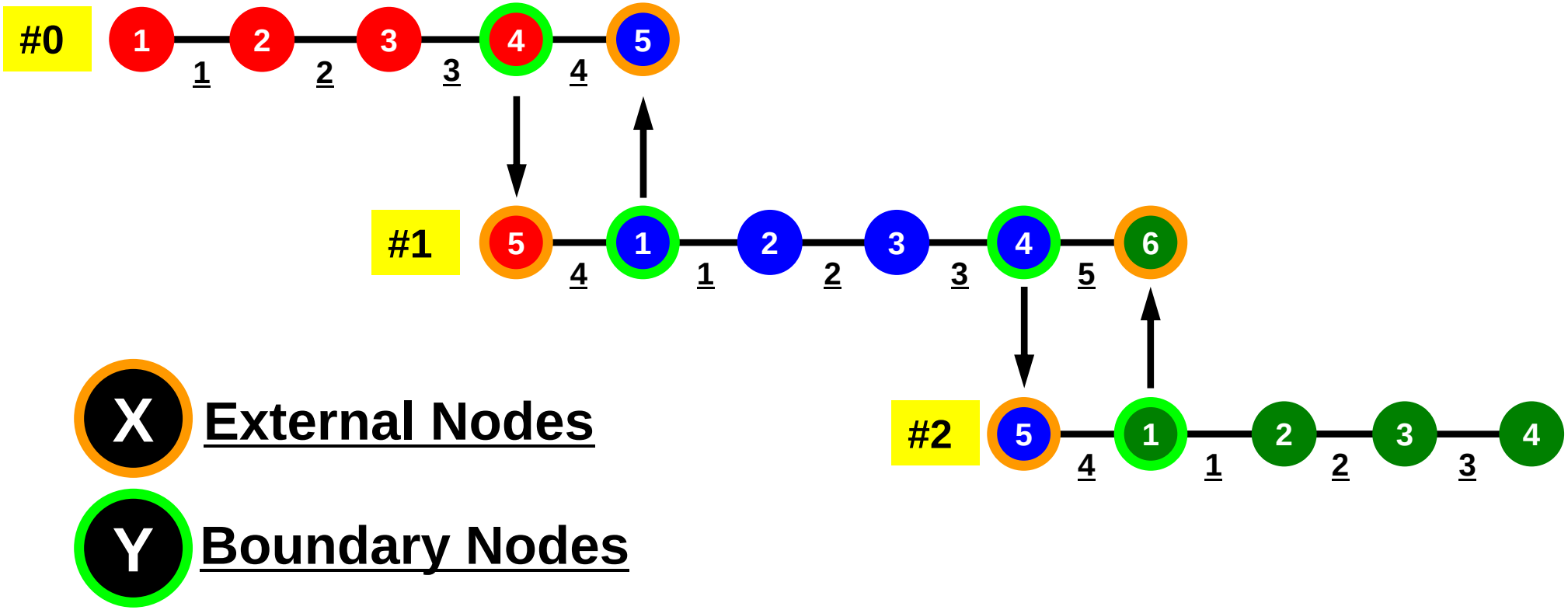
Local ID: Starting from 1 for node and elem at each domain



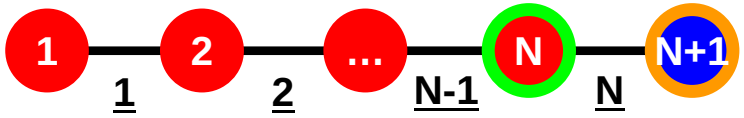
1D FEM: 12 nodes/11 elem's/3 domains

Internal/External/Boundary Nodes

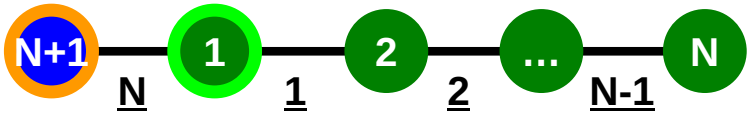
Boundary Nodes: Part of Internal Nodes, and External Nodes of Other Domains



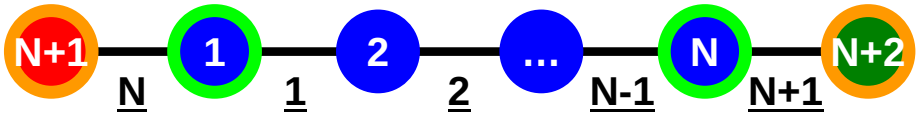
1D FEM: Numbering of Local ID



#0:
N+1 nodes
N elements



#PETot-1:
N+1 nodes
N elements

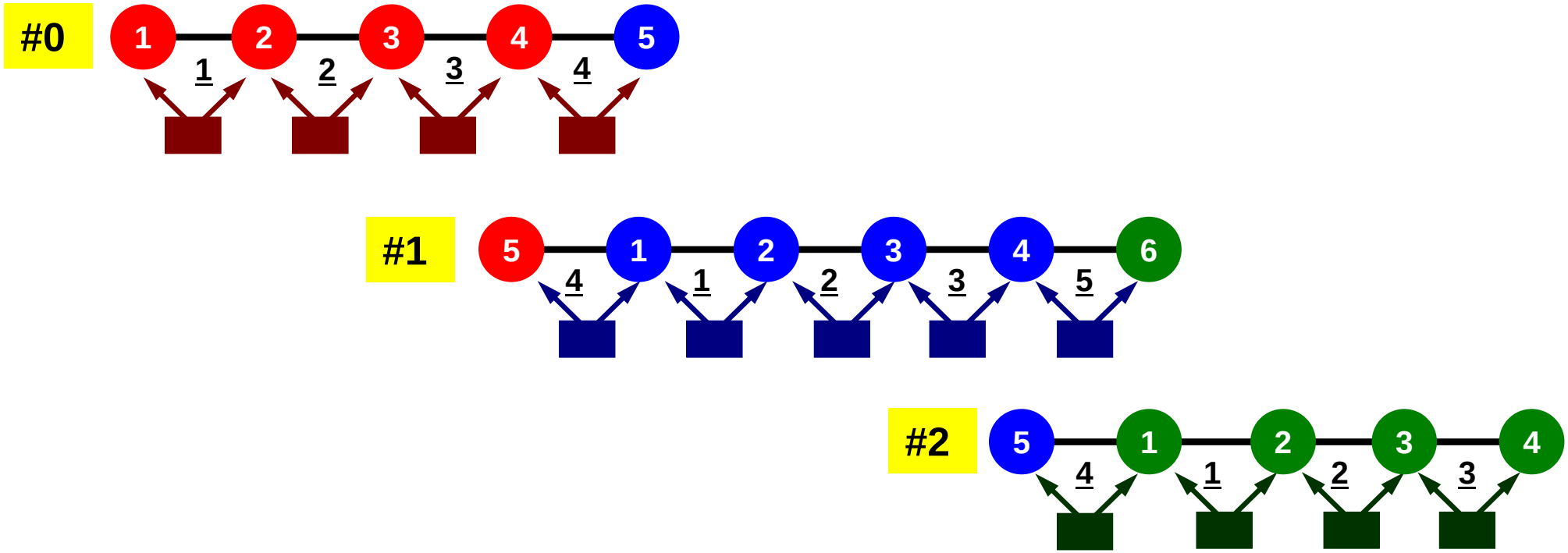


Others (General):
N+2 nodes
N+1 elements

1D FEM: 12 nodes/11 elem's/3 domains

Integration on each element, element matrix -> global matrix

Operations can be done by info. of internal/external nodes and elements which include these nodes



Preconditioned Conjugate Gradient Method (CG)

Compute $\mathbf{r^{(0)} = b - [A]x^{(0)}}$
for i= 1, 2, …
  solve $\mathbf{[M]z^{(i-1)} = r^{(i-1)}}$
  $\rho_{i-1} = \mathbf{r^{(i-1)}\ z^{(i-1)}}$
  if i=1
    $p^{(1)} = z^{(0)}$
  else
    $\beta_{i-1} = \rho_{i-1}/\rho_{i-2}$
    $p^{(i)} = z^{(i-1)} + \beta_{i-1}\ p^{(i-1)}$
  endif
  $\mathbf{q^{(i)} = [A]p^{(i)}}$
  $\alpha_i = \rho_{i-1}/\mathbf{p^{(i)}q^{(i)}}$
  $x^{(i)} = x^{(i-1)} + \alpha_i p^{(i)}$
  $r^{(i)} = r^{(i-1)} - \alpha_i q^{(i)}$
  check convergence $\mathbf{|r|}$
end

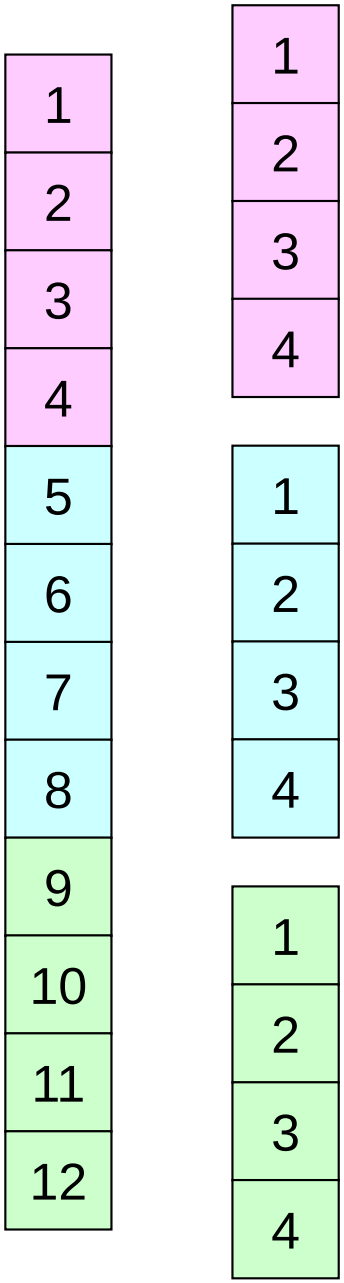
Preconditioning:
Diagonal Scaling
(or Point Jacobi)

Preconditioning, DAXPY

Local Operations by Only Internal Points: Parallel Processing is possible

```
!C
!C-- {z}= [Minv]{r}

      do i= 1, N
        W(i,Z)= W(i,DD) * W(i,R)
      enddo
```

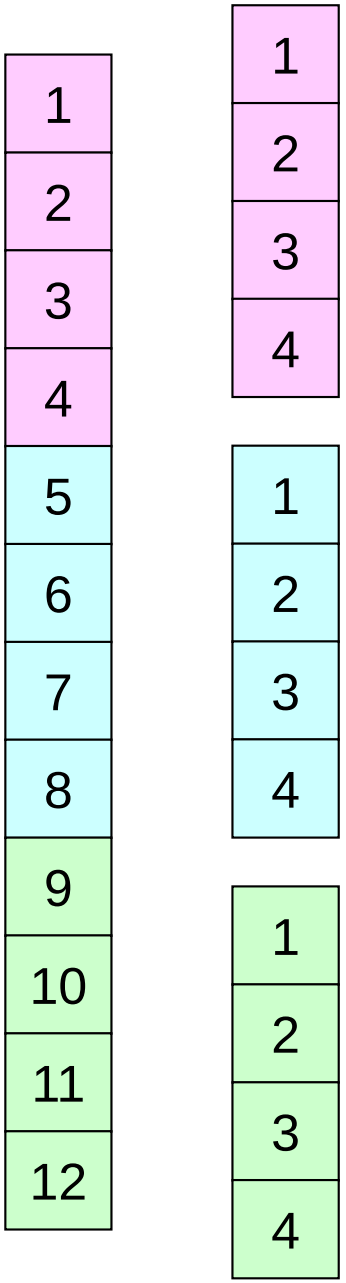
```
!C
!C-- {x}= {x} + ALPHA*{p}
!C   {r}= {r} - ALPHA*{q}

      do i= 1, N
        PHI(i)= PHI(i) + ALPHA * W(i,P)
        W(i,R)= W(i,R) - ALPHA * W(i,Q)
      enddo
```



Dot Products

Global Summation needed: Communication ?

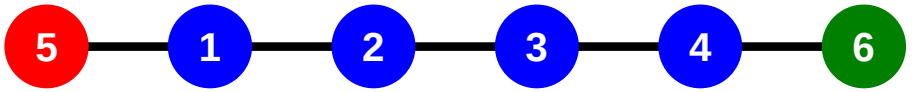
```
!C
!C-- ALPHA= RHO / {p}{q}

      C1= 0.d0
      do i= 1, N
        C1= C1 + W(i,P)*W(i,Q)
      enddo
      ALPHA= RHO / C1
```



Matrix-Vector Products

Values at External Points: P-to-P Communication

```
!C
!C-- {q}= [A]{p}

      do i= 1, N
        W(i,Q) = DIAG(i)*W(i,P)
        do j= INDEX(i-1)+1, INDEX(i)
          W(i,Q) = W(i,Q) + AMAT(j)*W(ITEM(j),P)
        enddo
      enddo
```



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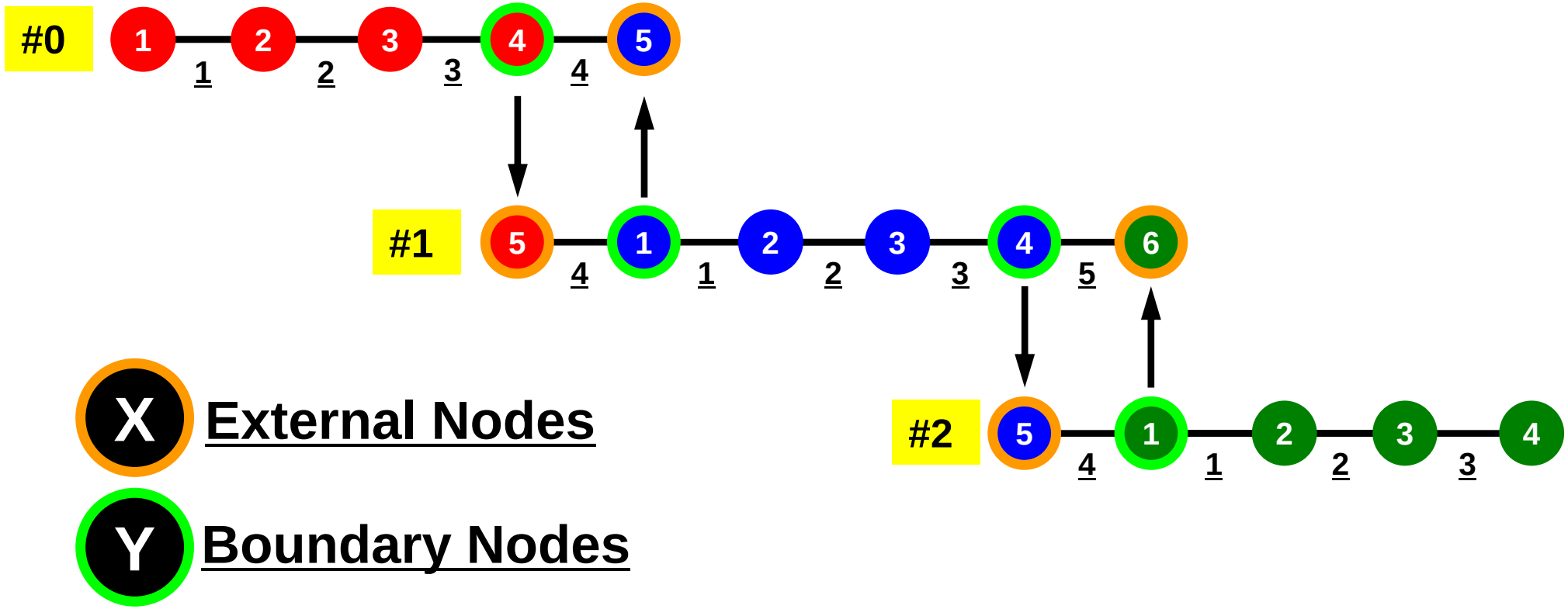
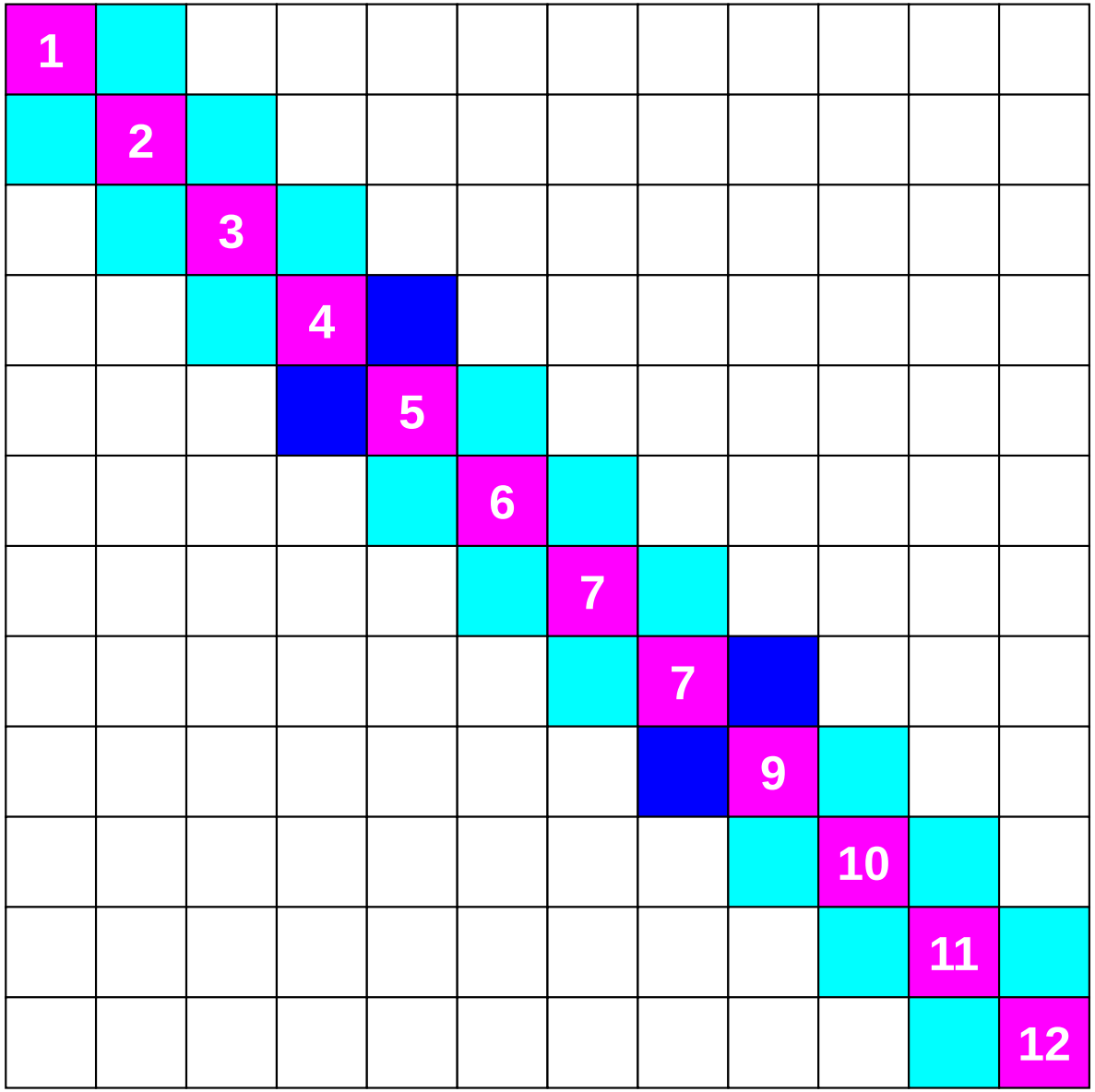
Internal/External/Boundary Nodes

Boundary Nodes: Part of Internal Nodes, and External Nodes of Other Domains



Mat-Vec Products: Local Op. Possible



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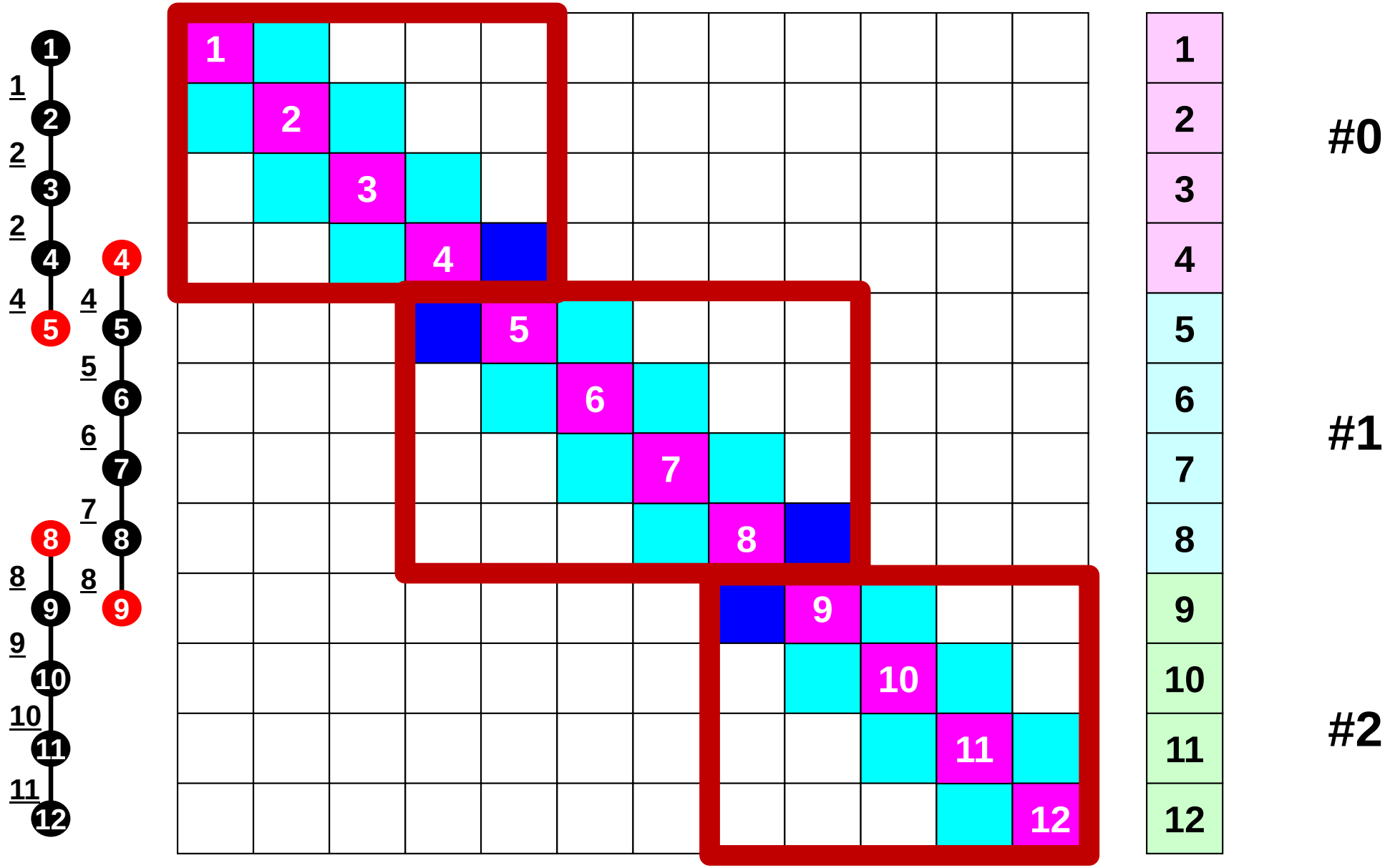
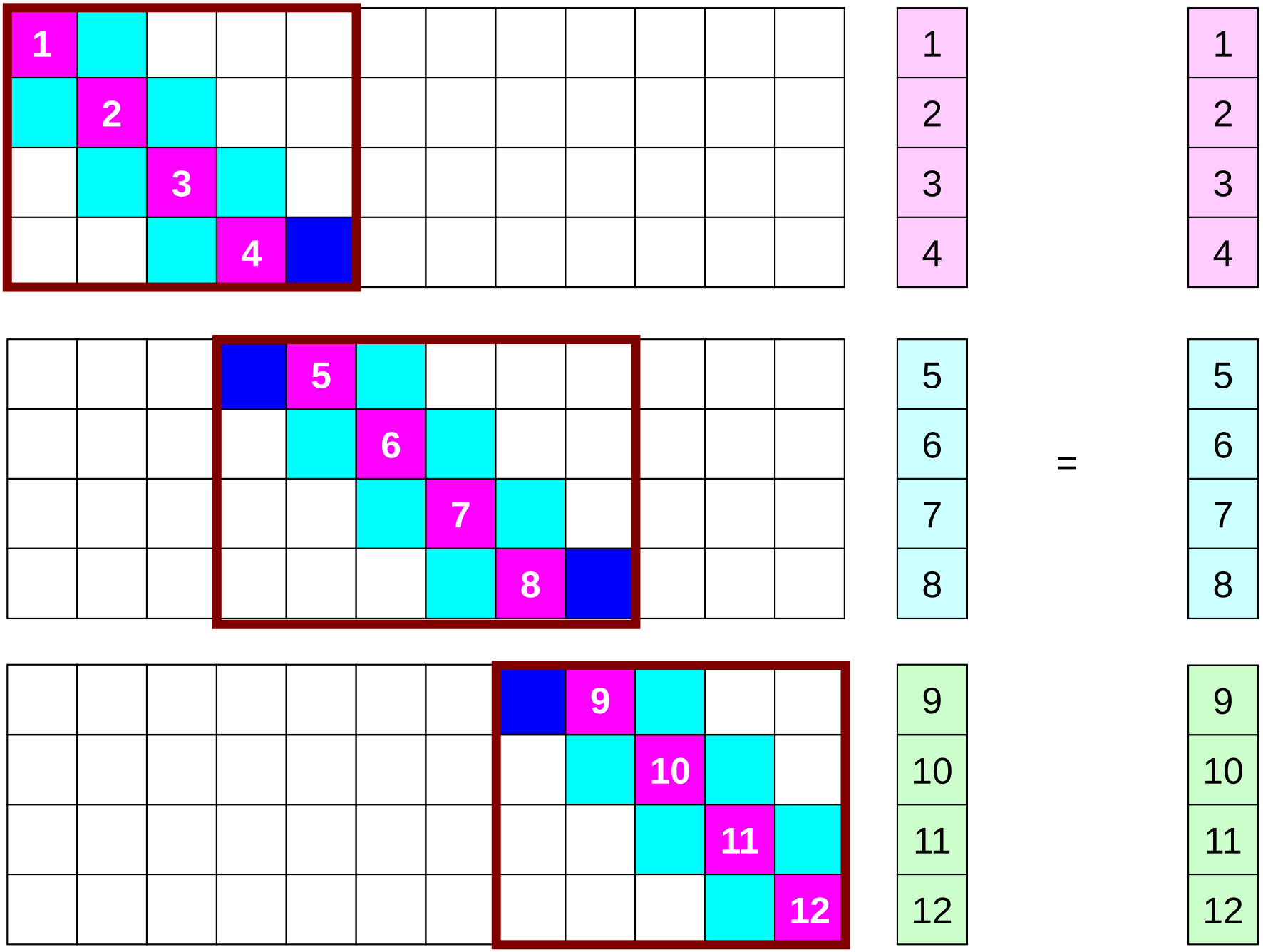
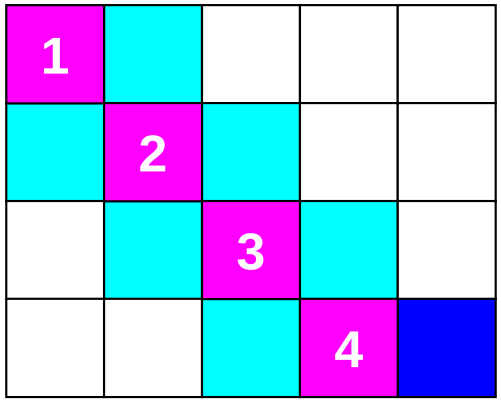
Because the matrix is sparse, the union of the local matrices forms the global matrix !



Mat-Vec Products: Local Op. Possible



Mat-Vec Products: Local Op. Possible



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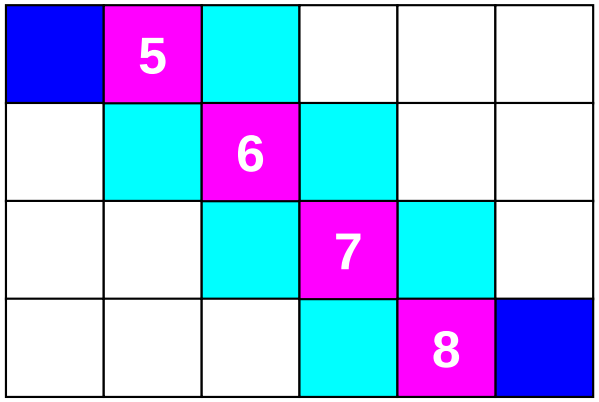

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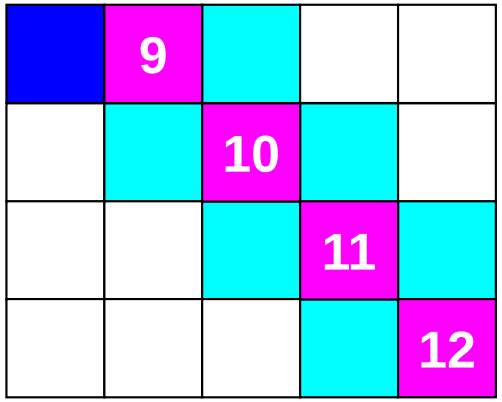

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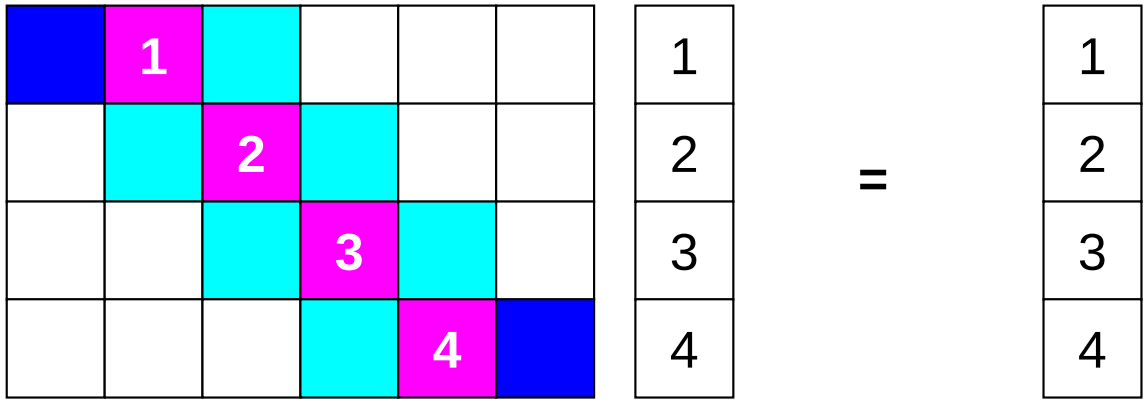
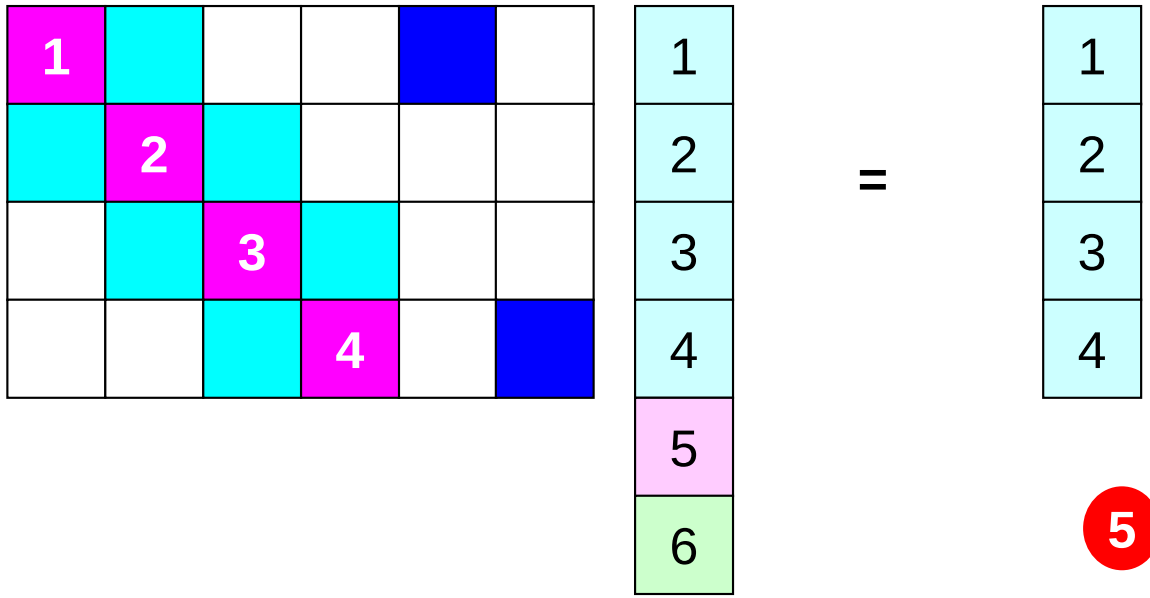
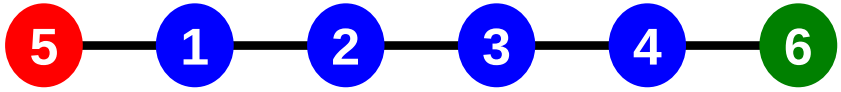
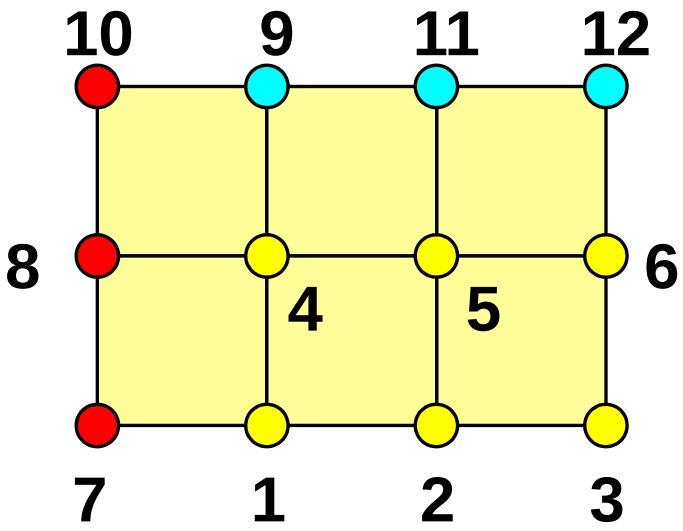
Mat-Vec Products: Local Op. #1







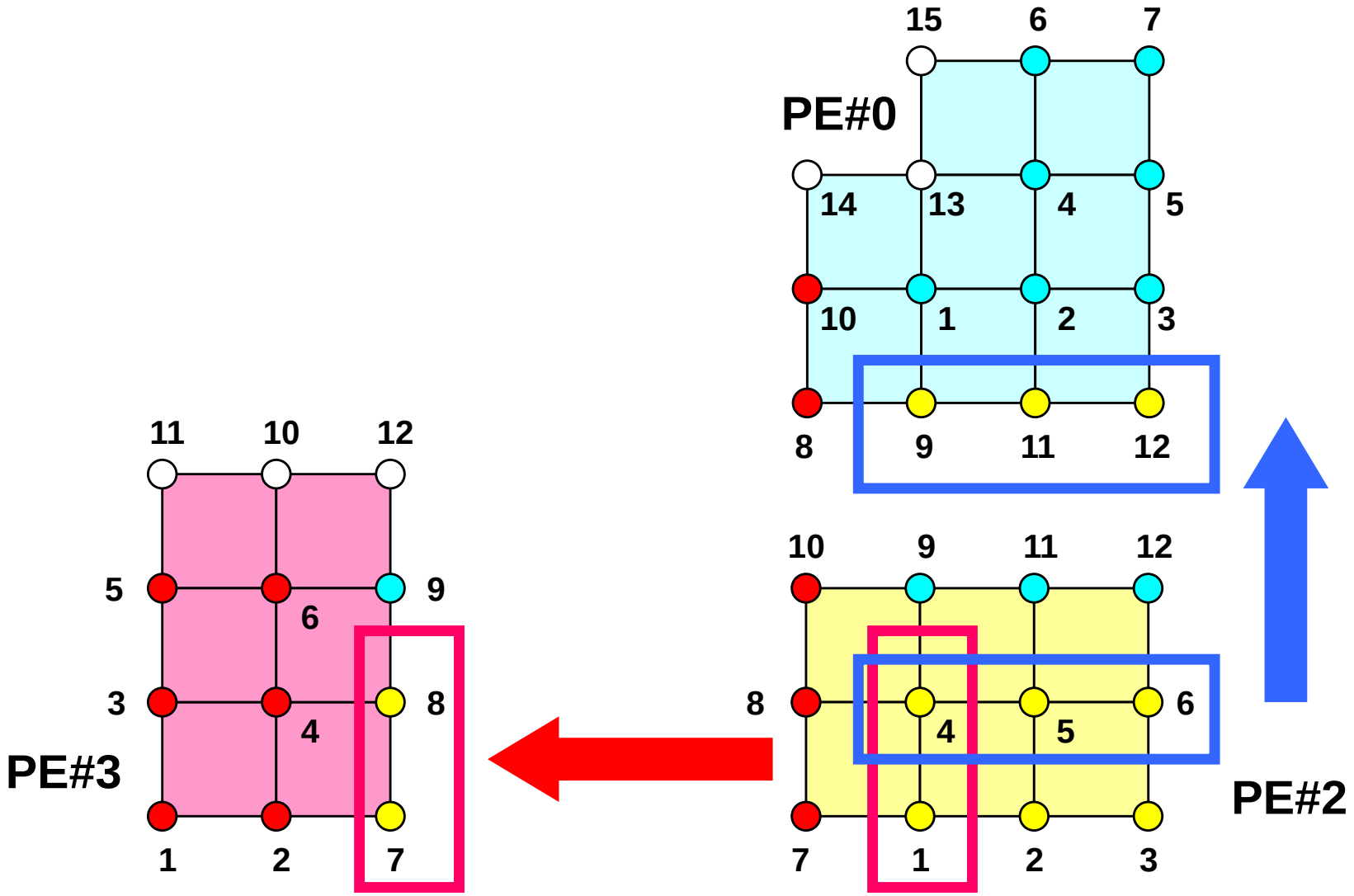
Description of Distributed Local Data



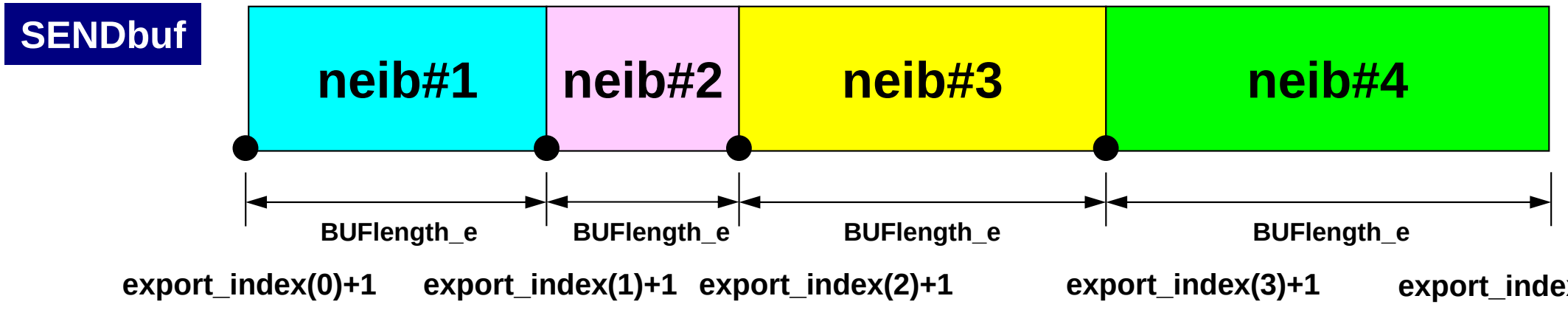
- Internal/External Points
  - Numbering: Starting from internal pts, then external pts after that
- Neighbors
  - Shares overlapped meshes
  - Number and ID of neighbors
- External Points
  - From where, how many, and which external points are received/imported ?
- Boundary Points
  - To where, how many and which boundary points are sent/exported ?

Boundary Nodes（境界点）：SEND

PE#2 : send information on “boundary nodes”



SEND: MPI_Isend/Irecv/Waitall

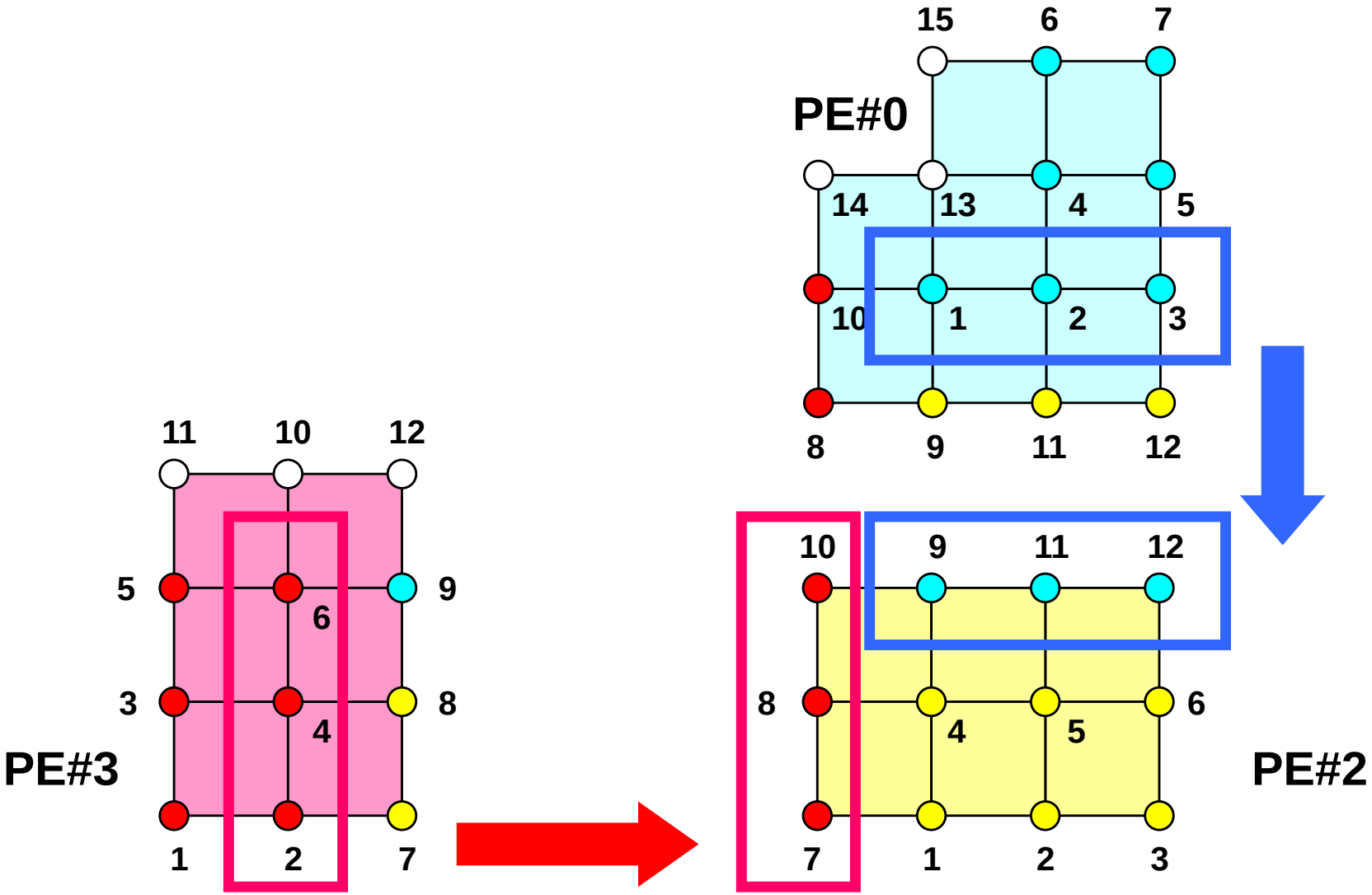



```
 do neib= 1, NEIBPETOT
   do k= export_index(neib-1)+1, export_index(neib)
     kk= export_item(k)
     SENDbuf(k)= VAL(kk)
   enddo
 enddo

 do neib= 1, NEIBPETOT
   iS_e= export_index(neib-1) + 1
   iE_e= export_index(neib  )
   BUFlength_e= iE_e + 1 - iS_e

   call MPI_ISEND                                                  &
&          (SENDbuf(iS_e), BUFlength_e, MPI_INTEGER, NEIBPE(neib), 0,&
&           MPI_COMM_WORLD, request_send(neib), ierr)
 enddo

 call MPI_WAITALL (NEIBPETOT, request_send, stat_send, ierr)
```

Copied to sending buffers

External Nodes（外点）：RECEIVE

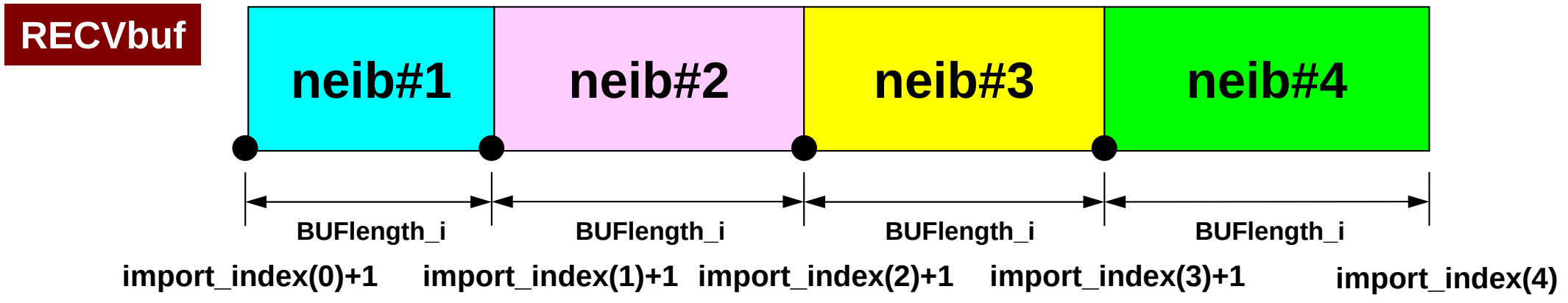
PE#2 : receive information for “external nodes”



RECV: MPI_Isend/Irecv/Waitall

```
 do neib= 1, NEIBPETOT
   iS_i= import_index(neib-1) + 1
   iE_i= import_index(neib  )
   BUFlength_i= iE_i + 1 - iS_i

   call MPI_IRECV                                                  &
&         (RECVbuf(iS_i), BUFlength_i, MPI_INTEGER, NEIBPE(neib), 0,&
&          MPI_COMM_WORLD, request_recv(neib), ierr)
  enddo

  call MPI_WAITALL (NEIBPETOT, request_recv, stat_recv, ierr)

  do neib= 1, NEIBPETOT
     do k= import_index(neib-1)+1, import_index(neib)
       kk= import_item(k)
       VAL(kk)= RECVbuf(k)
     enddo
  enddo
```

Copied from receiving buffer



- Overview
- Distributed Local Data
- **Program**
- Results

Program: 1d.f (1/11)

Variables

```
      program heat1Dp
      implicit REAL*8 (A-H,O-Z)
      include 'mpif.h'

      integer :: N, NPLU, ITERmax
      integer :: R, Z, P, Q, DD

      real(kind=8) :: dX, RESID, EPS
      real(kind=8) :: AREA, QV, COND
      real(kind=8), dimension(:), allocatable :: PHI, RHS
      real(kind=8), dimension(:  ), allocatable :: DIAG, AMAT
      real(kind=8), dimension(:,:), allocatable :: W

      real(kind=8), dimension(2,2) :: KMAT, EMAT

      integer, dimension(:), allocatable :: ICELNOD
      integer, dimension(:), allocatable :: INDEX, ITEM
      integer(kind=4) :: NEIBPETOT, BUFlength, PETOT
      integer(kind=4), dimension(2) :: NEIBPE

      integer(kind=4), dimension(0:2) :: import_index, export_index
      integer(kind=4), dimension(  2) :: import_item , export_item

      real(kind=8), dimension(2) :: SENDbuf, RECVbuf

      integer(kind=4), dimension(:,:), allocatable :: stat_send
      integer(kind=4), dimension(:,:), allocatable :: stat_recv
      integer(kind=4), dimension(:  ), allocatable :: request_send
      integer(kind=4), dimension(:  ), allocatable :: request_recv
```

Program: 1d.f (2/11)

Control Data

```
!C
!C +-------+
!C | INIT. |
!C +-------+
!C===
!C
!C-- MPI init.

      call MPI_Init      (ierr)                                  Initialization
      call MPI_Comm_size (MPI_COMM_WORLD, PETOT, ierr )          Entire Process #: PETOT
      call MPI_Comm_rank (MPI_COMM_WORLD, my_rank, ierr )        Rank ID (0-PETot-1): my_rank

!C
!C-- CTRL data
      if (my_rank.eq.0) thenn
        open  (11, file='input.dat', status='unknown')
         read (11,*) NEg
         read (11,*) dX, QV, AREA, COND
         read (11,*) ITERmax
         read (11,*) EPS
        close (11)
      endif

      call MPI_Bcast (NEg    , 1, MPI_INTEGER, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (ITERmax, 1, MPI_INTEGER, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (dX     , 1, MPI_DOUBLE_PRECISION, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (QV     , 1, MPI_DOUBLE_PRECISION, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (AREA   , 1, MPI_DOUBLE_PRECISION, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (COND   , 1, MPI_DOUBLE_PRECISION, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (EPS    , 1, MPI_DOUBLE_PRECISION, 0, MPI_COMM_WORLD, ierr)
```

Program: 1d.f (2/11)

Control Data

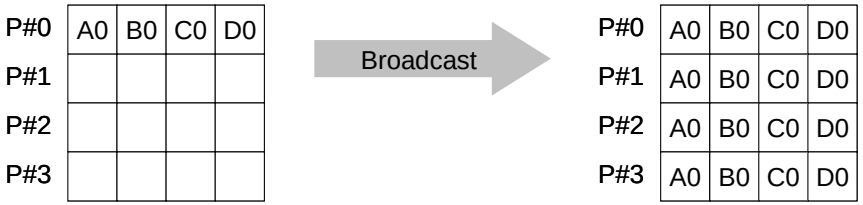
```
!C
!C +-------+
!C | INIT. |
!C +-------+
!C===
!C
!C-- MPI init.

      call MPI_Init      (ierr)                                  Initialization
      call MPI_Comm_size (MPI_COMM_WORLD, PETOT, ierr )          Entire Process #: PETOT
      call MPI_Comm_rank (MPI_COMM_WORLD, my_rank, ierr )        Rank ID (0-PETot-1): my_rank

!C
!C-- CTRL data
      if (my_rank.eq.0) then                                     Reading control file if my_rank=0
        open  (11, file='input.dat', status='unknown')
         read (11,*) Neg                                         Neg: Global Number of Elements
         read (11,*) dX, QV, AREA, COND
         read (11,*) ITERmax
         read (11,*) EPS
        close (11)
      endif

      call MPI_Bcast (NEg    , 1, MPI_INTEGER, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (ITERmax, 1, MPI_INTEGER, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (dX     , 1, MPI_DOUBLE_PRECISION, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (QV     , 1, MPI_DOUBLE_PRECISION, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (AREA   , 1, MPI_DOUBLE_PRECISION, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (COND   , 1, MPI_DOUBLE_PRECISION, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (EPS    , 1, MPI_DOUBLE_PRECISION, 0, MPI_COMM_WORLD, ierr)
```

Program: 1d.f (2/11)

Control Data

```
!C
!C +-------+
!C | INIT. |
!C +-------+
!C===
!C
!C-- MPI init.

      call MPI_Init      (ierr)                                  Initialization
      call MPI_Comm_size (MPI_COMM_WORLD, PETOT, ierr )          Entire Process #: PETOT
      call MPI_Comm_rank (MPI_COMM_WORLD, my_rank, ierr )        Rank ID (0-PETot-1): my_rank

!C
!C-- CTRL data
      if (my_rank.eq.0) then                                     Reading control file if my_rank=0
        open  (11, file='input.dat', status='unknown')
         read (11,*) Neg                                         Neg: Global Number of Elements
         read (11,*) dX, QV, AREA, COND
         read (11,*) ITERmax
         read (11,*) EPS
        close (11)
      endif

      call MPI_Bcast (NEg    , 1, MPI_INTEGER, 0, MPI_COMM_WORLD, ierr) Parameters are sent to each proces
      call MPI_Bcast (ITERmax, 1, MPI_INTEGER, 0, MPI_COMM_WORLD, ierr) from Process #0.
      call MPI_Bcast (dX     , 1, MPI_DOUBLE_PRECISION, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (QV     , 1, MPI_DOUBLE_PRECISION, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (AREA   , 1, MPI_DOUBLE_PRECISION, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (COND   , 1, MPI_DOUBLE_PRECISION, 0, MPI_COMM_WORLD, ierr)
      call MPI_Bcast (EPS    , 1, MPI_DOUBLE_PRECISION, 0, MPI_COMM_WORLD, ierr)
```

MPI_BCAST



- Broadcasts a message from the process with rank "root" to all other processes of the communicator

- **`call MPI_BCAST (buffer,count,datatype,root,comm,ierr)`**
  - **`buffer`** choice I/O — starting address of buffer
    type is defined by "**`datatype`**"
  - **`count`** I I — number of elements in send/recv buffer
  - **`datatype`** I I — data type of elements of send/recv buffer

```
FORTRAN  MPI_INTEGER, MPI_REAL, MPI_DOUBLE_PRECISION, MPI_CHARACTER etc.
C        MPI_INT, MPI_FLOAT, MPI_DOUBLE, MPI_CHAR etc.
```

  - **`root`** I I — rank of root process
  - **`comm`** I I — communicator
  - **`ierr`** I O — completion code

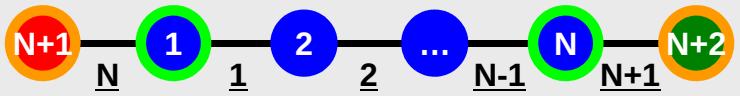
Program: 1d.f (3/11)
Distributed Local Mesh

```
!C
!C-- Local Mesh Size

      Ng= NEg + 1                          Global Number of Nodes
      N = Ng / PETOT                       Local Number of Nodes

      nr = Ng - N*PETOT                    mod(Ng,PETOT) .ne. 0
      if (my_rank.lt.nr) N= N+1

      NE= N - 1 + 2
      NP= N + 2

      if (my_rank.eq.0) NE= N - 1 + 1
      if (my_rank.eq.0) NP= N + 1

      if (my_rank.eq.PETOT-1) NE= N - 1 + 1
      if (my_rank.eq.PETOT-1) NP= N + 1

      if (PETOT.eq.1) NE= N-1
      if (PETOT.eq.1) NP= N

!C
!C- ARRAYs

      allocate (PHI(NP), DIAG(NP), AMAT(2*NP-2), RHS(NP))
      allocate (ICELNOD(2*NE))
      allocate (INDEX(0:NP), ITEM(2*NP-2), W(NP,4))
       PHI= 0.d0
      AMAT= 0.d0
      DIAG= 0.d0
       RHS= 0.d0
```

Program: 1d.f (3/11)

Distributed Local Mesh, Uniform Elements

```
!C
!C-- Local Mesh Size

      Ng= NEg + 1                          Global Number of Nodes
      N = Ng / PETOT                       Local Number of Nodes

      nr = Ng - N*PETOT                    mod(Ng,PETOT) .ne. 0
      if (my_rank.lt.nr) N= N+1

      NE= N - 1 + 2                        Number of Elements (Local)
      NP= N + 2                            Total Number of Nodes (Local) (Internal + External Nodes)

      if (my_rank.eq.0) NE= N - 1 + 1
      if (my_rank.eq.0) NP= N + 1

      if (my_rank.eq.PETOT-1) NE= N - 1 + 1
      if (my_rank.eq.PETOT-1) NP= N + 1

      if (PETOT.eq.1) NE= N-1
      if (PETOT.eq.1) NP= N
```



Others (General):
N+2 nodes
N+1 elements

```
!C
!C- ARRAYs

      allocate (PHI(NP), DIAG(NP), AMAT(2*NP-2), RHS(NP))
      allocate (ICELNOD(2*NE))
      allocate (INDEX(0:NP), ITEM(2*NP-2), W(NP,4))
       PHI= 0.d0
      AMAT= 0.d0
      DIAG= 0.d0
       RHS= 0.d0
```

Program: 1d.f (3/11)

Distributed Local Mesh, Uniform Elements

```
!C
!C-- Local Mesh Size

      Ng= NEg + 1                          Global Number of Nodes
      N = Ng / PETOT                       Local Number of Nodes

      nr = Ng - N*PETOT                    mod(Ng,PETOT) .ne. 0
      if (my_rank.lt.nr) N= N+1

      NE= N - 1 + 2                        Number of Elements (Local)
      NP= N + 2                            Total Number of Nodes (Local) (Internal + External Nodes)

      if (my_rank.eq.0) NE= N - 1 + 1
      if (my_rank.eq.0) NP= N + 1

      if (my_rank.eq.PETOT-1) NE= N - 1 + 1
      if (my_rank.eq.PETOT-1) NP= N + 1

      if (PETOT.eq.1) NE= N-1
      if (PETOT.eq.1) NP= N

!C
!C- ARRAYs

      allocate (PHI(NP), DIAG(NP), AMAT(2*NP-2), RHS(NP))
      allocate (ICELNOD(2*NE))
      allocate (INDEX(0:NP), ITEM(2*NP-2), W(NP,4))
       PHI= 0.d0
      AMAT= 0.d0
      DIAG= 0.d0
       RHS= 0.d0
```

#0:
N+1 nodes
N elements

Program: 1d.f (3/11)
Distributed Local Mesh, Uniform Elem

```
!C
!C-- Local Mesh Size

      Ng= NEg + 1                          Global Number of Nodes
      N = Ng / PETOT                       Local Number of Nodes

      nr = Ng - N*PETOT                    mod(Ng,PETOT) .ne. 0
      if (my_rank.lt.nr) N= N+1

      NE= N - 1 + 2                        Number of Elements (Local)
      NP= N + 2                            Total Number of Nodes (Local) (Internal + External Nodes)

      if (my_rank.eq.0) NE= N - 1 + 1
      if (my_rank.eq.0) NP= N + 1

      if (my_rank.eq.PETOT-1) NE= N - 1 + 1
      if (my_rank.eq.PETOT-1) NP= N + 1

      if (PETOT.eq.1) NE= N-1
      if (PETOT.eq.1) NP= N

!C
!C- ARRAYs

      allocate (PHI(NP), DIAG(NP), AMAT(2*NP-2), RHS(NP))
      allocate (ICELNOD(2*NE))
      allocate (INDEX(0:NP), ITEM(2*NP-2), W(NP,4))
       PHI= 0.d0
      AMAT= 0.d0
      DIAG= 0.d0
       RHS= 0.d0
```

#PETot-1:
N+1 nodes
N elements

Program: 1d.f (3/11)

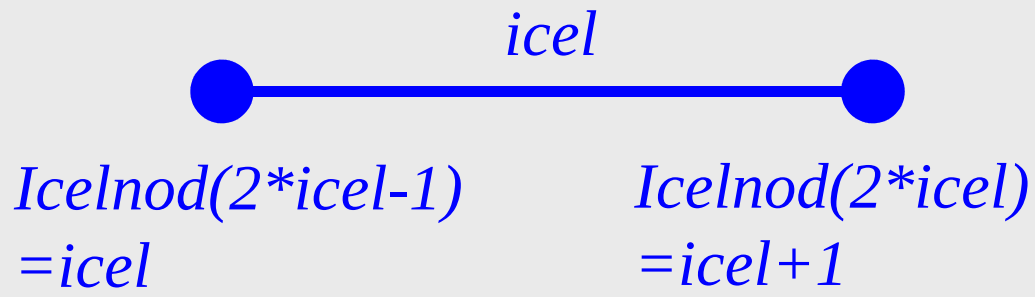
Distributed Local Mesh, Uniform Elements

```
!C
!C-- Local Mesh Size

      Ng= NEg + 1                          Global Number of Nodes
      N = Ng / PETOT                       Local Number of Nodes

      nr = Ng - N*PETOT                    mod(Ng,PETOT) .ne. 0
      if (my_rank.lt.nr) N= N+1

      NE= N - 1 + 2                        Number of Elements (Local)
      NP= N + 2                            Total Number of Nodes (Local) (Internal + External Nodes)

      if (my_rank.eq.0) NE= N - 1 + 1
      if (my_rank.eq.0) NP= N + 1

      if (my_rank.eq.PETOT-1) NE= N - 1 + 1
      if (my_rank.eq.PETOT-1) NP= N + 1

      if (PETOT.eq.1) NE= N-1
      if (PETOT.eq.1) NP= N

!C
!C- ARRAYs

      allocate (PHI(NP), DIAG(NP), AMAT(2*NP-2), RHS(NP))    Size of arrays is "NP", not "N"
      allocate (ICELNOD(2*NE))
      allocate (INDEX(0:NP), ITEM(2*NP-2), W(NP,4))
       PHI= 0.d0
      AMAT= 0.d0
      DIAG= 0.d0
       RHS= 0.d0
```

Program: 1d.f (4/11)

Initialization of Arrays, Elements-Nodes

```
do icel= 1, NE
  ICELNOD(2*icel-1)= icel
  ICELNOD(2*icel  )= icel + 1
enddo

if (PETOT.gt.1) then

if (my_rank.eq.0) then
  icel= NE
  ICELNOD(2*icel-1)= N
  ICELNOD(2*icel  )= N + 1


 else if (my_rank.eq.PETOT-1) then
  icel= NE
  ICELNOD(2*icel-1)= N + 1
  ICELNOD(2*icel  )= 1


 else
  icel= NE - 1
  ICELNOD(2*icel-1)= N + 1
  ICELNOD(2*icel  )= 1
  icel= NE
  ICELNOD(2*icel-1)= N
  ICELNOD(2*icel  )= N + 2

endif
endif
```

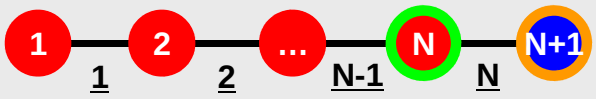


Program: 1d.f (4/11)

Initialization of Arrays, Elements-Nodes

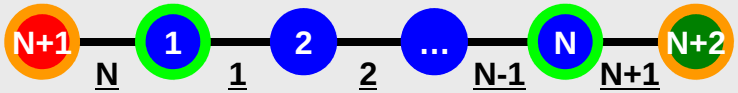
```
do icel= 1, NE
  ICELNOD(2*icel-1)= icel
  ICELNOD(2*icel  )= icel + 1
enddo

if (PETOT.gt.1) then

if (my_rank.eq.0) then
  icel= NE
  ICELNOD(2*icel-1)= N
  ICELNOD(2*icel  )= N + 1


 else if (my_rank.eq.PETOT-1) then
  icel= NE
  ICELNOD(2*icel-1)= N + 1
  ICELNOD(2*icel  )= 1


 else
  icel= NE - 1
  ICELNOD(2*icel-1)= N + 1
  ICELNOD(2*icel  )= 1
  icel= NE
  ICELNOD(2*icel-1)= N
  ICELNOD(2*icel  )= N + 2

endif
endif
```

e.g. Element-1 includes node-1 and node-2



#0:
N+1 nodes
N elements

#PETot-1:
N+1 nodes
N elements



Others (General):
N+2 nodes
N+1 elements

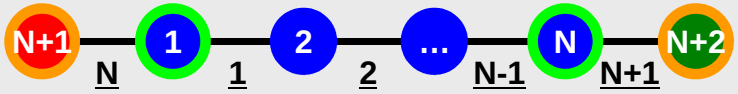
Program: 1d.f (5/11)

"Index"

```
 KMAT(1,1)= +1.d0
      KMAT(1,2)= -1.d0
      KMAT(2,1)= -1.d0
      KMAT(2,2)= +1.d0
!C===

!C
!C +-------------+
!C | CONNECTIVITY |
!C +-------------+
!C===
      INDEX   = 2

      INDEX(0)= 0

      INDEX(N+1)= 1
      INDEX(NP )= 1

      if (my_rank.eq.0)         INDEX(1)= 1
      if (my_rank.eq.PETOT-1)   INDEX(N)= 1

      do i= 1, NP
        INDEX(i)= INDEX(i) + INDEX(i-1)
      enddo

      NPLU= INDEX(NP)
      ITEM= 0
```

#0:
N+1 nodes
N elements

#PETot-1:
N+1 nodes
N elements



Others (General):
N+2 nodes
N+1 elements

Program: 1d.f (6/11)

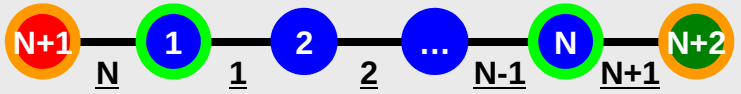
”Item”

```
do i= 1, N
  jS= INDEX(i-1)
  if (my_rank.eq.0.and.i.eq.1) then
    ITEM(jS+1)= i+1
   else if (my_rank.eq.PETOT-1.and.i.eq.N) then
    ITEM(jS+1)= i-1
   else
    ITEM(jS+1)= i-1
    ITEM(jS+2)= i+1
    if (i.eq.1) ITEM(jS+1)= N + 1
    if (i.eq.N) ITEM(jS+2)= N + 2
    if (my_rank.eq.0.and.i.eq.N) ITEM(jS+2)= N + 1
  endif
enddo
```

#0:
N+1 nodes
N elements

```
i = N + 1
jS= INDEX(i-1)
if (my_rank.eq.0) then
  ITEM(jS+1)= N
 else
  ITEM(jS+1)= 1
endif
```

#PETot-1:
N+1 nodes
N elements



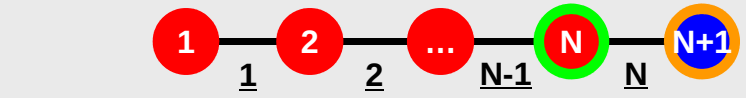
Others (General):
N+2 nodes
N+1 elements

```
i = N + 2
if (my_rank.ne.0.and.my_rank.ne.PETOT-1) then
  jS= INDEX(i-1)
  ITEM(jS+1)= N
endif
```

Program: 1d.f (7/11)

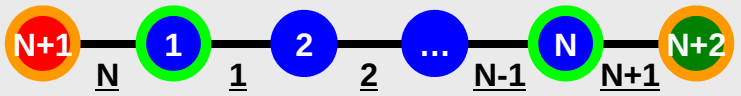
Communication Tables

```
!C
!C-- COMMUNICATION
      NEIBPETOT= 2
      if (my_rank.eq.0      ) NEIBPETOT= 1
      if (my_rank.eq.PETOT-1) NEIBPETOT= 1
      if (PETOT.eq.1)         NEIBPETOT= 0

      NEIBPE(1)= my_rank - 1
      NEIBPE(2)= my_rank + 1

      if (my_rank.eq.0      ) NEIBPE(1)= my_rank + 1
      if (my_rank.eq.PETOT-1) NEIBPE(1)= my_rank - 1

      BUFlength= 1

      import_index(1)= 1
      import_index(2)= 2
      import_item (1)= N+1
      import_item (2)= N+2

      export_index(1)= 1
      export_index(2)= 2
      export_item (1)= 1
      export_item (2)= N

      if (my_rank.eq.0) then
        import_item (1)= N+1
        export_item (1)= N
      endif
!C
!C-- INIT. arrays for MPI_Waitall
      allocate (stat_send(MPI_STATUS_SIZE,NEIBPETOT), stat_recv(MPI_STATUS_SIZE,NEIBPETOT))
      allocate (request_send(NEIBPETOT), request_recv(NEIBPETOT))
```



#0:
N+1 nodes
N elements

#PETot-1:
N+1 nodes
N elements



Others (General):
N+2 nodes
N+1 elements

MPI_ISEND

- Begins a non-blocking send
  - Send the contents of sending buffer (starting from **sendbuf**, number of messages: **count**) to **dest** with **tag** .
  - Contents of sending buffer cannot be modified before calling corresponding **MPI_Waitall**.

- **call MPI_ISEND (sendbuf,count,datatype,dest,tag,comm,request, ierr)**
  - **sendbuf** choice I starting address of sending buffer
  - **count** I I number of elements sent to each process
  - **datatype** I I data type of elements of sending buffer
  - **dest** I I rank of destination
  - **tag** I I message tag
    This integer can be used by the application to distinguish messages. Communication occurs if `tag's` of `MPI_Isend` and `MPI_Irecv` are matched.
    Usually tag is set to be “0” (in this class),
  - **comm** I I communicator
  - **request** I O communication request array used in `MPI_Waitall`
  - **ierr** I O completion code

MPI_IRECV

- Begins a non-blocking receive
  - Receiving the contents of receiving buffer (starting from **`recvbuf`**, number of messages: **`count`**) from **`source`** with **`tag`** .
  - Contents of receiving buffer cannot be used before calling corresponding **`MPI_Waitall`**.

- **`call MPI_IRECV`**
  `(recvbuf,count,datatype,dest,tag,comm,request, ierr)`
  - **`recvbuf`** choice I starting address of receiving buffer
  - **`count`** I I number of elements in receiving buffer
  - **`datatype`** I I data type of elements of receiving buffer
  - **`source`** I I rank of source
  - **`tag`** I I message tag
    This integer can be used by the application to distinguish messages. Communication occurs if `tag's` of `MPI_Isend` and `MPI_Irecv` are matched.
    Usually tag is set to be “0” (in this class),
  - **`comm`** I I communicator
  - **`request`** I O communication request used in `MPI_Waitall`
  - **`ierr`** I O completion code

MPI_WAITALL

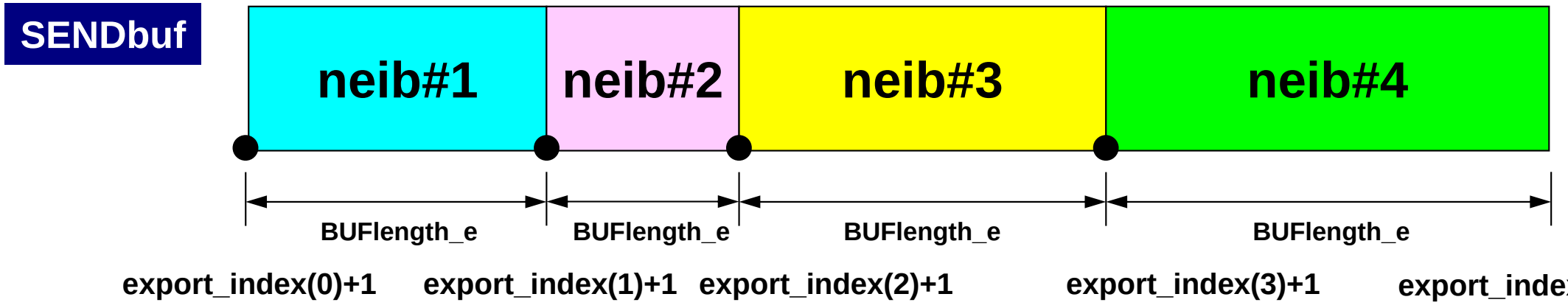
- **MPI_Waitall** blocks until all comm's, associated with **request** in the array, complete. It is used for synchronizing **MPI_Isend** and **MPI_Irecv** in this class.
- At sending phase, contents of sending buffer cannot be modified before calling corresponding **MPI_Waitall**. At receiving phase, contents of receiving buffer cannot be used before calling corresponding **MPI_Waitall**.
- **MPI_Isend** and **MPI_Irecv** can be synchronized simultaneously with a single **MPI_Waitall** if it is consitent.
  - Same **request** should be used in **MPI_Isend** and **MPI_Irecv**.
- Its operation is similar to that of **MPI_Barrier** but, **MPI_Waitall** can not be replaced by **MPI_Barrier.**
  - Possible troubles using **MPI_Barrier** instead of **MPI_Waitall**: Contents of **request** and **status** are not updated properly, very slow operations etc.

- **call MPI_WAITALL (count,request,status,ierr)**
  - **count** I I number of processes to be synchronized
  - **request** I I/O comm. request used in `MPI_Waitall` (array size: `count`)
  - **status** I O array of status objects
    `MPI_STATUS_SIZE: defined in 'mpif.h', 'mpi.h'`
  - **ierr** I O completion code

Generalized Comm. Table: Send

- Neighbors
  - NEIBPETOT, NEIBPE(neib)
- Message size for each neighbor
  - export_index(neib), neib= 0, NEIBPETOT
- ID of **boundary** points
  - export_item(k), k= 1, export_index(NEIBPETOT)
- Messages to each neighbor
  - SENDbuf(k), k= 1, export_index(NEIBPETOT)

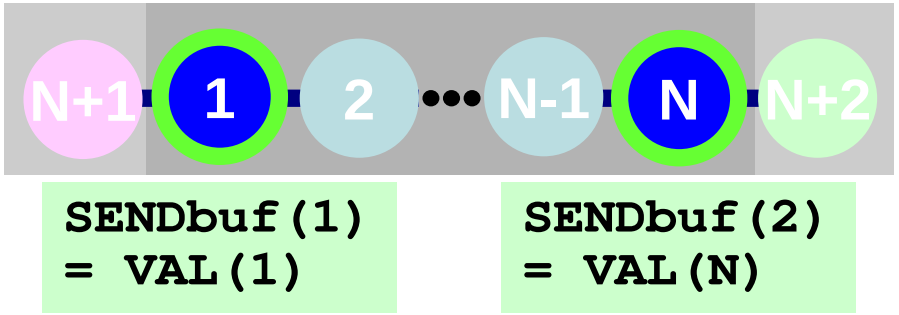
SEND: MPI_Isend/Irecv/Waitall




```
 do neib= 1, NEIBPETOT
   do k= export_index(neib-1)+1, export_index(neib)
     kk= export_item(k)
     SENDbuf(k)= VAL(kk)
   enddo
 enddo

 do neib= 1, NEIBPETOT
   iS_e= export_index(neib-1) + 1
   iE_e= export_index(neib  )
   BUFlength_e= iE_e + 1 - iS_e

   call MPI_ISEND                                                  &
&         (SENDbuf(iS_e), BUFlength_e, MPI_INTEGER, NEIBPE(neib), 0,&
&          MPI_COMM_WORLD, request_send(neib), ierr)
 enddo

 call MPI_WAITALL (NEIBPETOT, request_send, stat_send, ierr)
```

Copied to sending buffers

SEND/Export: 1D Problem



- Neighbors
  - NEIBPETOT, NEIBPE(neib)
    - NEIBPETOT=2, NEIB(1)= my_rank-1, NEIB(2)= my_rank+1
- Message size for each neighbor
  - export_index(neib), neib= 0, NEIBPETOT
    - export_index(0)=0, export_index(1)= 1, export_index(2)= 2
- ID of **boundary** points
  - export_item(k), k= 1, export_index(NEIBPETOT)
    - export_item(1)= 1, export_item(2)= N
- Messages to each neighbor
  - SENDbuf(k), k= 1, export_index(NEIBPETOT)
    - SENDbuf(1)= BUF(1), SENDbuf(2)= BUF(N)

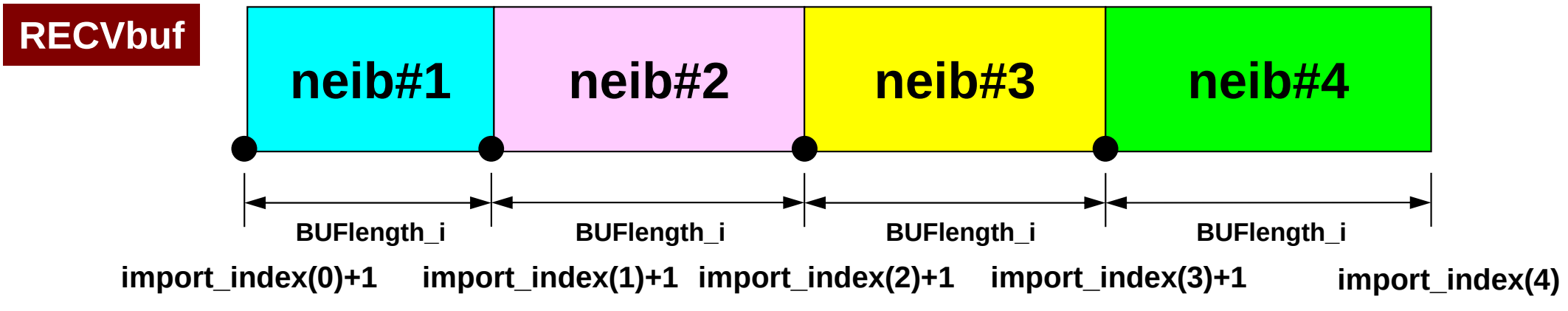
Generalized Comm. Table: Receive

- Neighbors
  - NEIBPETOT, NEIBPE(neib)
- Message size for each neighbor
  - import_index(neib), neib= 0, NEIBPETOT
- ID of **external** points
  - import_item(k), k= 1, import_index(NEIBPETOT)
- Messages from each neighbor
  - RECVbuf(k), k= 1, import_index(NEIBPETOT)

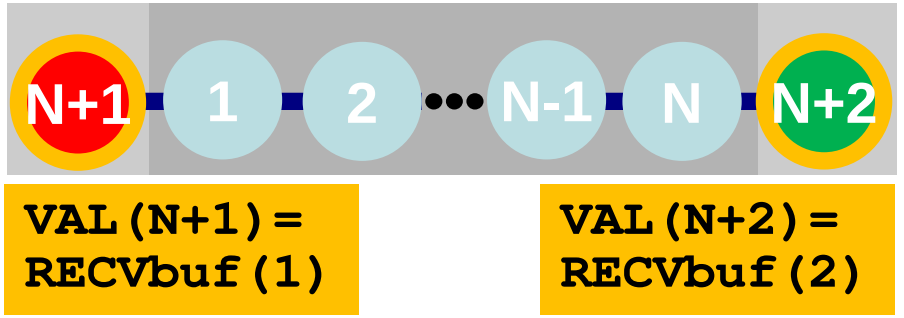
RECV: MPI_Isend/Irecv/Waitall

```
 do neib= 1, NEIBPETOT
   iS_i= import_index(neib-1) + 1
   iE_i= import_index(neib  )
   BUFlength_i= iE_i + 1 - iS_i

   call MPI_IRECV                                                  &
&         (RECVbuf(iS_i), BUFlength_i, MPI_INTEGER, NEIBPE(neib), 0,&
&          MPI_COMM_WORLD, request_recv(neib), ierr)
  enddo

  call MPI_WAITALL (NEIBPETOT, request_recv, stat_recv, ierr)

  do neib= 1, NEIBPETOT
     do k= import_index(neib-1)+1, import_index(neib)
       kk= import_item(k)
       VAL(kk)= RECVbuf(k)
     enddo
  enddo
```

Copied from receiving buffer

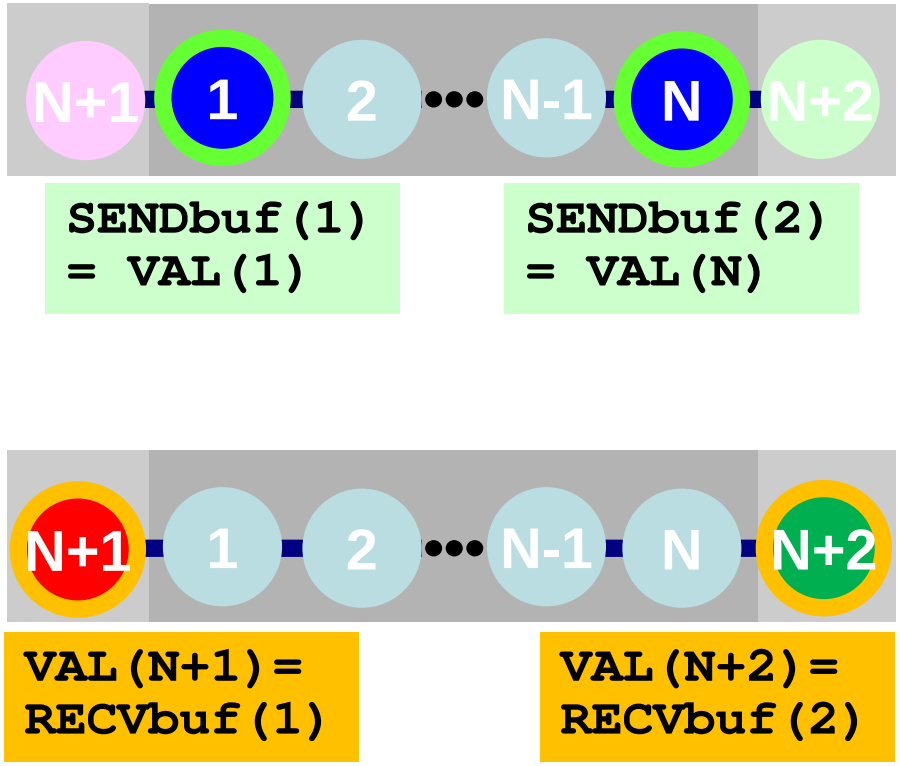


RECV/Import: 1D Problem



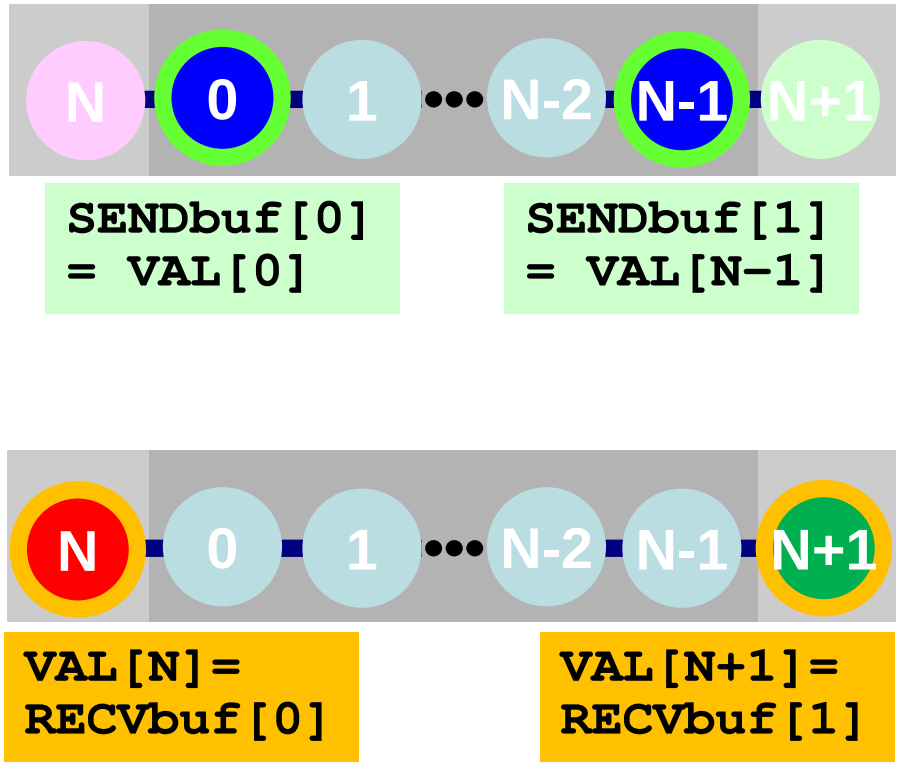
- Neighbors
  - NEIBPETOT, NEIBPE(neib)
    - NEIBPETOT=2, NEIB(1)= my_rank-1, NEIB(2)= my_rank+1
- Message size for each neighbor
  - import_index(neib), neib= 0, NEIBPETOT
    - import_index(0)=0, import_index(1)= 1, import_index(2)= 2
- ID of **external** points
  - import_item(k), k= 1, import_index(NEIBPETOT)
    - import_item(1)= N+1, import_item(2)= N+2
- Messages from each neighbor
  - RECVbuf(k), k= 1, import_index(NEIBPETOT)
    - BUF(N+1)=RECVbuf(1), BUF(N+2)=RECVbuf(2)

Generalized Comm. Table: Fortran




```
NEIBPETOT= 2
NEIBPE(1)= my_rank - 1
NEIBPE(2)= my_rank + 1

import_index(1)= 1
import_index(2)= 2
import_item  (1)= N+1
import_item  (2)= N+2

export_index(1)= 1
export_index(2)= 2
export_item  (1)= 1
export_item  (2)= N

if (my_rank.eq.0) then
  import_item  (1)= N+1
  export_item  (1)= N
  NEIBPE(1)= my_rank+1
endif
```

Generalized Comm. Table: C

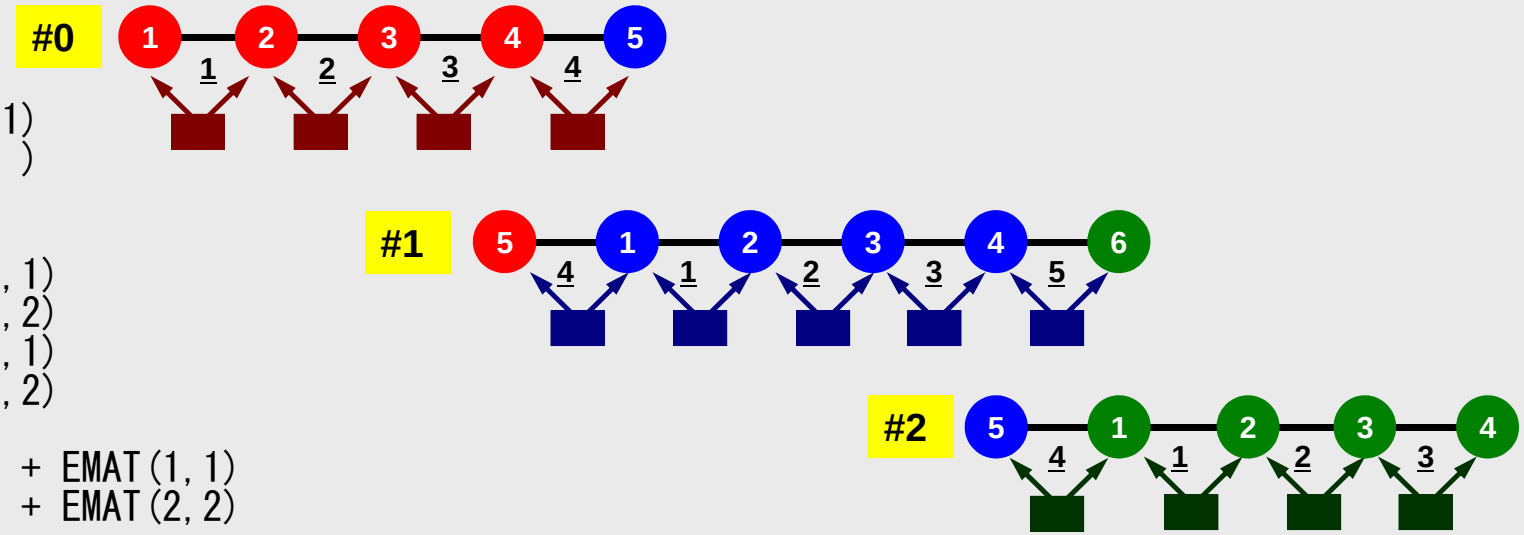



```
NEIBPETOT= 2
NEIBPE[0]= my_rank - 1
NEIBPE[1]= my_rank + 1

import_index[1]= 1
import_index[2]= 2
import_item [0]= N
import_item [1]= N+1

export_index[1]= 1
export_index[2]= 2
export_item [0]= 0
export_item [1]= N-1

if (my_rank.eq.0) then
  import_item [0]= N
  export_item [0]= N-1
  NEIBPE[0]= my_rank+1
endif
```

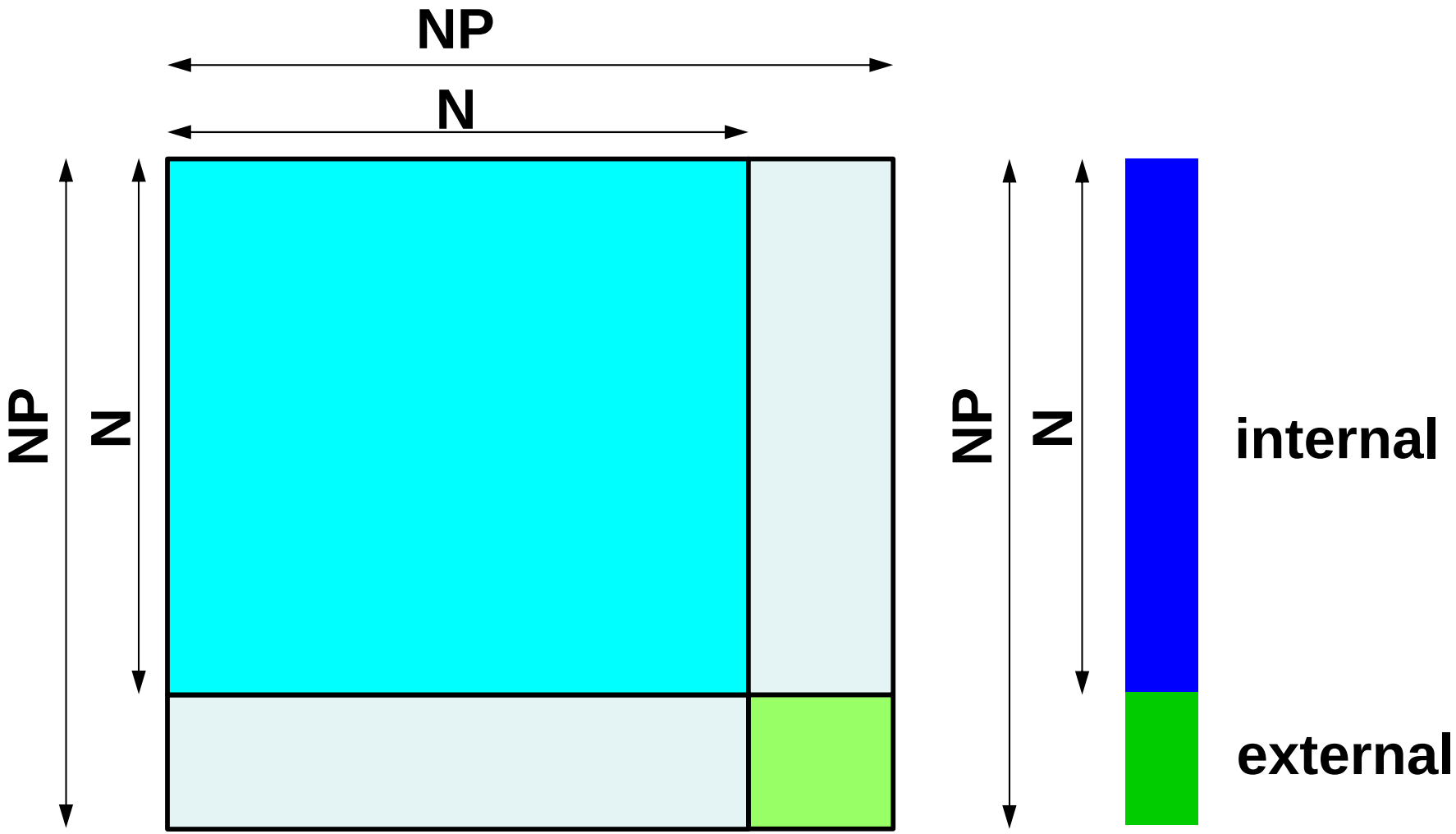
Program: 1d.f (8/11)

Matrix Assembling, NO changes from 1-CPU co

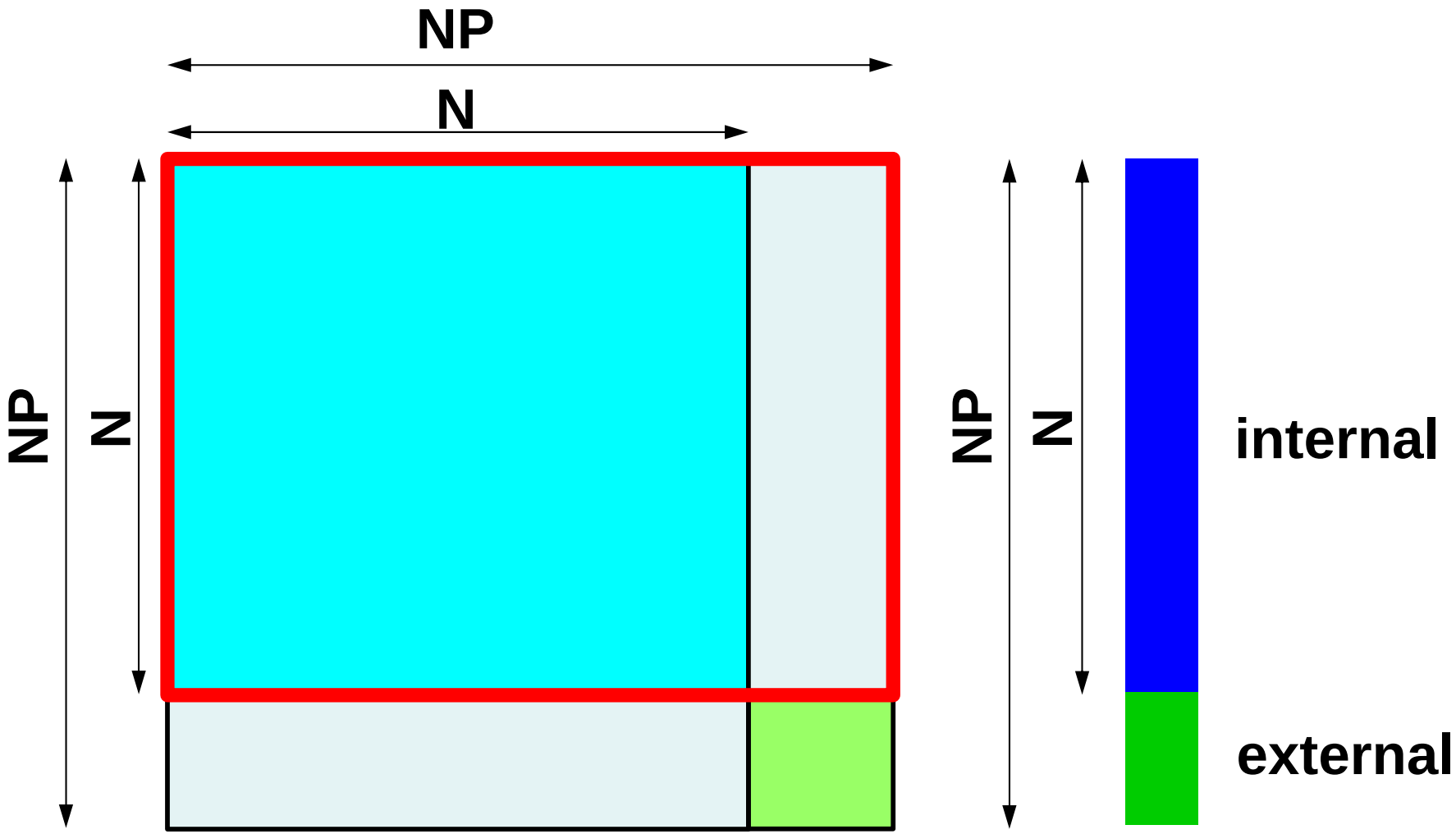
```
!C
!C +----------------+
!C | MATRIX ASSEMBLE |
!C +----------------+
!C===
      do icel= 1, NE
        in1= ICELNOD(2*icel-1)
        in2= ICELNOD(2*icel  )
        DL = dX
        cK= AREA*COND/DL
        EMAT(1,1)= Ck*KMAT(1,1)
        EMAT(1,2)= Ck*KMAT(1,2)
        EMAT(2,1)= Ck*KMAT(2,1)
        EMAT(2,2)= Ck*KMAT(2,2)

        DIAG(in1)= DIAG(in1) + EMAT(1,1)
        DIAG(in2)= DIAG(in2) + EMAT(2,2)

        if (my_rank.eq.0.and.icel.eq.1) then
          k1= INDEX(in1-1) + 1
         else
          k1= INDEX(in1-1) + 2
        endif
        k2= INDEX(in2-1) + 1

        AMAT(k1)= AMAT(k1) + EMAT(1,2)
        AMAT(k2)= AMAT(k2) + EMAT(2,1)

        QN= 0.50d0*QV*AREA*DL
        RHS(in1)= RHS(in1) + QN
        RHS(in2)= RHS(in2) + QN
      enddo
!C===
```



Local Matrix



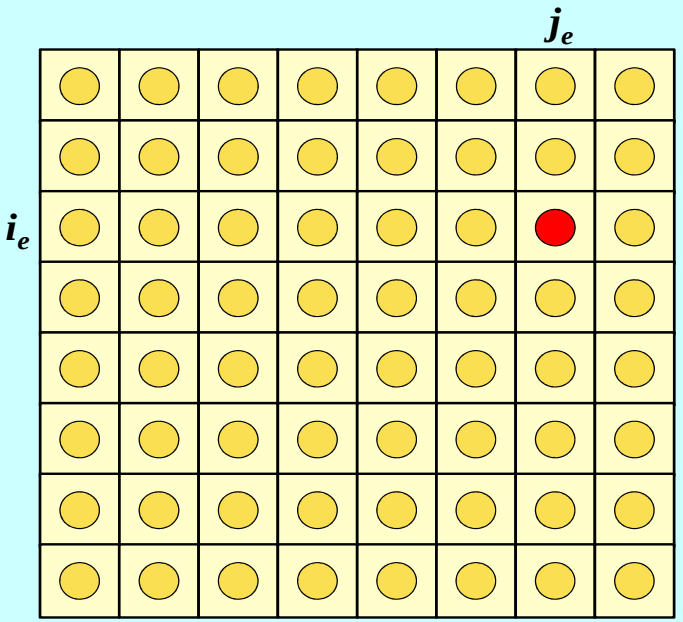
We really need these parts:



MAT_ASS_MAIN: Overview

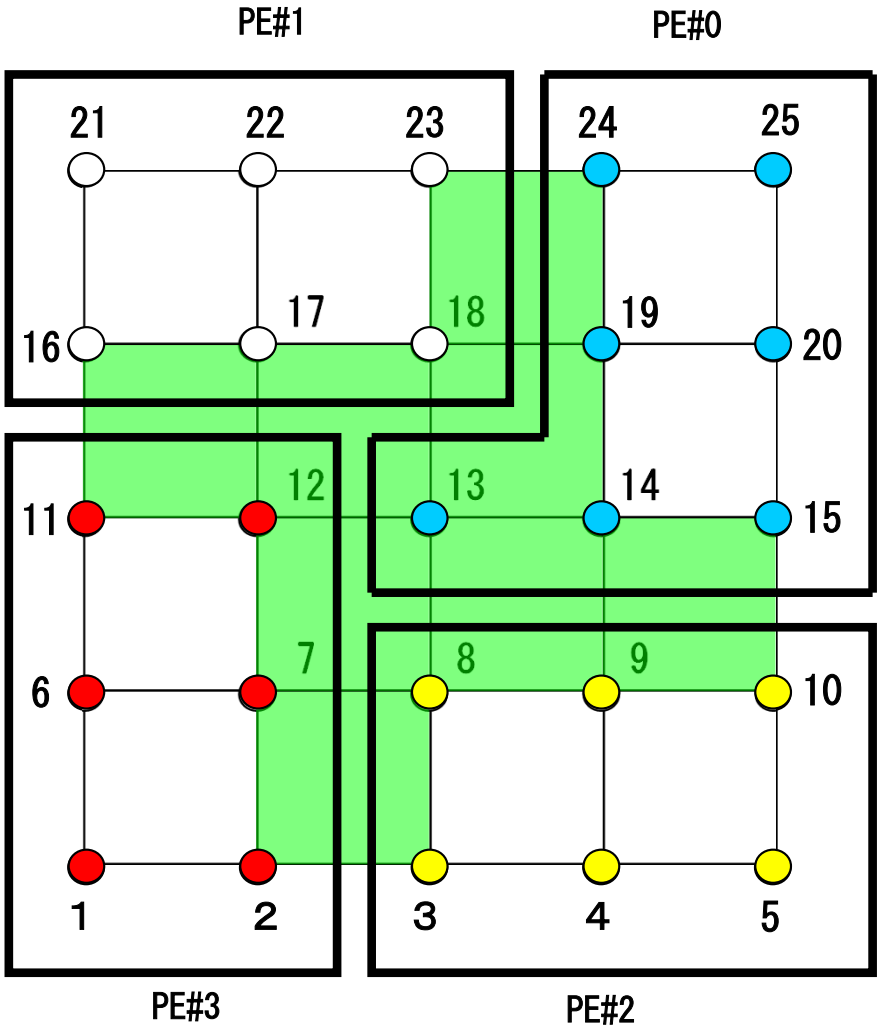
```
do kpn= 1, 2          Gaussian Quad. points in ζ-direction
  do jpn= 1, 2        Gaussian Quad. points in η-direction
    do ipn= 1, 2      Gaussian Quad. Pointe in ξ-direction
      Define Shape Function at Gaussian Quad. Points (8-points)
      Its derivative on natural/local coordinate is also defined.
    enddo
  enddo
enddo
```

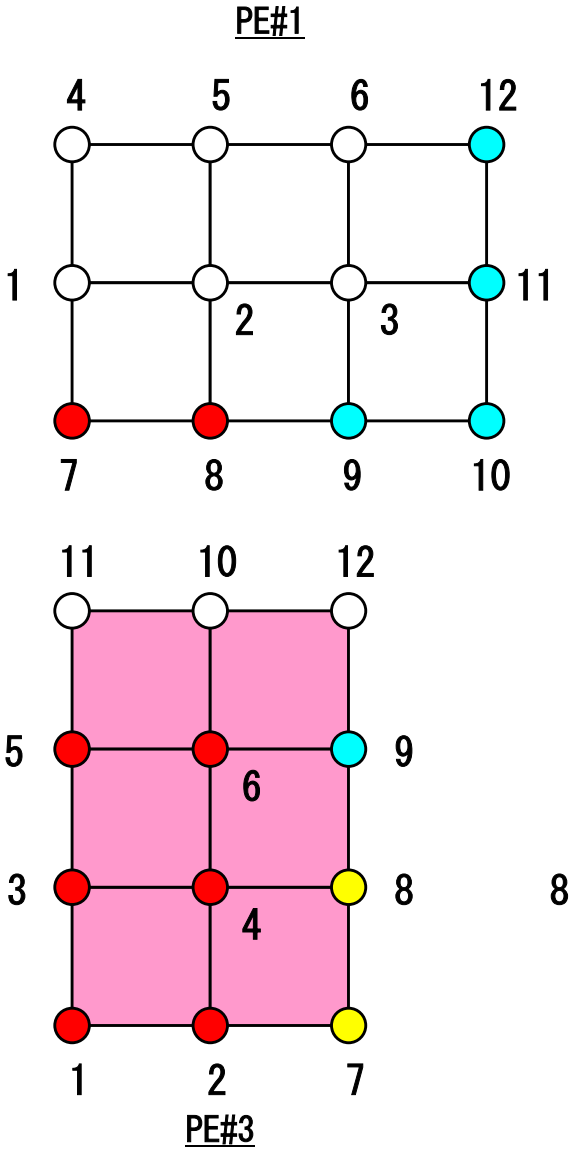
```
do icel= 1, ICELTOT    Loop for Element
  Jacobian and derivative on global coordinate of shape functions at
  Gaussian Quad. Points are defined according to coordinates of 8 nodes. (JACOBI)

  do ie= 1, 8          Local Node ID
    do je= 1, 8        Local Node ID
      Global Node ID: ip, jp
      Address of A_ip,jp in "item" : kk

      do kpn= 1, 2       Gaussian Quad. points in ζ-direction
        do jpn= 1, 2     Gaussian Quad. points in η-direction
          do ipn= 1, 2   Gaussian Quad. points in ξ-direction
            integration on each element
            coefficients of element matrices
            accumulation to global matrix
          enddo
        enddo
      enddo
    enddo
  enddo
enddo
```



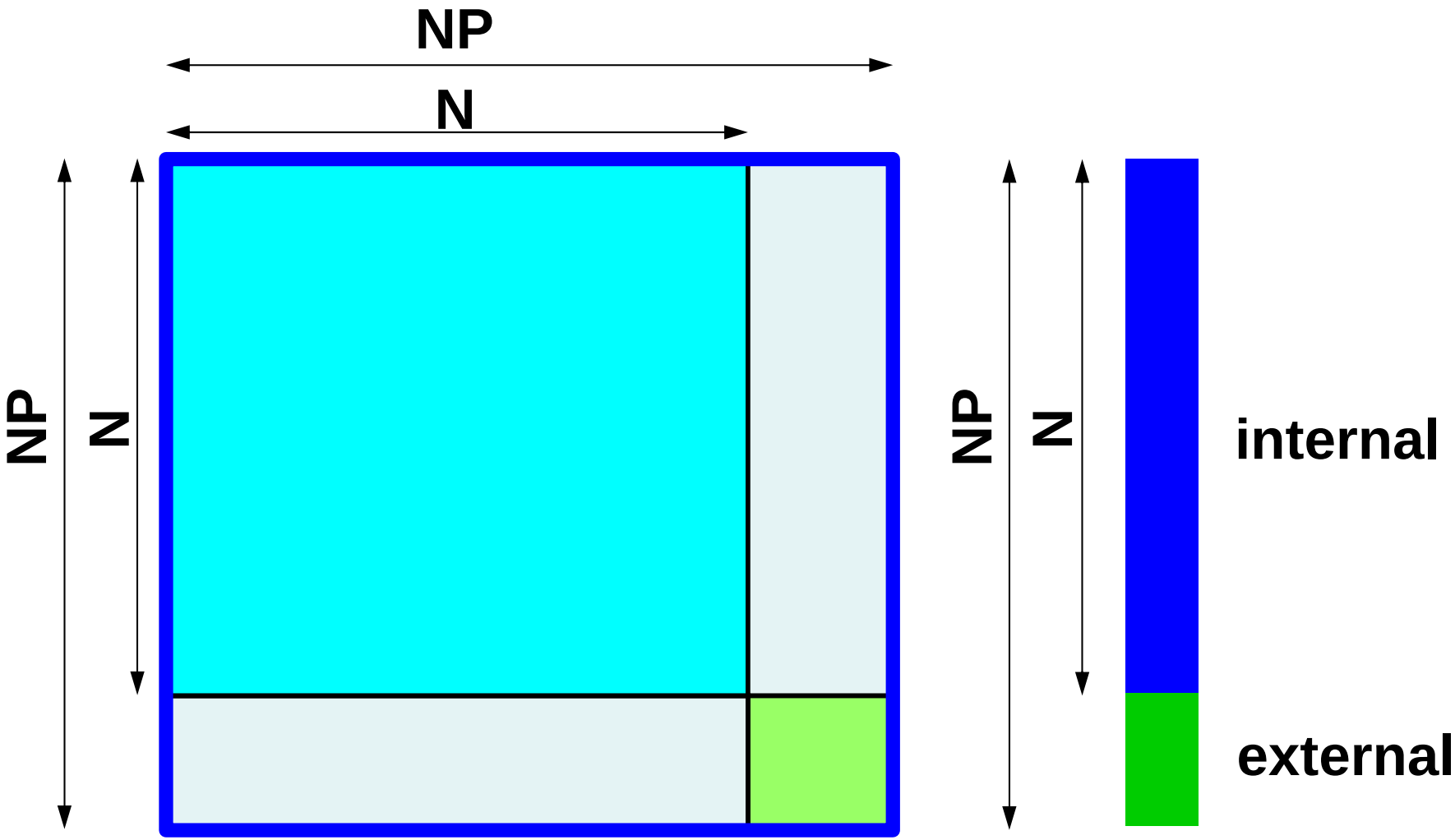
MAT_ASS_MAIN visits all elements

including overlapped elements with external nodes

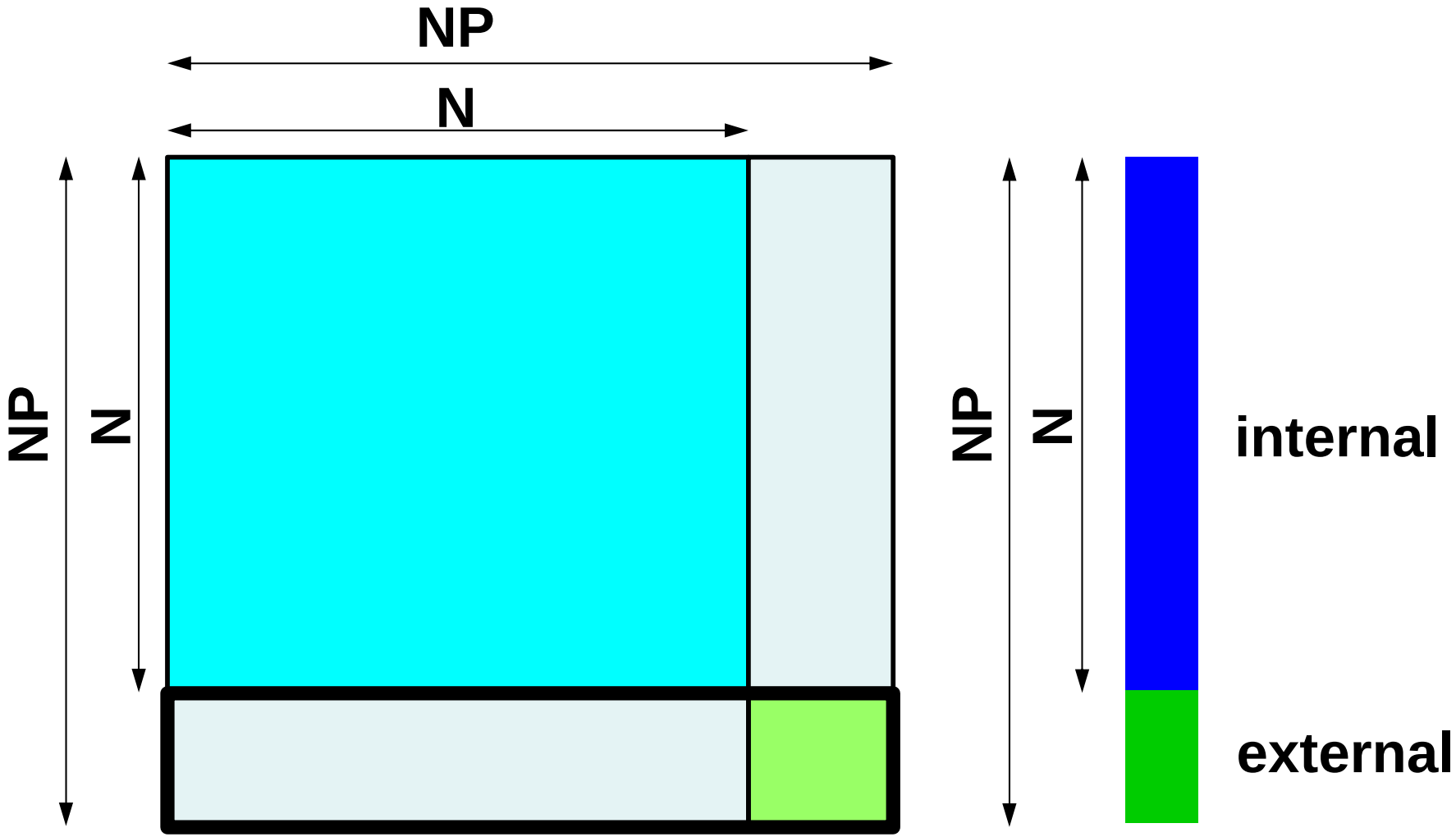




Therefore, we have this matrix



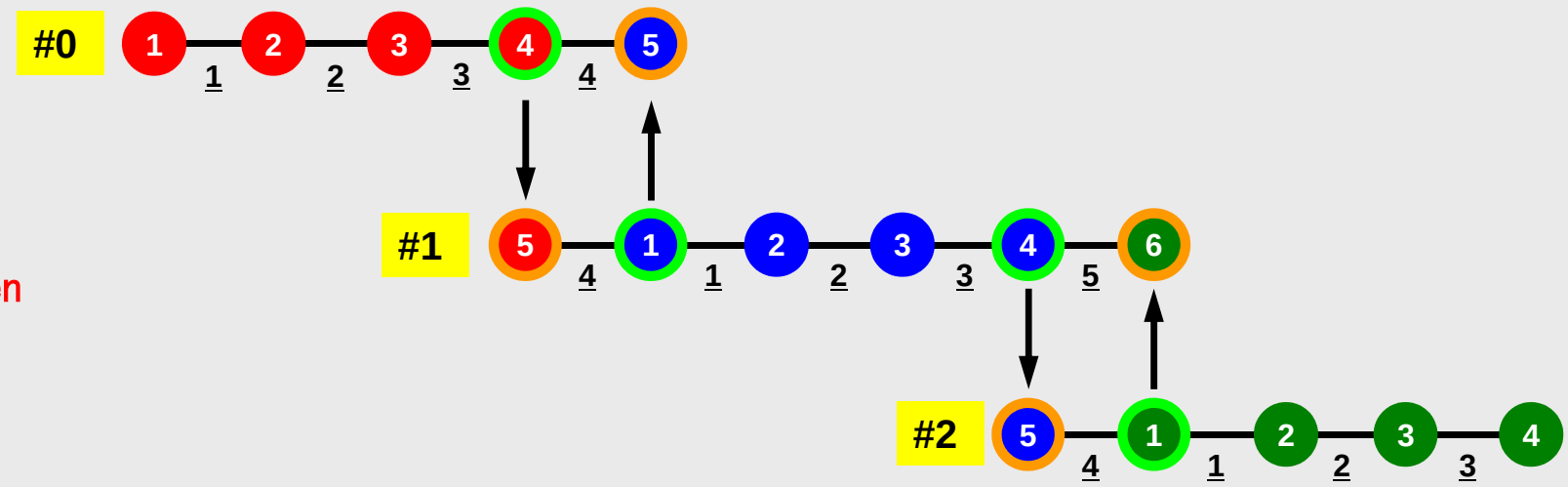
But components of this part are not complete, and not used in computation



Program: 1d.f (9/11)

Boundary Cond., ALMOST NO changes from 1-CPU code

```
!C
!C +--------------------+
!C | BOUNDARY CONDITIONS |
!C +--------------------+
!C===

!C
!C-- X=Xmin

      if (my_rank.eq.0) then
        i = 1
        jS= INDEX(i-1)

        AMAT(jS+1)= 0.d0
        DIAG(i)= 1.d0
        RHS (i)= 0.d0
        do k= 1, NPLU
          if (ITEM(k).eq.1) AMAT(k)= 0.d0
        enddo
      endif

!C===
```



Program: 1d.c(10/11)

Conjugate Gradient Method

```
!C
!C +--------------+
!C | CG iterations |
!C +--------------+
!C===
      R = 1
      Z = 2
      Q = 2
      P = 3
      DD= 4

      do i= 1, N
        W(i,DD)= 1.0D0 / DIAG(i)
      enddo

!C
!C-- {r0}= {b} - [A]{xini} |

!C-    init

        do neib= 1, NEIBPETOT
          do k= export_index(neib-1)+1, export_index(neib)
            kk= export_item(k)
            SENDbuf(k)= PHI(kk)
          enddo
        enddo
```

Compute $r^{(0)}= b-[A]x^{(0)}$

for i= 1, 2, …

- solve $[M]z^{(i-1)}= r^{(i-1)}$
- $\rho_{i-1}= r^{(i-1)}\ z^{(i-1)}$
- if i=1
  - $p^{(1)}= z^{(0)}$
- else
  - $\beta_{i-1}= \rho_{i-1}/\rho_{i-2}$
  - $p^{(i)}= z^{(i-1)} + \beta_{i-1}\ p^{(i-1)}$
- endif
- $q^{(i)}= [A]p^{(i)}$
- $\alpha_i = \rho_{i-1}/p^{(i)}q^{(i)}$
- $x^{(i)}= x^{(i-1)} + \alpha_i p^{(i)}$
- $r^{(i)}= r^{(i-1)} - \alpha_i q^{(i)}$
- check convergence $|r|$

end

Conjugate Gradient Method (CG)

- Matrix-Vector Multiply
- Dot Product
- Preconditioning: in the same way as 1CPU code
- DAXPY: in the same way as 1CPU code

Preconditioning, DAXPY

```
!C
!C-- {z}= [Minv]{r}

      do i= 1, N
        W(i,Z)= W(i,DD) * W(i,R)
      enddo
```

```
!C
!C-- {x}= {x} + ALPHA*{p}
!C   {r}= {r} - ALPHA*{q}

      do i= 1, N
        PHI(i)= PHI(i) + ALPHA * W(i,P)
        W(i,R)= W(i,R) - ALPHA * W(i,Q)
      enddo
```

Matrix-Vector Multiply (1/2)

Using Comm. Table, {p} is updated before computation

```
!C
!C-- {q}= [A]{p}

      do neib= 1, NEIBPETOT
        do k= export_index(neib-1)+1, export_index(neib)
          kk= export_item(k)
          SENDbuf(k)= W(kk,P)
        enddo
      enddo

      do neib= 1, NEIBPETOT
        is   = export_index(neib-1) + 1
        len_s= export_index(neib) - export_index(neib-1)
        call MPI_Isend (SENDbuf(is), len_s, MPI_DOUBLE_PRECISION,        &
     &                  NEIBPE(neib), 0, MPI_COMM_WORLD, request_send(neib), ierr)
      enddo

      do neib= 1, NEIBPETOT
        ir   = import_index(neib-1) + 1
        len_r= import_index(neib) - import_index(neib-1)
        call MPI_Irecv (RECVbuf(ir), len_r, MPI_DOUBLE_PRECISION,        &
     &                  NEIBPE(neib), 0, MPI_COMM_WORLD, request_recv(neib), ierr)
      enddo

      call MPI_Waitall (NEIBPETOT, request_recv, stat_recv, ierr)

      do neib= 1, NEIBPETOT
        do k= import_index(neib-1)+1, import_index(neib)
          kk= import_item(k)
          W(kk,P)= RECVbuf(k)
        enddo
      enddo
```

Matrix-Vector Multiply (2/2)

{q}= [A]{p}

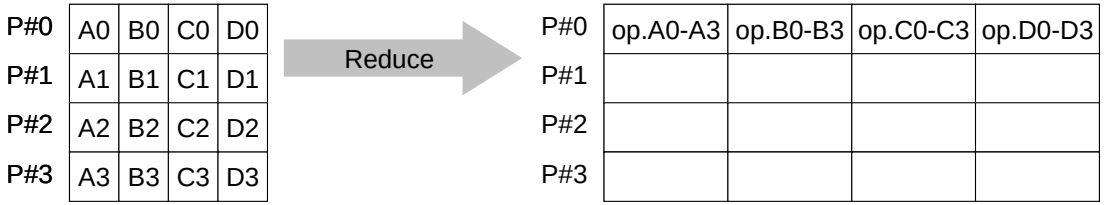
```
call MPI_Waitall (NEIBPETOT, request_send, stat_send, ierr)

do i= 1, N
  W(i,Q) = DIAG(i)*W(i,P)
  do j= INDEX(i-1)+1, INDEX(i)
    W(i,Q) = W(i,Q) + AMAT(j)*W(ITEM(j),P)
  enddo
enddo
```

Dot Product

Global Summation by MPI_Allreduce

```
!C
!C-- RHO= {r}{z}

      RHO0= 0.d0
      do i= 1, N
        RHO0= RHO0 + W(i,R)*W(i,Z)
      enddo
      call MPI_Allreduce (RHO0, RHO, 1, MPI_DOUBLE_PRECISION,          &
     &                    MPI_SUM, MPI_COMM_WORLD, ierr)
```

MPI_REDUCE



- Reduces values on all processes to a single value
  - Summation, Product, Max, Min etc.

- **`call MPI_REDUCE`**
  `(sendbuf,recvbuf,count,datatype,op,root,comm,ierr)`
  - **`sendbuf`** choice I starting address of send buffer
  - **`recvbuf`** choice O starting address receive buffer
    type is defined by "**`datatype`**"
  - **`count`** I I number of elements in send/receive buffer
  - **`datatype`** I I data type of elements of send/recive buffer
    FORTRAN MPI_INTEGER, MPI_REAL, MPI_DOUBLE_PRECISION, MPI_CHARACTER etc.
    C MPI_INT, MPI_FLOAT, MPI_DOUBLE, MPI_CHAR etc
  - **`op`** I I reduce operation
    MPI_MAX, MPI_MIN, MPI_SUM, MPI_PROD, MPI_LAND, MPI_BAND etc
    Users can define operations by **`MPI_OP_CREATE`**
  - **`root`** I I rank of root process
  - **`comm`** I I communicator
  - **`ierr`** I O completion code

Send/Receive Buffer (Sending/Receiving)

- Arrays of “send (sending) buffer” and “receive (receiving) buffer” often appear in MPI.

- Addresses of “send (sending) buffer” and “receive (receiving) buffer” must be different.

Example of MPI_Reduce (1/2)

```
call MPI_REDUCE
(sendbuf,recvbuf,count,datatype,op,root,comm,ierr)
```

```
real(kind=8):: X0, X1

call MPI_REDUCE
(X0, X1, 1, MPI_DOUBLE_PRECISION, MPI_MAX, 0, <comm>, ierr)
```

```
real(kind=8):: X0(4), XMAX(4)

call MPI_REDUCE
(X0, XMAX, 4, MPI_DOUBLE_PRECISION, MPI_MAX, 0, <comm>, ierr)
```

Global Max values of X0(i) go to XMAX(i) on #0 process（i=1~4）

Example of MPI_Reduce (2/2)

```
call MPI_REDUCE
(sendbuf,recvbuf,count,datatype,op,root,comm,ierr)
```

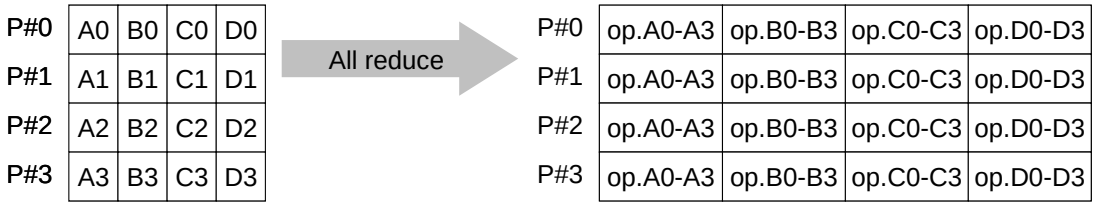
```
real(kind=8):: X0, XSUM

call MPI_REDUCE
(X0, XSUM, 1, MPI_DOUBLE_PRECISION, MPI_SUM, 0, <comm>, ierr)
```

Global summation of X0 goes to XSUM on #0 process.

```
real(kind=8):: X0(4)

call MPI_REDUCE
(X0(1), X0(3), 2, MPI_DOUBLE_PRECISION, MPI_SUM, 0, <comm>, ierr)
```

- Global summation of X0(1) goes to X0(3) on #0 process.
- Global summation of X0(2) goes to X0(4) on #0 process.

MPI_ALLREDUCE



- MPI_Reduce + MPI_Bcast
- Summation (of dot products) and MAX/MIN values are likely to utilized in each process

- **call MPI_ALLREDUCE**
  (sendbuf,recvbuf,count,datatype,op, comm,ierr)
  - **sendbuf** choice I starting address of send buffer
  - **recvbuf** choice O starting address receive buffer
    type is defined by "**datatype**"
  - **count** I I number of elements in send/recv buffer
  - **datatype** I I data type of elements in send/recv buffer
  - **op** I I reduce operation
  - **comm** I I commuinicator
  - **ierr** I O completion code

CG method (1/5)

```
!C
!C-- {r0}= {b} - [A]{xini}
    do neib= 1, NEIBPETOT
      do k= export_index(neib-1)+1, export_index(neib)
        kk= export_item(k)
        SENDbuf(k)= PHI(kk)
      enddo
    enddo

    do neib= 1, NEIBPETOT
      is   = export_index(neib-1) + 1
      len_s= export_index(neib) - export_index(neib-1)
      call MPI_Isend (SENDbuf(is), len_s,
     &                MPI_DOUBLE_PRECISION,
     &                NEIBPE(neib), 0, MPI_COMM_WORLD,
     &                request_send(neib), ierr)
    enddo

    do neib= 1, NEIBPETOT
      ir   = import_index(neib-1) + 1
      len_r= import_index(neib) - import_index(neib-1)
      call MPI_Irecv (RECVbuf(ir), len_r,
     &                MPI_DOUBLE_PRECISION,
     &                NEIBPE(neib), 0, MPI_COMM_WORLD,
     &                request_recv(neib), ierr)
    enddo
    call MPI_Waitall (NEIBPETOT, request_recv, stat_recv, ier

    do neib= 1, NEIBPETOT
      do k= import_index(neib-1)+1, import_index(neib)
        kk= import_item(k)
        PHI(kk)= RECVbuf(k)
      enddo
    enddo
    call MPI_Waitall (NEIBPETOT, request_send, stat_send, ierr)
```

Compute $r^{(0)}= b-[A]x^{(0)}$

for i= 1, 2, …

- solve $[M]z^{(i-1)}= r^{(i-1)}$
- $\rho_{i-1}= r^{(i-1)}\ z^{(i-1)}$
- if i=1
  - $p^{(1)}= z^{(0)}$
- else
  - $\beta_{i-1}= \rho_{i-1}/\rho_{i-2}$
  - $p^{(i)}= z^{(i-1)} + \beta_{i-1}\ p^{(i-1)}$
- endif
- $q^{(i)}= [A]p^{(i)}$
- $\alpha_i = \rho_{i-1}/p^{(i)}q^{(i)}$
- $x^{(i)}= x^{(i-1)} + \alpha_i p^{(i)}$
- $r^{(i)}= r^{(i-1)} - \alpha_i q^{(i)}$
- check convergence $|r|$

end

CG method (2/5)

```
      do i= 1, N
        W(i,R) = DIAG(i)*PHI(i)
        do j= INDEX(i-1)+1, INDEX(i)
          W(i,R) = W(i,R) + AMAT(j)*PHI(ITEM(j))
        enddo
      enddo

      BNRM20= 0.0D0
      do i= 1, N
        BNRM20 = BNRM20  + RHS(i)  **2
        W(i,R) = RHS(i) - W(i,R)
      enddo
      call MPI_Allreduce (BNRM20, BNRM2, 1,
     &                    MPI_DOUBLE_PRECISION,
     &                    MPI_SUM, MPI_COMM_WORLD, ierr)

!C***********************************************************
      do iter= 1, ITERmax

!C
!C-- {z}= [Minv]{r}

      do i= 1, N
        W(i,Z)= W(i,DD) * W(i,R)
      enddo
!C
!C-- RHO= {r}{z}

      RHO0= 0.d0
      do i= 1, N
        RHO0= RHO0 + W(i,R)*W(i,Z)
      enddo
      call MPI_Allreduce (RHO0, RHO, 1, MPI_DOUBLE_PRECISION,
     &                    MPI_SUM, MPI_COMM_WORLD, ierr)
```

Compute $r^{(0)}= b-[A]x^{(0)}$

for i= 1, 2, …

- **solve $[M]z^{(i-1)}= r^{(i-1)}$**
- $\rho_{i-1}= r^{(i-1)}\ z^{(i-1)}$
- if i=1
  - $p^{(1)}= z^{(0)}$
- else
  - $\beta_{i-1}= \rho_{i-1}/\rho_{i-2}$
  - $p^{(i)}= z^{(i-1)} + \beta_{i-1}\ p^{(i-1)}$
- endif
- $q^{(i)}= [A]p^{(i)}$
- $\alpha_i = \rho_{i-1}/p^{(i)}q^{(i)}$
- $x^{(i)}= x^{(i-1)} + \alpha_i p^{(i)}$
- $r^{(i)}= r^{(i-1)} - \alpha_i q^{(i)}$
- check convergence $|r|$

end

CG method (3/5)

```
!C
!C-- {p} = {z} if      ITER=1
!C   BETA= RHO / RHO1  otherwise

      if ( iter.eq.1 ) then
        do i= 1, N
          W(i,P)= W(i,Z)
        enddo
       else
        BETA= RHO / RHO1
        do i= 1, N
          W(i,P)= W(i,Z) + BETA*W(i,P)
        enddo
      endif


!C
!C-- {q}= [A]{p}

      do neib= 1, NEIBPETOT
        do k= export_index(neib-1)+1, export_index(neib)
          kk= export_item(k)
          SENDbuf(k)= W(kk,P)
        enddo
      enddo

      do neib= 1, NEIBPETOT
        is   = export_index(neib-1) + 1
        len_s= export_index(neib) - export_index(neib-1)
        call MPI_Isend (SENDbuf(is), len_s, MPI_DOUBLE_PRECISION,       &
     &                  NEIBPE(neib), 0, MPI_COMM_WORLD,                &
     &                  request_send(neib), ierr)
      enddo
```

Compute $r^{(0)}= b-[A]x^{(0)}$

for i= 1, 2, …

- solve $[M]z^{(i-1)}= r^{(i-1)}$
- $\rho_{i-1}= r^{(i-1)}\ z^{(i-1)}$
- **if i=1**
  - **$p^{(1)}= z^{(0)}$**
  - **else**
    - **$\beta_{i-1}= \rho_{i-1}/\rho_{i-2}$**
    - **$p^{(i)}= z^{(i-1)} + \beta_{i-1}\ p^{(i-1)}$**
  - **endif**
- **$q^{(i)}= [A]p^{(i)}$**
- $\alpha_i = \rho_{i-1}/p^{(i)}q^{(i)}$
- $x^{(i)}= x^{(i-1)} + \alpha_i p^{(i)}$
- $r^{(i)}= r^{(i-1)} - \alpha_i q^{(i)}$
- check convergence $|r|$

end

CG method (4/5)

```
  do neib= 1, NEIBPETOT
    ir   = import_index(neib-1) + 1
    len_r= import_index(neib) - import_index(neib-1)
    call MPI_Irecv (RECVbuf(ir), len_r,
&                   MPI_DOUBLE_PRECISION,
&                   NEIBPE(neib), 0, MPI_COMM_WORLD,
&                   request_recv(neib), ierr)
  enddo
  call MPI_Waitall (NEIBPETOT, request_recv, stat_recv, ierr)

  do neib= 1, NEIBPETOT
    do k= import_index(neib-1)+1, import_index(neib)
      kk= import_item(k)
      W(kk,P)= RECVbuf(k)
    enddo
  enddo
  call MPI_Waitall (NEIBPETOT, request_send, stat_send, ierr)

  do i= 1, N
    W(i,Q) = DIAG(i)*W(i,P)
    do j= INDEX(i-1)+1, INDEX(i)
      W(i,Q) = W(i,Q) + AMAT(j)*W(ITEM(j),P)
    enddo
  enddo

!C
!C-- ALPHA= RHO / {p}{q}

  C10= 0.d0
  do i= 1, N
    C10= C10 + W(i,P)*W(i,Q)
  enddo
  call MPI_Allreduce (C10, C1, 1, MPI_DOUBLE_PRECISION, MPI_SUM, MPI_COMM_WORLD, ierr)
  ALPHA= RHO / C1
```

Compute $r^{(0)}= b-[A]x^{(0)}$

for i= 1, 2, …

- solve $[M]z^{(i-1)}= r^{(i-1)}$
- $\rho_{i-1}= r^{(i-1)}\ z^{(i-1)}$
- if i=1
  - $p^{(1)}= z^{(0)}$
- else
  - $\beta_{i-1}= \rho_{i-1}/\rho_{i-2}$
  - $p^{(i)}= z^{(i-1)} + \beta_{i-1}\ p^{(i-1)}$
- endif
- **$q^{(i)}= [A]p^{(i)}$**
- **$\alpha_i = \rho_{i-1}/p^{(i)}q^{(i)}$**
- $x^{(i)}= x^{(i-1)} + \alpha_i p^{(i)}$
- $r^{(i)}= r^{(i-1)} - \alpha_i q^{(i)}$
- check convergence $|r|$

end

CG method (5/5)

```
!C
!C-- {x}= {x} + ALPHA*{p}
!C   {r}= {r} - ALPHA*{q}

      do i= 1, N
        PHI(i)= PHI(i) + ALPHA * W(i,P)
        W(i,R)= W(i,R) - ALPHA * W(i,Q)
      enddo

      DNRM20 = 0.0
      do i= 1, N
        DNRM20= DNRM20 + W(i,R)**2
      enddo

      call MPI_Allreduce (DNRM20, DNRM2, 1,
     &                    MPI_DOUBLE_PRECISION,
     &                    MPI_SUM, MPI_COMM_WORLD, ierr)

        RESID= dsqrt(DNRM2/BNRM2)

        if (my_rank.eq.0.and.mod(iter,1000).eq.0) then
          write (*, '(i8,1pe16.6)') iter, RESID
        endif

        if ( RESID.le.EPS) goto 900
        RHO1 = RHO

      enddo
```

Compute $r^{(0)}= b-[A]x^{(0)}$

for i= 1, 2, …

- solve $[M]z^{(i-1)}= r^{(i-1)}$
- $\rho_{i-1}= r^{(i-1)}\ z^{(i-1)}$
- if i=1
  - $p^{(1)}= z^{(0)}$
- else
  - $\beta_{i-1}= \rho_{i-1}/\rho_{i-2}$
  - $p^{(i)}= z^{(i-1)} + \beta_{i-1}\ p^{(i-1)}$
- endif
- $q^{(i)}= [A]p^{(i)}$
- $\alpha_i = \rho_{i-1}/p^{(i)}q^{(i)}$
- **$x^{(i)}= x^{(i-1)} + \alpha_i p^{(i)}$**
- **$r^{(i)}= r^{(i-1)} - \alpha_i q^{(i)}$**
- **check convergence $|r|$**

end

Program: 1d.f (11/11)

Output by Each Proces

```
!C
!C-- OUTPUT
        if (my_rank.eq.0) then
          write (*,'(2(1pe16.6))') E1Time-S1Time, E2Time-E1Time
        endif

      write (*,'(/a)') '### TEMPERATURE'
      do i= 1, N
        write (*,'(2i8, 2(1pe16.6))') my_rank, i, PHI(i)
      enddo

      call MPI_FINALIZE (ierr)
      end program heat1Dp
```

- Overview
- Distributed Local Data
- Program
- **Results**

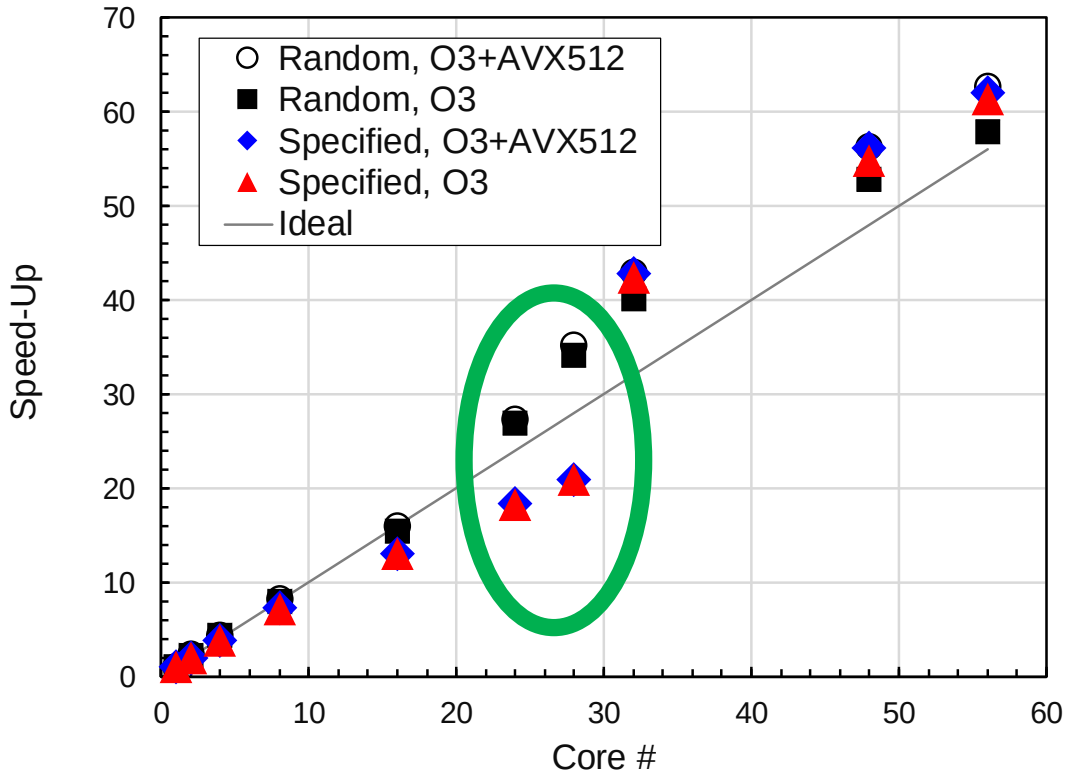
Results: Time for CG Solver, N=10^6

Time for 1,000 iterations, Strong Scaling

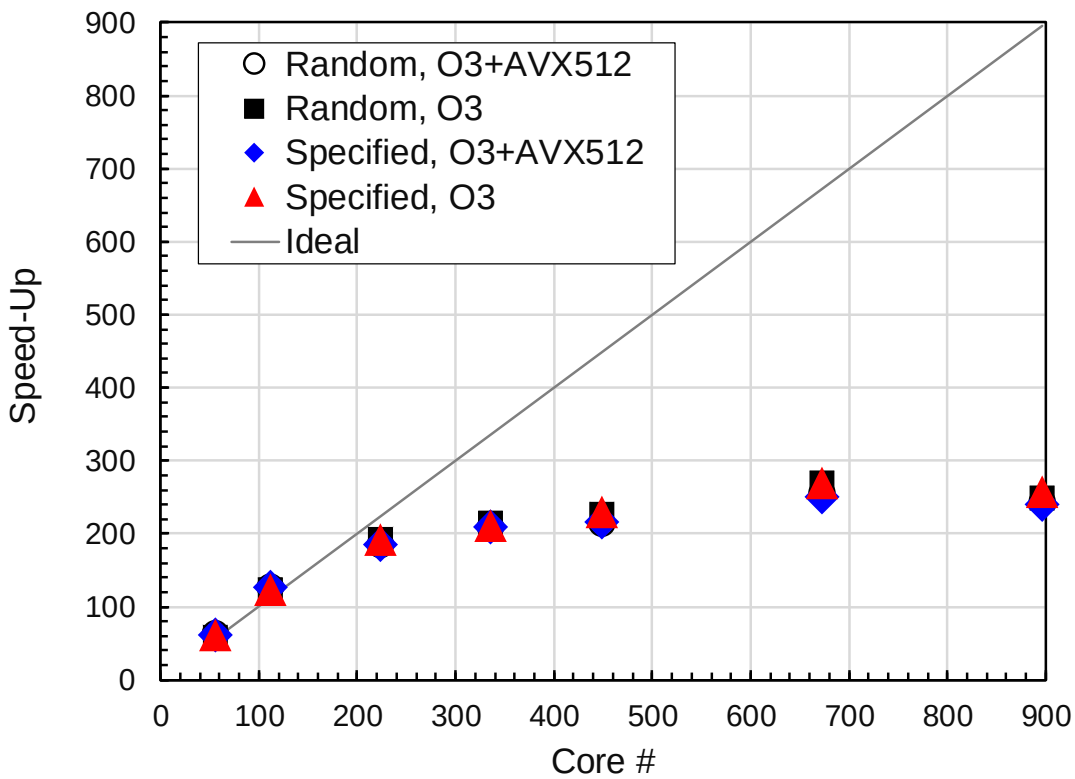
All 56 cores are used if number of nodes is more than 1

Perf. at a single core= 1.00, 5 measurements, Best Case

up to 1 node, 56cores



up to 16 nodes, 896 cores



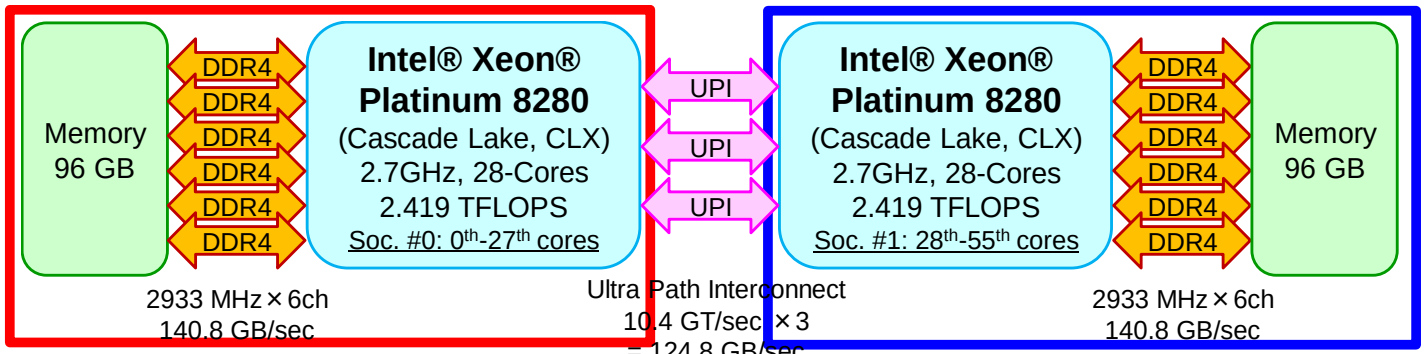
1-node, 24-cores

```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture2
#PJM -L node=1
#PJM --mpi proc=24
#PJM -L elapse=00:15:00
#PJM -g gt62
#PJM -j
#PJM -e err
#PJM -o test.lst

mpiexec.hydra -n ${PJM_MPI_PROC} ./1da
mpiexec.hydra -n ${PJM_MPI_PROC} ./1db
export I_MPI_PIN_PROCESSOR_LIST=0-23
mpiexec.hydra -n ${PJM_MPI_PROC} ./1da
mpiexec.hydra -n ${PJM_MPI_PROC} ./1db
```

1da: -O3 + AVX512
1db: -O3 Only

24-cores are randomly selected from 56-cores on the node
RANDOM

24-cores on Socket #0 are assigned.
SPECIFIED



Socket #0: 0th-27th Cores

Socket #1: 28th-55th Cores

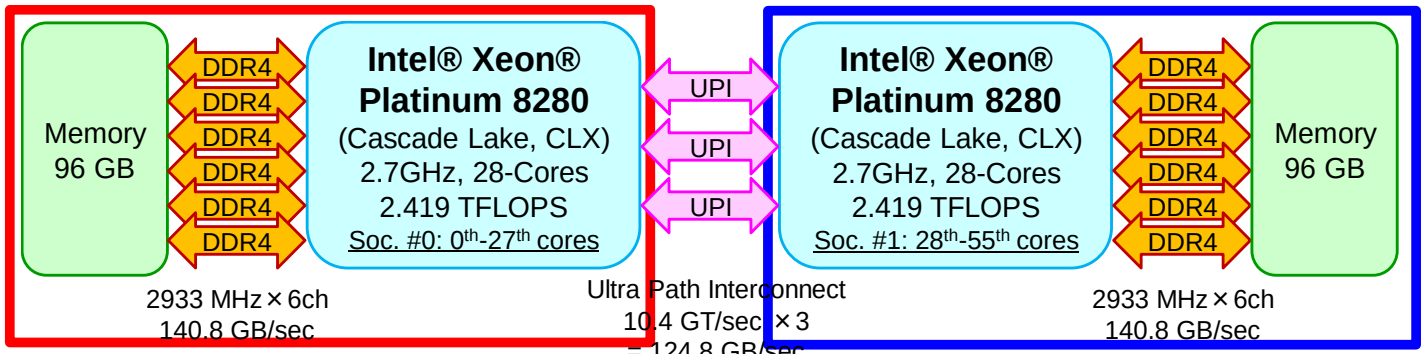
1-node, 28-cores

```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture2
#PJM -L node=1
#PJM --mpi proc=28
#PJM -L elapse=00:15:00
#PJM -g gt62
#PJM -j
#PJM -e err
#PJM -o test.lst

mpiexec.hydra -n ${PJM_MPI_PROC} ./1da
mpiexec.hydra -n ${PJM_MPI_PROC} ./1db
export I_MPI_PIN_PROCESSOR_LIST=0-27
mpiexec.hydra -n ${PJM_MPI_PROC} ./1da
mpiexec.hydra -n ${PJM_MPI_PROC} ./1db
```

1da: -O3 + AVX512
1db: -O3 Only

28-cores are randomly selected from 56-cores on the node
RANDOM

28-cores on Socket #0 are assigned.
SPECIFIED



Socket #0: 0th-27th Cores

Socket #1: 28th-55th Cores

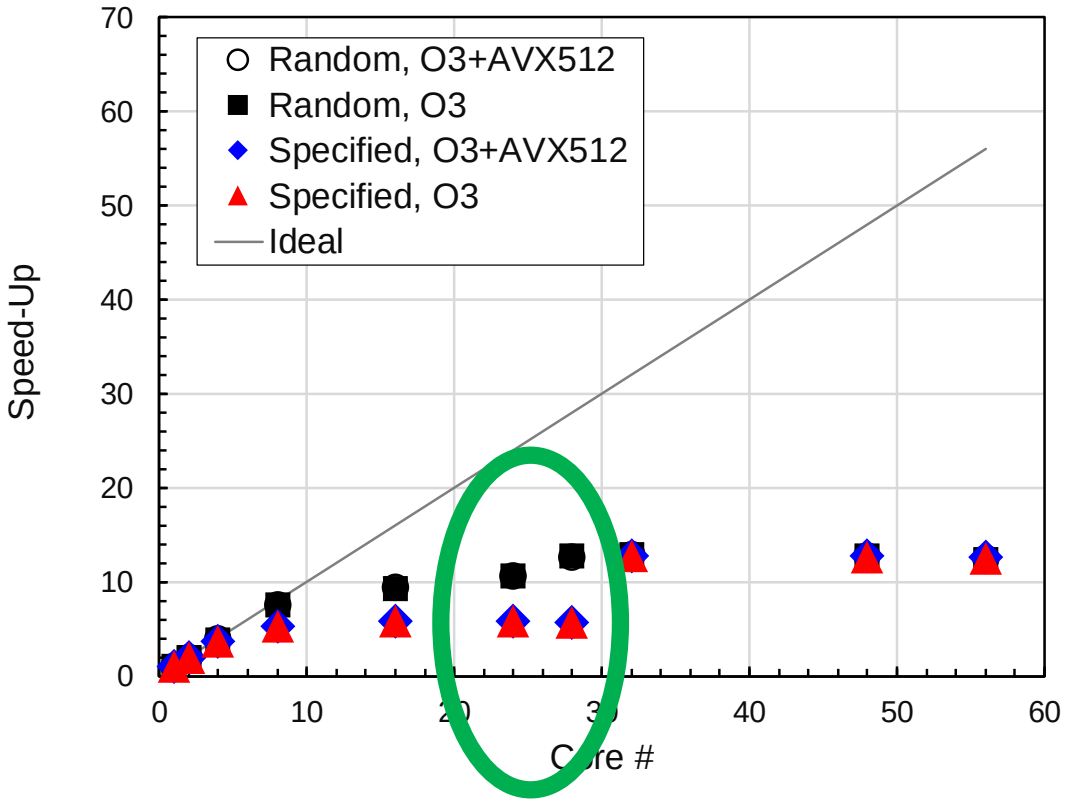
Results: Time for CG Solver, N=10^7

Time for 200 iterations, Strong Scaling

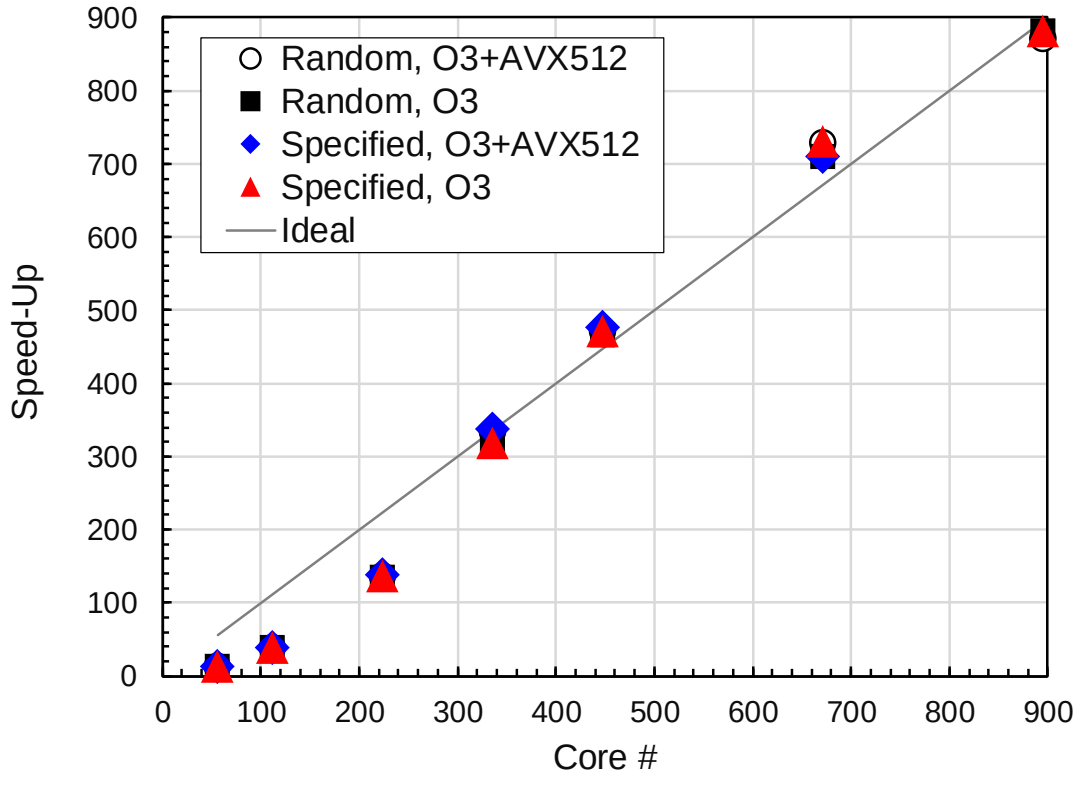
All 56 cores are used if number of nodes is more than 1

Perf. at a single core= 1.00, 5 measurements, Best Case

up to 1 node, 56cores



up to 16 nodes, 896 cores

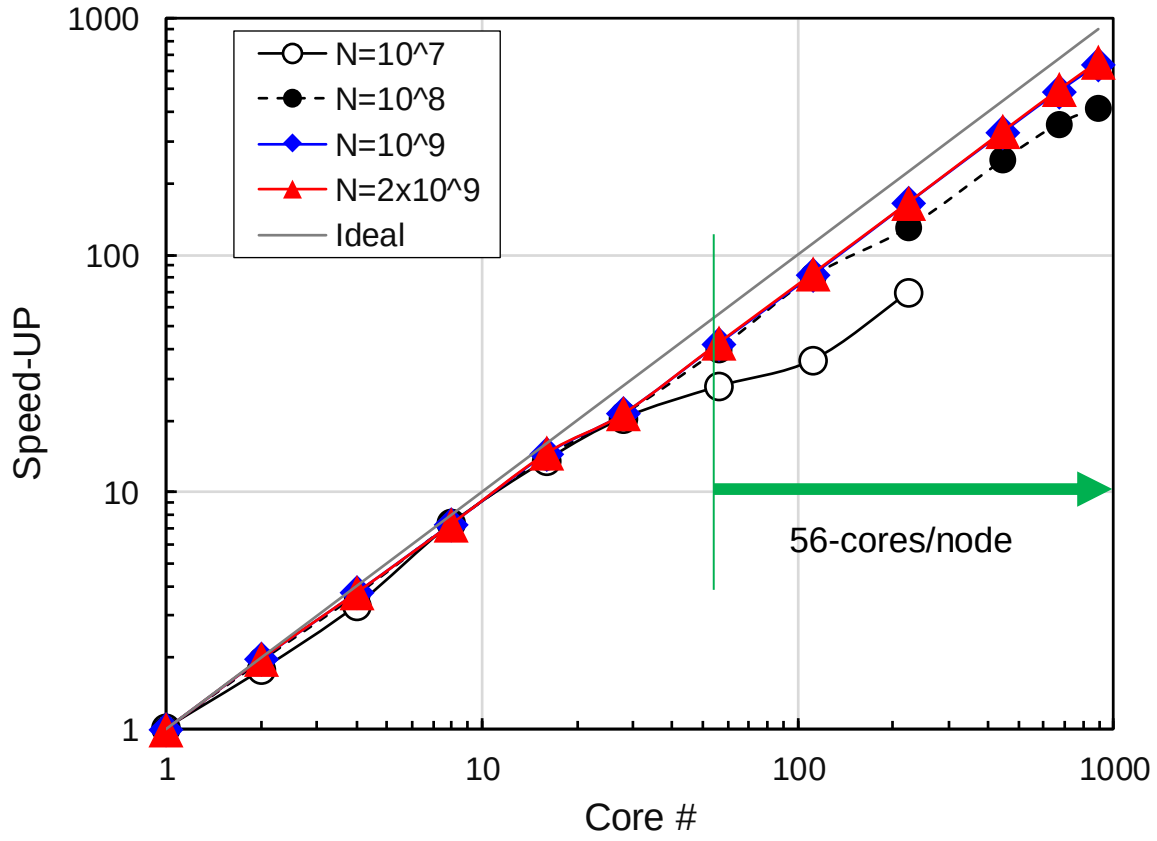


Performance is lower than ideal one

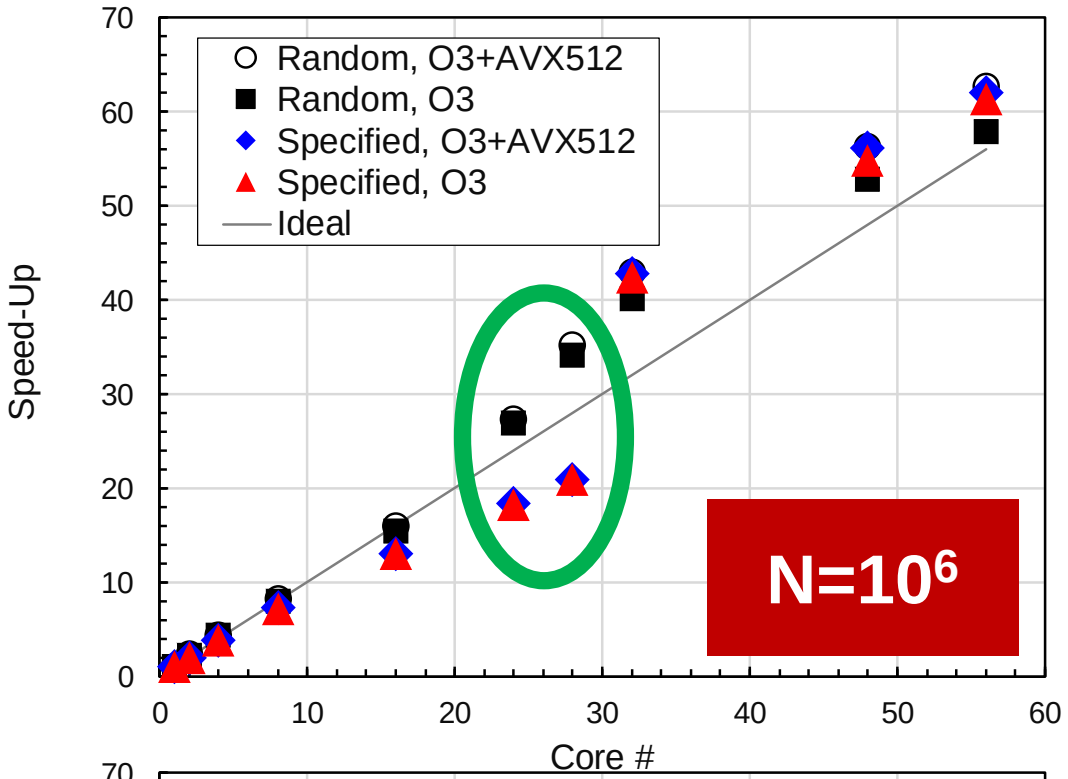
- Time for MPI communication
  - Time for sending data
  - Communication bandwidth between nodes
  - Time is proportional to size of sending/receiving buffers
- Time for starting MPI
  - latency
  - does not depend on size of buffers
    - depends on number of calling, increases according to process #
  - $O(10^0)$-$O(10^1)$ μsec.
- Synchronization of MPI
  - Increases according to number of processes

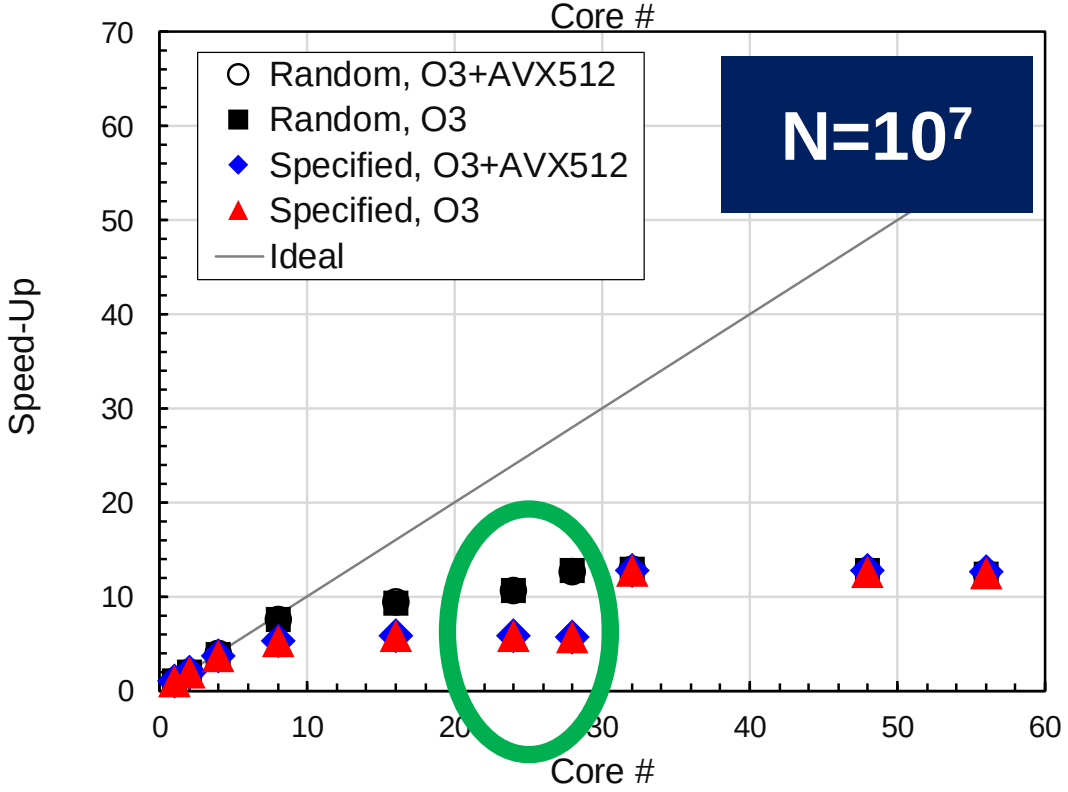
Performance is lower than ideal one (cont.): S1-3

- If computation time is relatively small (N is small in S1-3), these effects are not negligible.
  - If the size of messages is small, effect of “latency” is significant.



Up to 56 cores (single node)

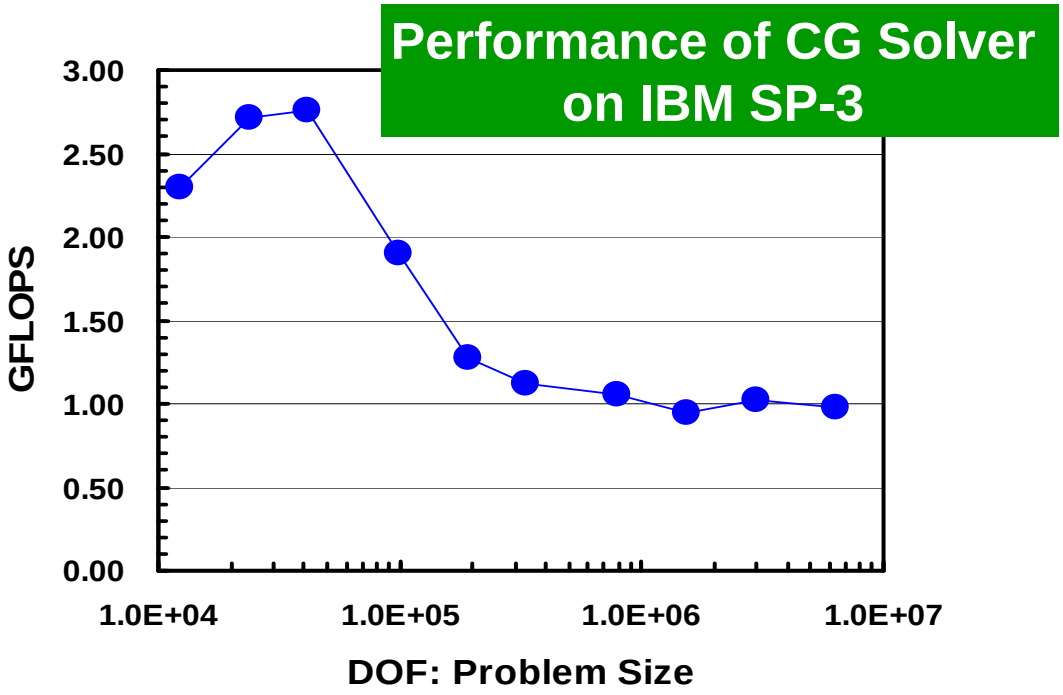


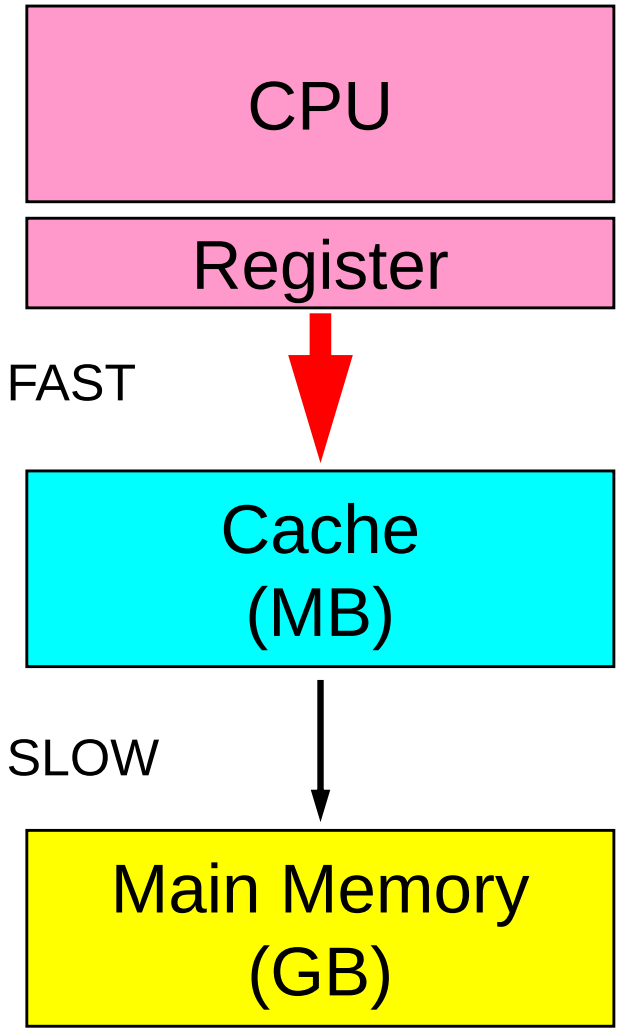


- in a single node (< 56 cores), performance of small cases ($N=10^6$) are rather better.
- Effect of memory contention/saturation (not communication)
- Memory throughput on each soc. is constant for 8+ cores (Stream)
- If problem size is small, cache can be well-utilized. Therefore, effect of memory bandwidth is small.
- Memory/cache of two sockets may be more efficiently utilized in “Random”

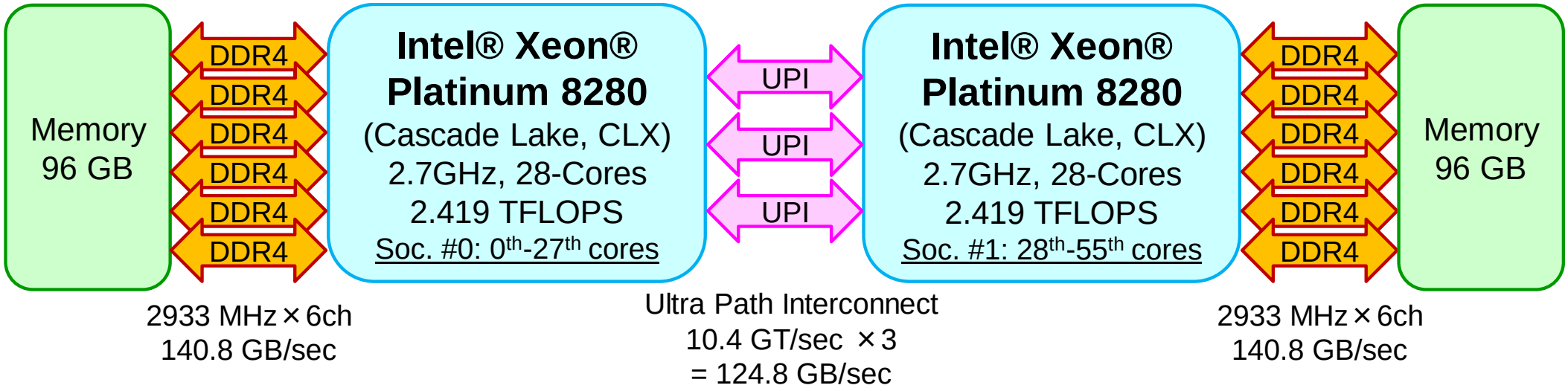
If problem size is small, cache can be well-utilized

- In scalar processors, performance for smaller problem is generally better.
  - Cache is well-utilized.

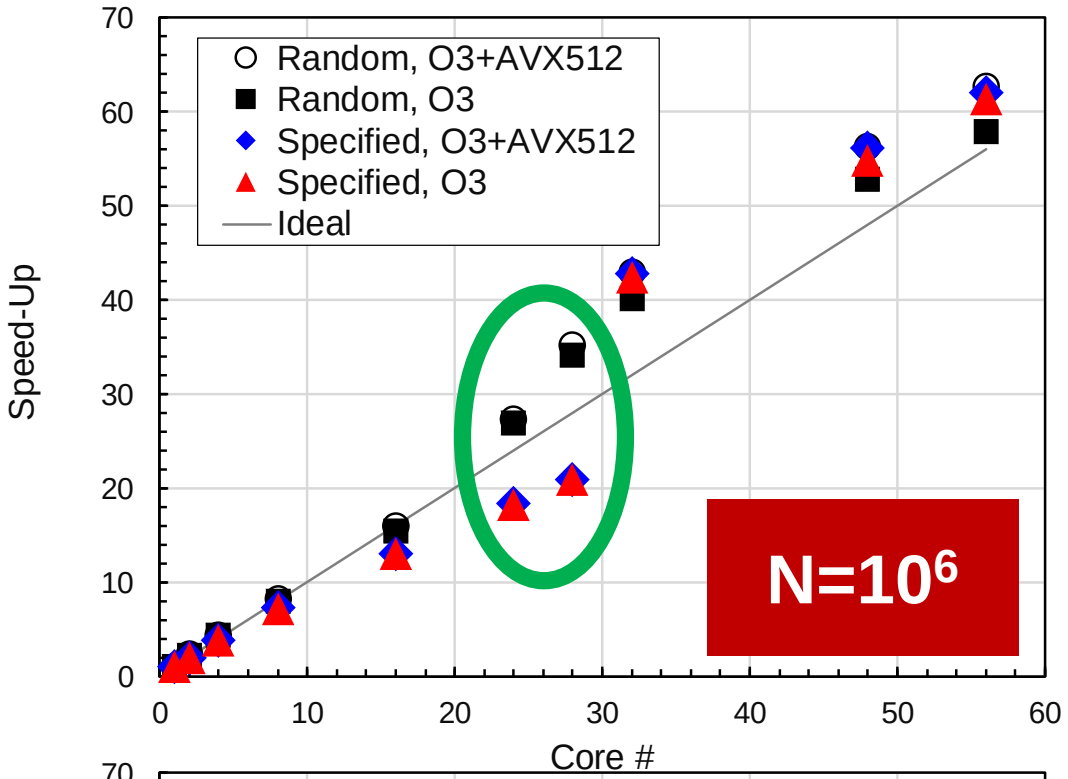


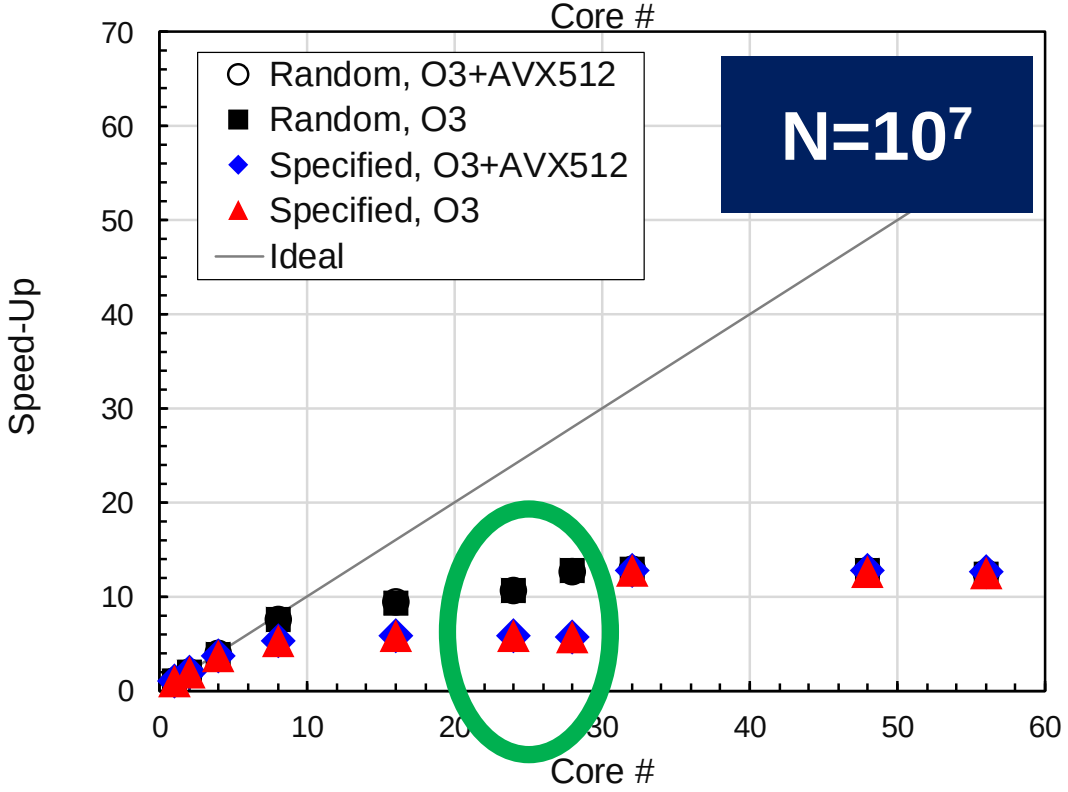


Category	Capacity	X-Way Set Associative	Cache Line
L1$Data	32 KB/core	8-Way	64B
L1$Instruction	32 KB/core	8-Way	64B
L2	1.00 MB/core	16-Way	64B
L3	38.5 MB/socket	11-Way	64B



Up to 56 cores (single node)





- Required Memory
- $N=10^6$: 80MB
  - Very Close to Cache Size
- $N=10^7$: 800MB

- Memory throughput on each socket is constant for 8+ cores (Stream)
- If problem size is small, cache can be well-utilized. Therefore, effect of memory bandwidth is small.

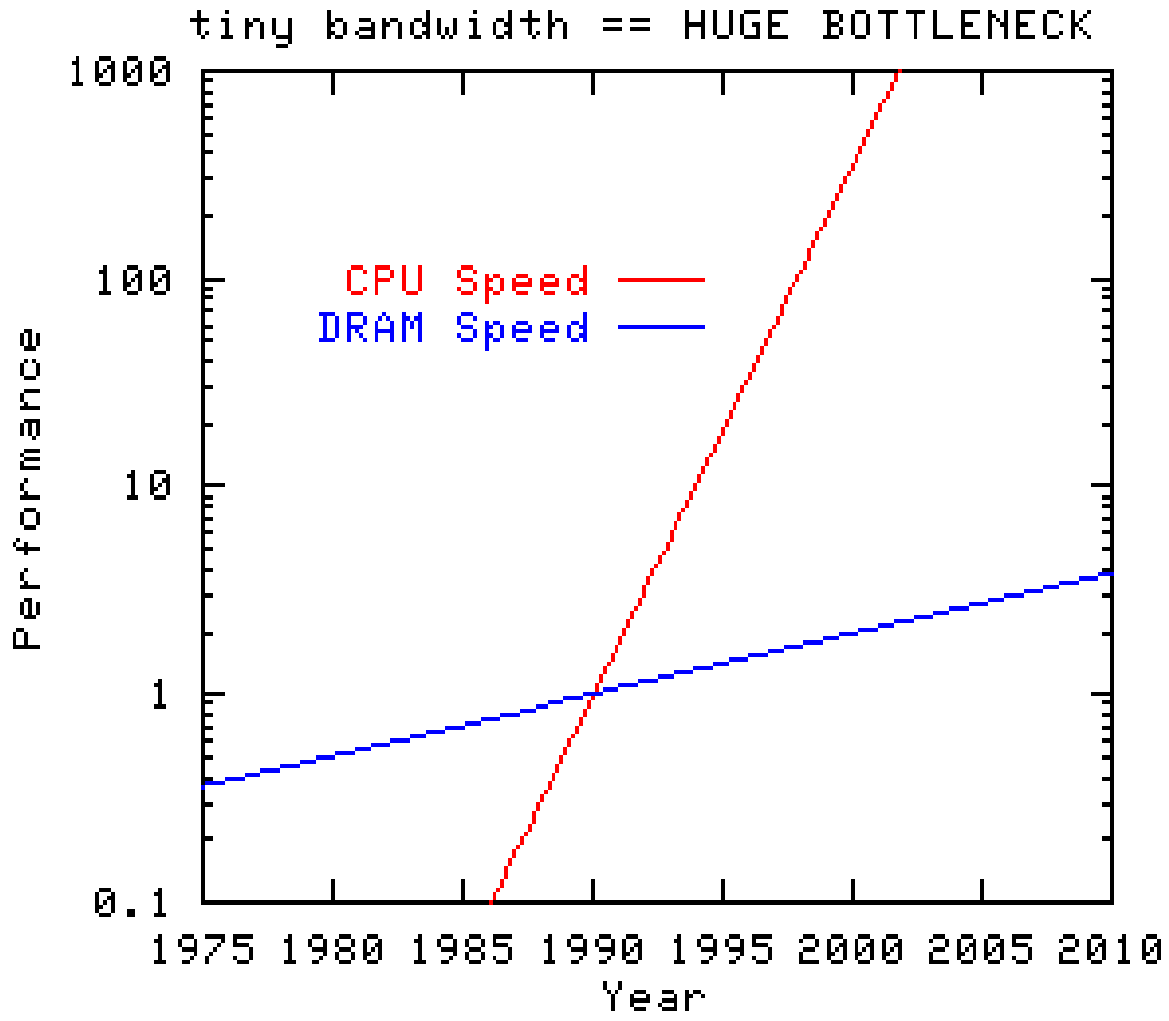
STREAM benchmark

http://www.cs.virginia.edu/stream/

- Benchmarks for Memory Bandwidth
  - Copy: c(i)= a(i)
  - Scale: c(i)= s*b(i)
  - Add: c(i)= a(i) + b(i)
  - Triad: c(i)= a(i) + s*b(i)

```
-------------------------------------------------
 Double precision appears to have 16 digits of accuracy
 Assuming 8 bytes per DOUBLE PRECISION word
-------------------------------------------------
 Number of processors =           16
 Array size =    2000000
 Offset     =          0
 The total memory requirement is     732.4 MB (
45.8MB/task)
 You are running each test  10 times
 --
 The *best* time for each test is used
 *EXCLUDING* the first and last iterations
 -----------------------------------------------------
 -----------------------------------------------------
Function     Rate (MB/s)  Avg time   Min time  Max time
Copy:      18334.1898      0.0280      0.0279      0.0280
Scale:     18035.1690      0.0284      0.0284      0.0285
Add:       18649.4455      0.0412      0.0412      0.0413
Triad:     19603.8455      0.0394      0.0392      0.0398
```

Gap between performance of CPU and Memory



http://www.cs.virginia.edu/stream/

Sparse/Dense Matrices

```
do i= 1, N
  Y(i)= D(i)*X(i)
  do k= index(i-1)+1, index(i)
    Y(i)= Y(i) + AMAT(k)*X(item(k))
  enddo
enddo
```

```
do j= 1, N
  do i= 1, N
    Y(j)= Y(j) + A(i,j)*X(i)
  enddo
enddo
```

- “X” in RHS
  - Dense: continuous on memory, easy to utilize cache
  - Sparse: continuity is not assured (in-direct access), difficult to utilize cache
    - more “memory-bound”

GeoFEM Benchmark

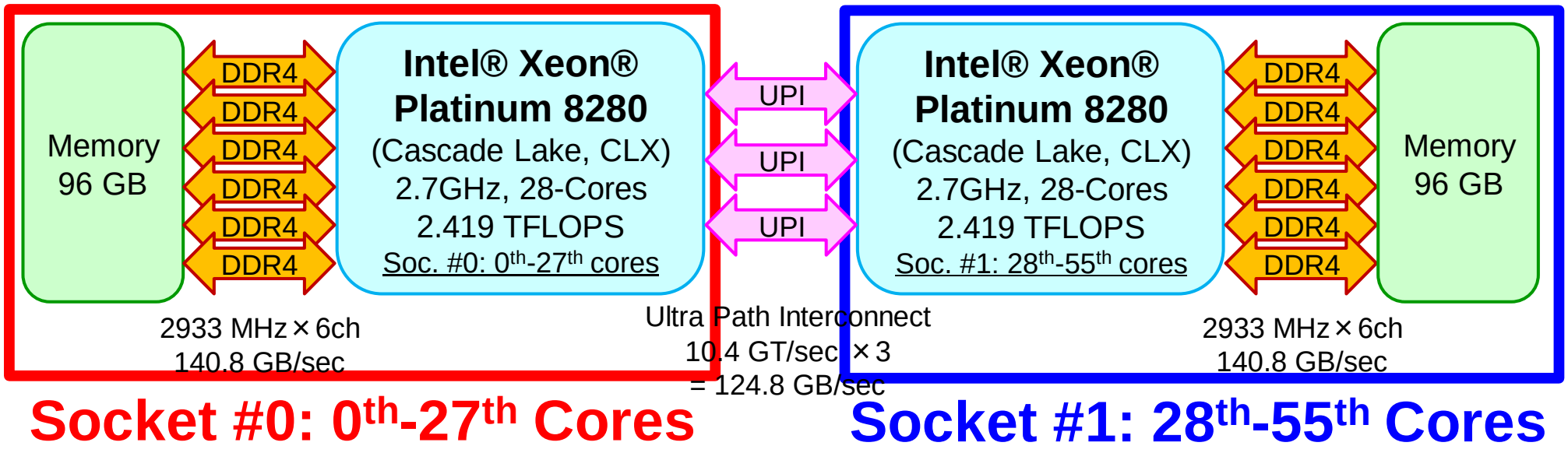
ICCG in FEM for Solid Mechanics

	SR11K/J2	SR16K/M1	T2K	FX10	京
Core #/Node	16	32	16	16	8
Peak Performance (GFLOPS)	147.2	980.5	147.2	236.5	128.0
STREAM Triad (GB/s)	101.0	264.2	20.0	64.7	43.3
B/F	0.686	0.269	0.136	0.274	0.338
GeoFEM (GFLOPS)	19.0	72.7	4.69	16.0	11.0
% to Peak	12.9	7.41	3.18	6.77	8.59
LLC/core (MB)	18.0	4.00	2.00	0.75	0.75

Sparse Linear Solver: Memory-Bound

Copy, Compile and Run

```
>$ cd /work/gt62/t62XXX/pFEM
>$ cp /work/gt62/z30088/pFEM/F/stream.tar .
>$ tar xvf stream.tar
>$ cd mpi/stream

>$ mpiifort -align array64byte -O3 -axCORE-AVX512 stream.f -o stream
>$ pjsub XXX.sh
```

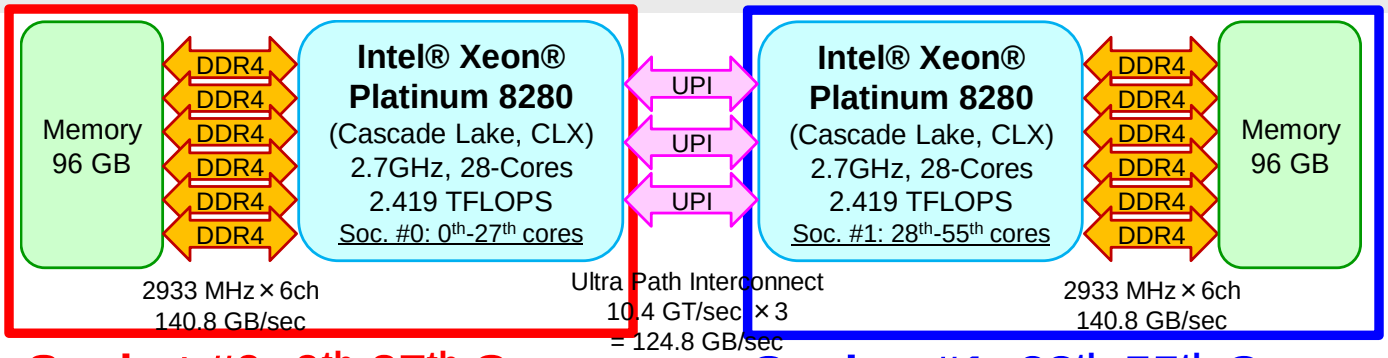


s01.sh: Use 1 core

```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture2
#PJM -L node=1
#PJM --mpi proc=1
#PJM -L elapse=00:15:00
#PJM -g gt62
#PJM -j
#PJM -e err
#PJM -o s01.lst

mpiexec.hydra -n ${PJM_MPI_PROC} numactl -l ./stream
mpiexec.hydra -n ${PJM_MPI_PROC} ./stream

export I_MPI_PIN_PROCESSOR_LIST=0
mpiexec.hydra -n ${PJM_MPI_PROC} numactl -l ./stream
mpiexec.hydra -n ${PJM_MPI_PROC} ./stream
```

Cores are randomly selected

Core are specified

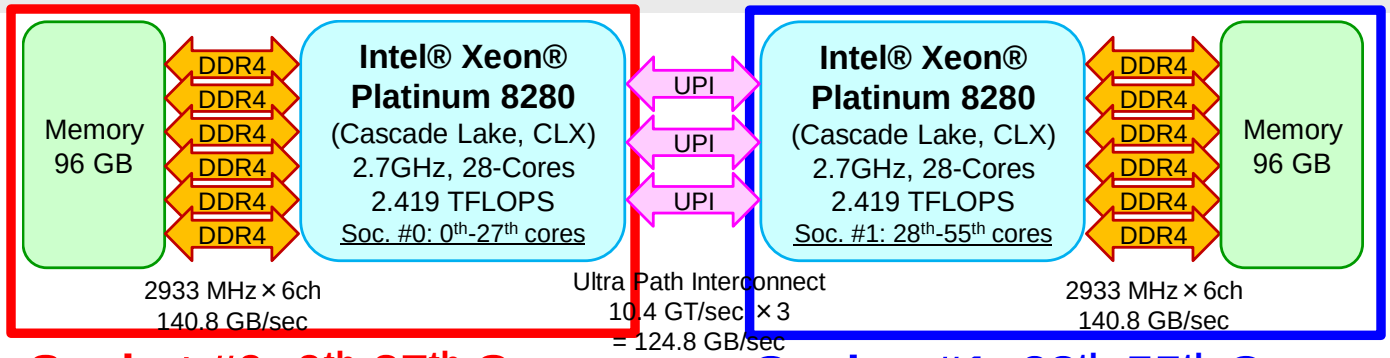


s16.sh: Use 16 cores

```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture2
#PJM -L node=1
#PJM --mpi proc=16
#PJM -L elapse=00:15:00
#PJM -g gt62
#PJM -j
#PJM -e err
#PJM -o s16.lst

mpiexec.hydra -n ${PJM_MPI_PROC} numactl -l ./stream
mpiexec.hydra -n ${PJM_MPI_PROC} ./stream

export I_MPI_PIN_PROCESSOR_LIST=0-15
mpiexec.hydra -n ${PJM_MPI_PROC} numactl -l ./stream
mpiexec.hydra -n ${PJM_MPI_PROC} ./stream
```

Cores are randomly selected

Core are specified

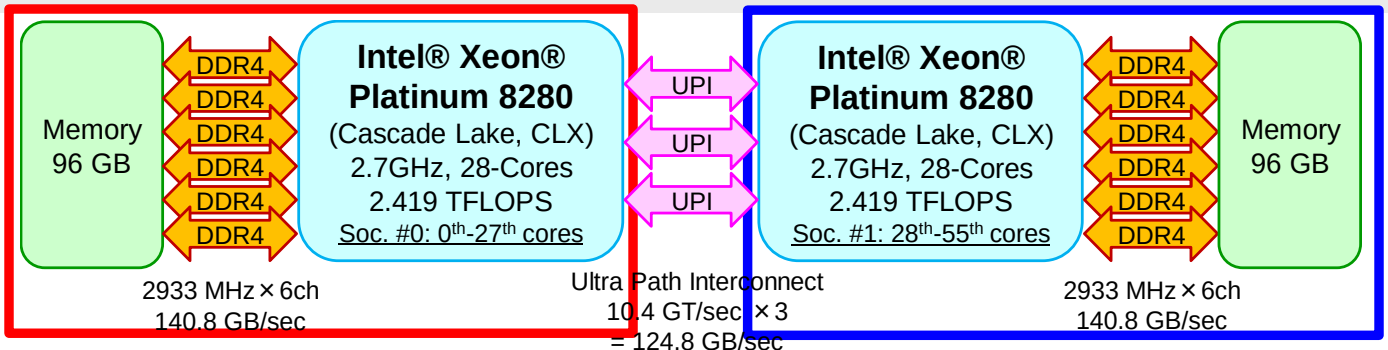


s32.sh: Use 32 cores

```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture2
#PJM -L node=1
#PJM --mpi proc=32
#PJM -L elapse=00:15:00
#PJM -g gt62
#PJM -j
#PJM -e err
#PJM -o s32.lst

mpiexec.hydra -n ${PJM_MPI_PROC} numactl -l ./stream
mpiexec.hydra -n ${PJM_MPI_PROC} ./stream

export I_MPI_PIN_PROCESSOR_LIST=0-15,28-43
mpiexec.hydra -n ${PJM_MPI_PROC} numactl -l ./stream
mpiexec.hydra -n ${PJM_MPI_PROC} ./stream
```

Cores are randomly selected

Core are specified

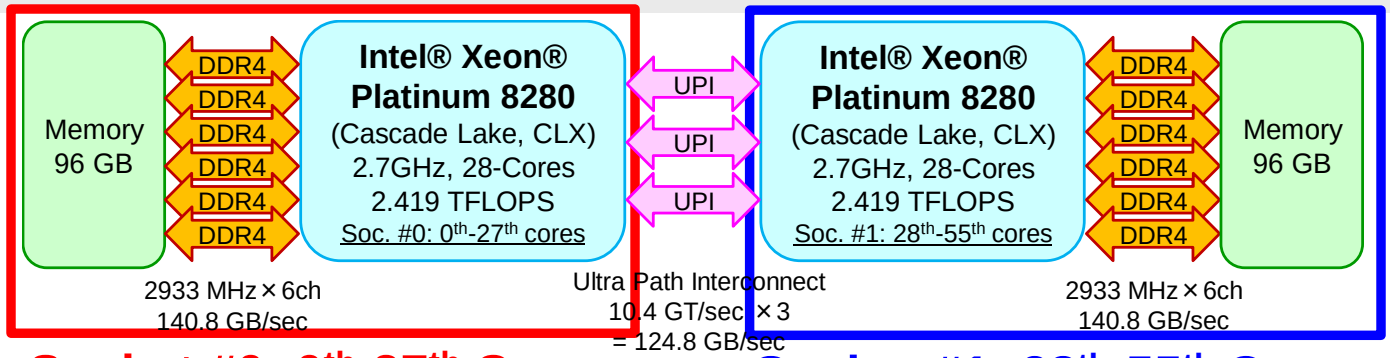


s48.sh: Use 48 cores

```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture2
#PJM -L node=1
#PJM --mpi proc=48
#PJM -L elapse=00:15:00
#PJM -g gt62
#PJM -j
#PJM -e err
#PJM -o s48.lst

mpiexec.hydra -n ${PJM_MPI_PROC} numactl -l ./stream
mpiexec.hydra -n ${PJM_MPI_PROC} ./stream

export I_MPI_PIN_PROCESSOR_LIST=0-23,28-51
mpiexec.hydra -n ${PJM_MPI_PROC} numactl -l ./stream
mpiexec.hydra -n ${PJM_MPI_PROC} ./stream
```

Cores are randomly selected

Core are specified

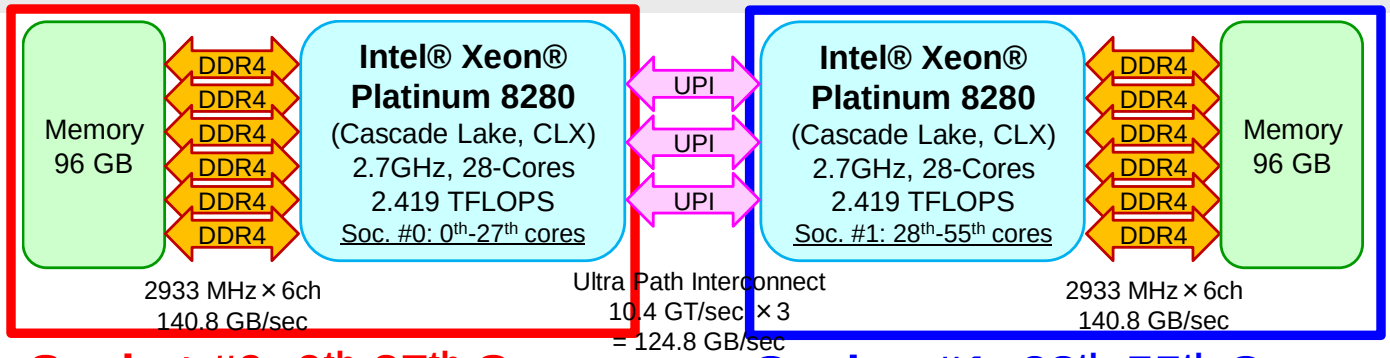


s56.sh: Use 56 cores

```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture2
#PJM -L node=1
#PJM --mpi proc=56
#PJM -L elapse=00:15:00
#PJM -g gt62
#PJM -j
#PJM -e err
#PJM -o s56.lst

mpiexec.hydra -n ${PJM_MPI_PROC} numactl -l ./stream
mpiexec.hydra -n ${PJM_MPI_PROC} ./stream

export I_MPI_PIN_PROCESSOR_LIST=0-55
mpiexec.hydra -n ${PJM_MPI_PROC} numactl -l ./stream
mpiexec.hydra -n ${PJM_MPI_PROC} ./stream
```

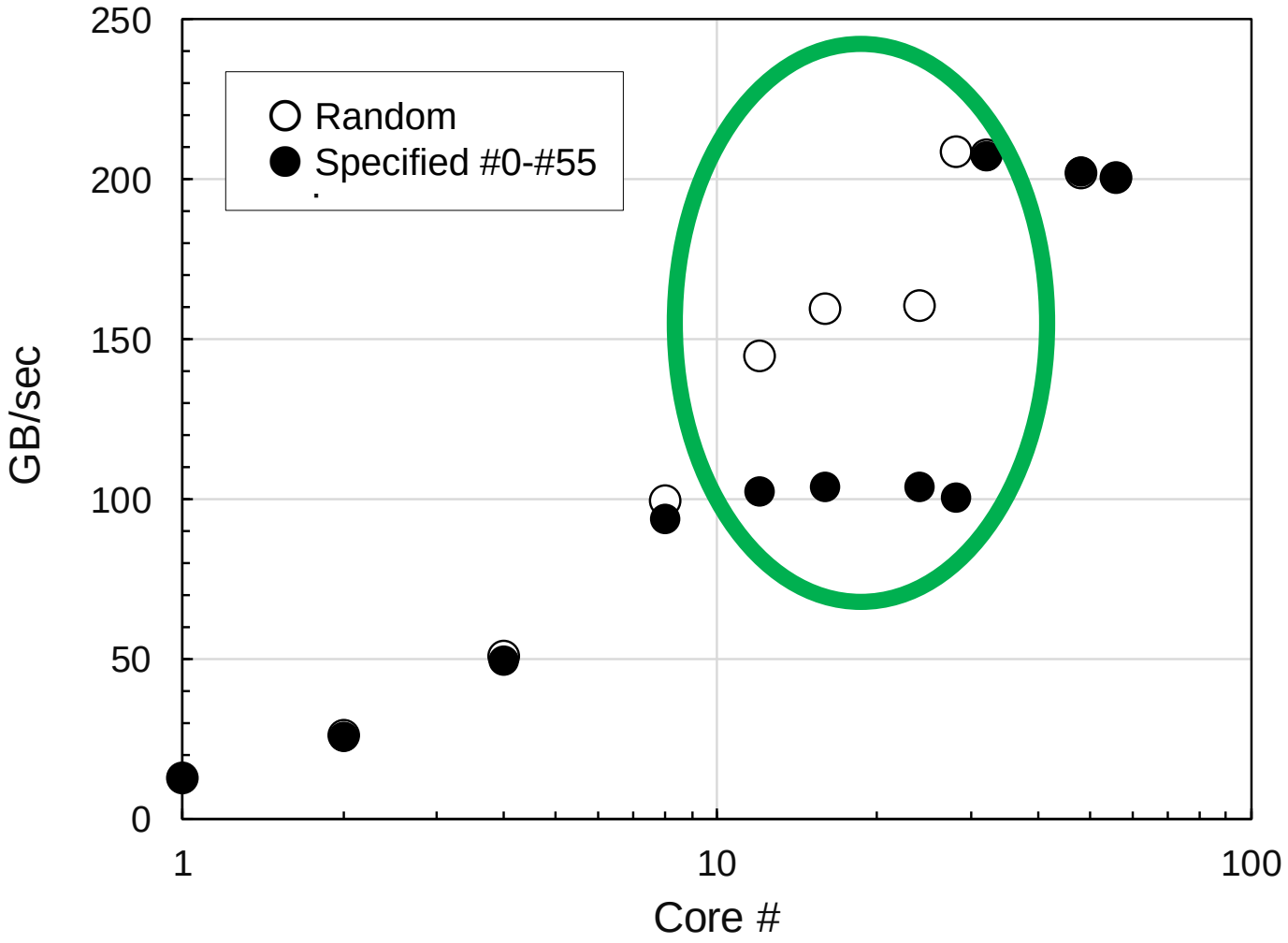
Cores are randomly selected

Core are specified



Triad on a Single Node of OBCX Peak is 281.57 GB/sec.

Memory/cache may be more efficiently utilized in “Random” (12-28 cores)



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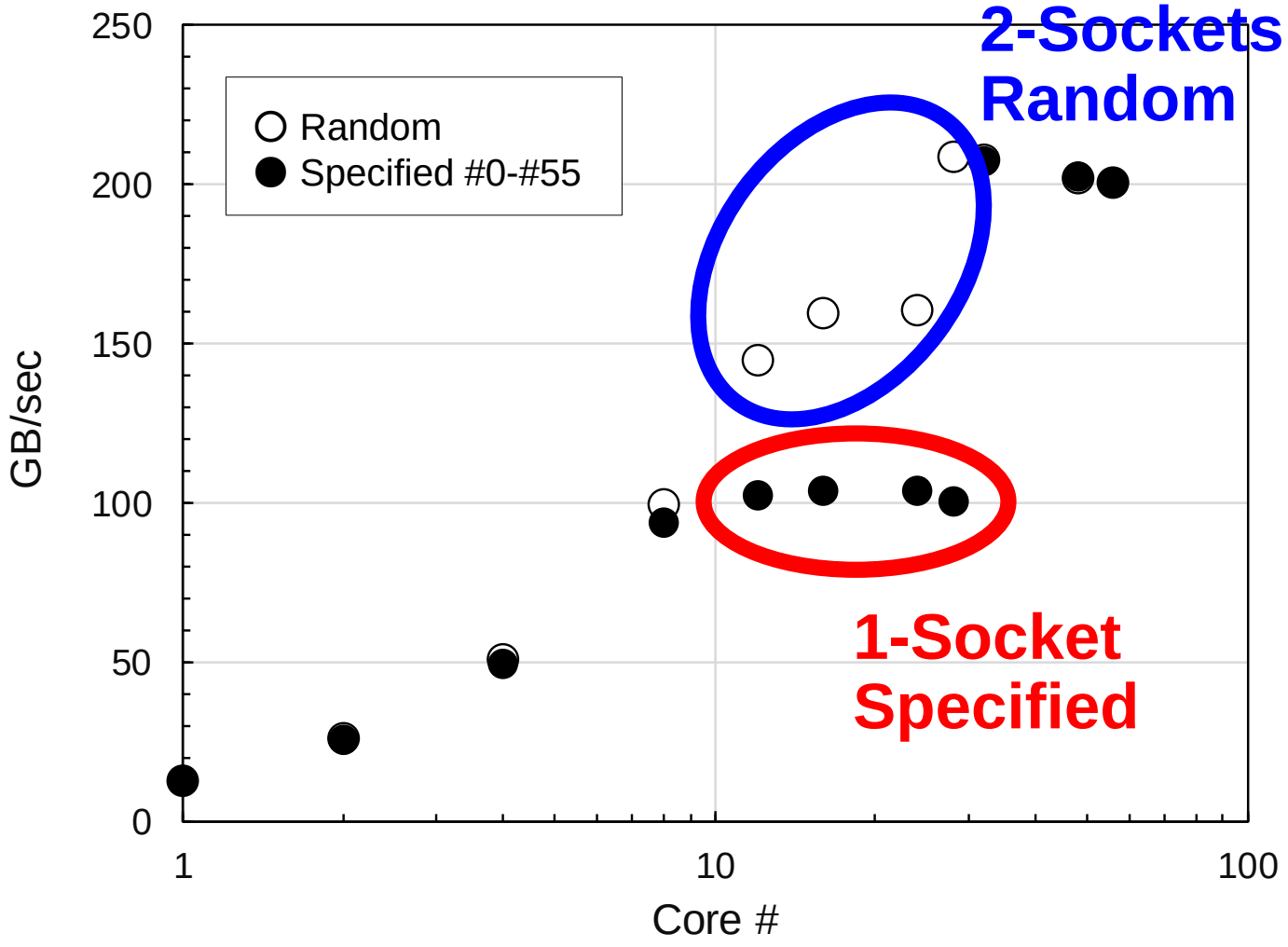
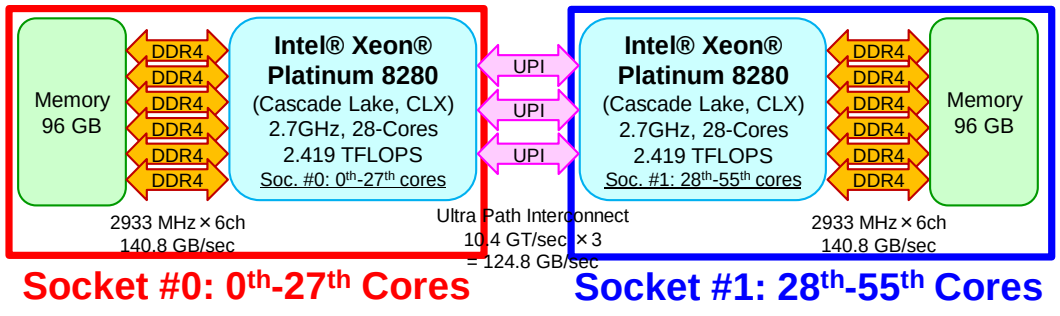
Memory/cache may be more efficiently utilized in “Random” (12-28 cores)





Socket #	Core #	Random	Specified
1	1	1.000	1.000
	2	1.998	1.963
	4	3.912	3.809
	6	5.792	5.682
	8	7.615	7.177
	12	11.089	7.844
	16	12.214	7.968
	24	12.278	7.945
	28	15.962	7.705
2	32	15.901	15.862
	48	15.452	15.497
	56	15.370	15.358

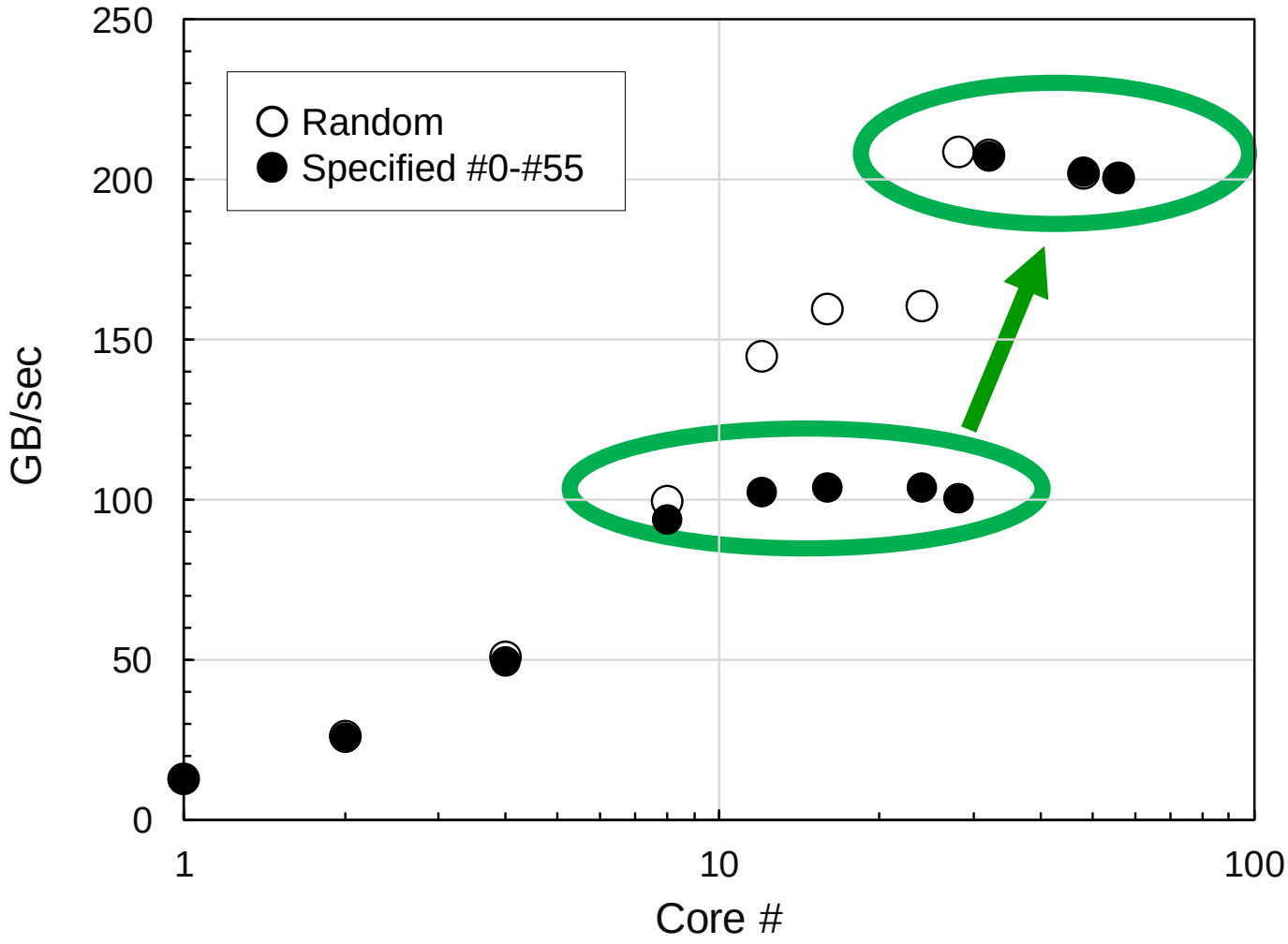
Triad on a Single Node of OBCX

Speed-Up based on Performance with 1 core

- Number of Memory Channels per Socket = 6
  - 6 memory chips can be loaded on each socket
  - Linear speed-up up to 6 cores (processes) on each socket

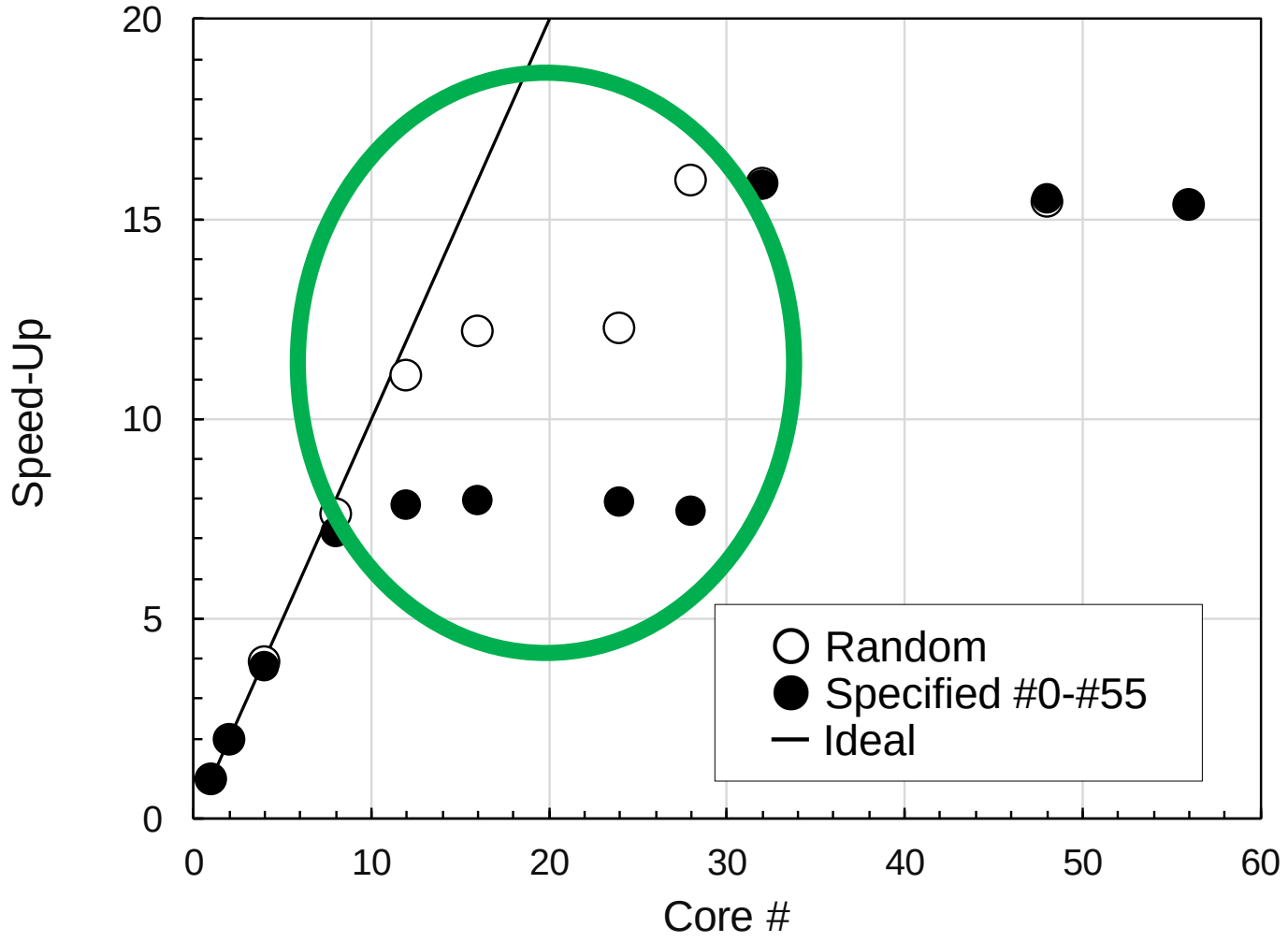
Triad on a Single Node of OBCX Peak is 281.57 GB/sec.

● : Memory BW is constant with 8-28 cores (saturated), Doubled with 32-56 cores



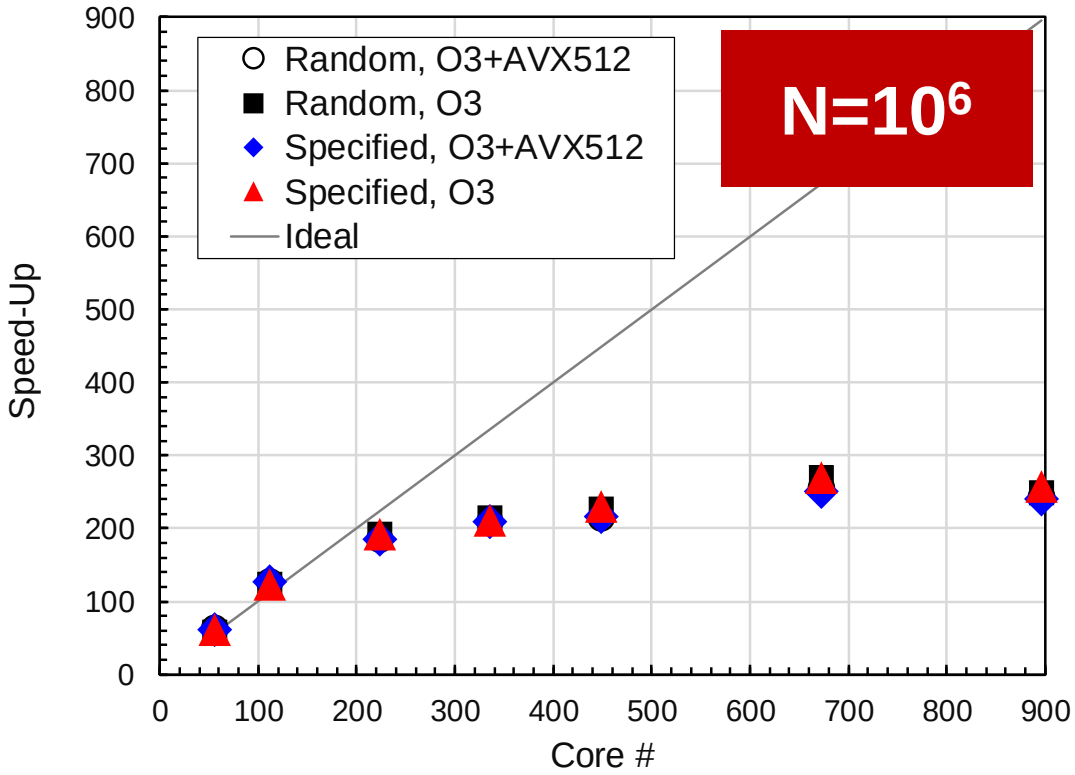
Triad on a Single Node of OBCX Peak is 281.57 GB/sec.

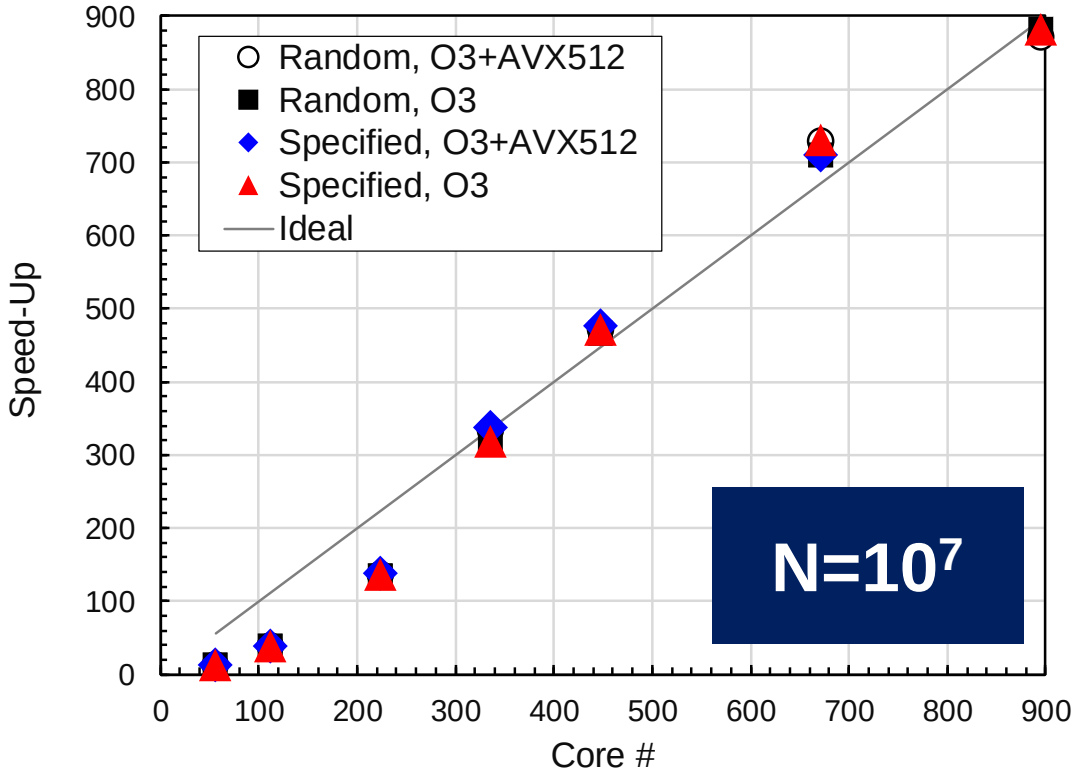
Memory/cache may be more efficiently utilized in “Random” (12-28 cores)



Exercises

- Running the code
- Try various number of processes (1-56)
- OpenMP-version and Single PE version are available
  - Fortran, C
  - Web-site of STREAM
  - http://www.cs.virginia.edu/stream/

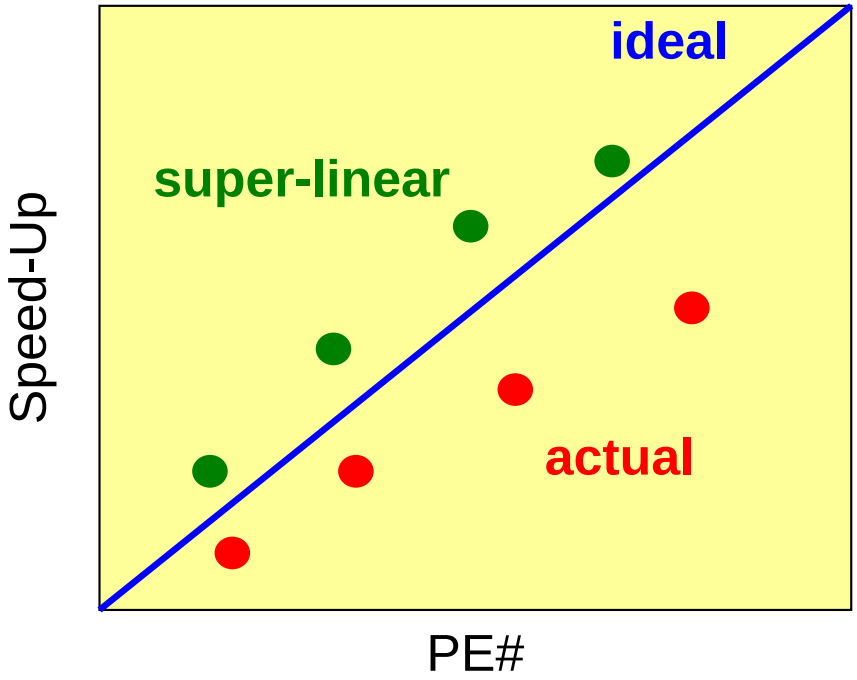




Up to 896 cores (16-nodes)

- Performance at a Single Core= 1.00
- Performance of $N=10^6$ case decreases, as node# increases.
- Performance of $N=10^7$ becomes close to ideal one gradually, and superlinear situation occurs at 672 cores (12 nodes).

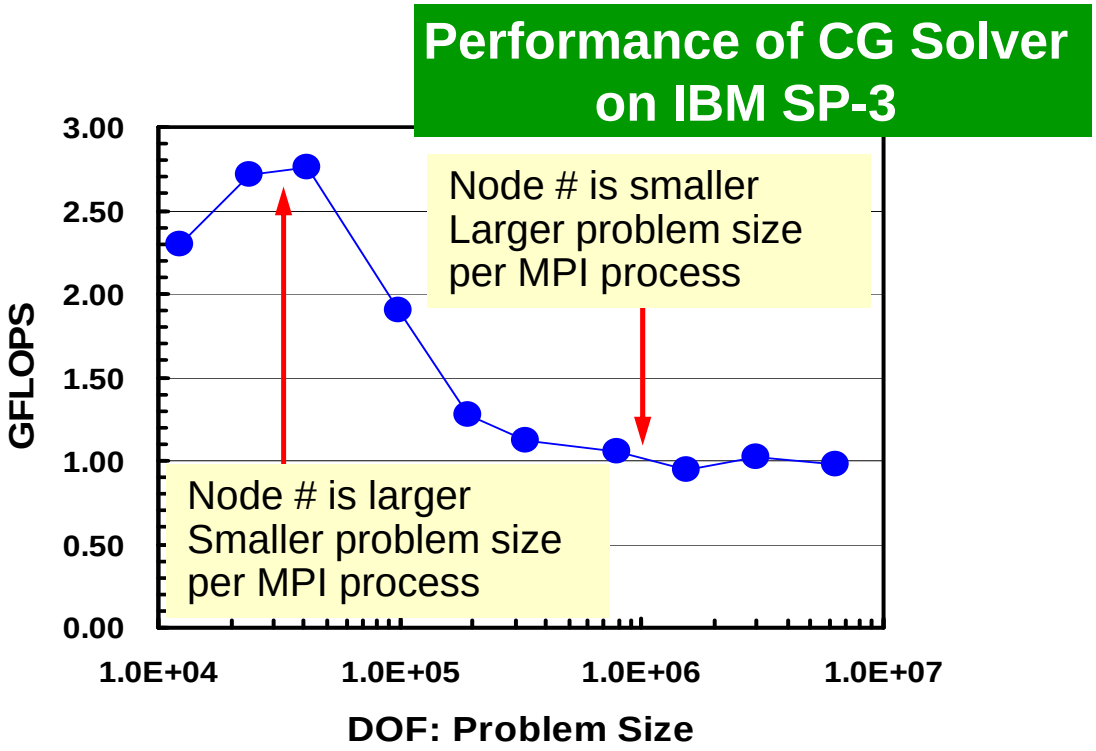
Super-Linear in Strong Scaling

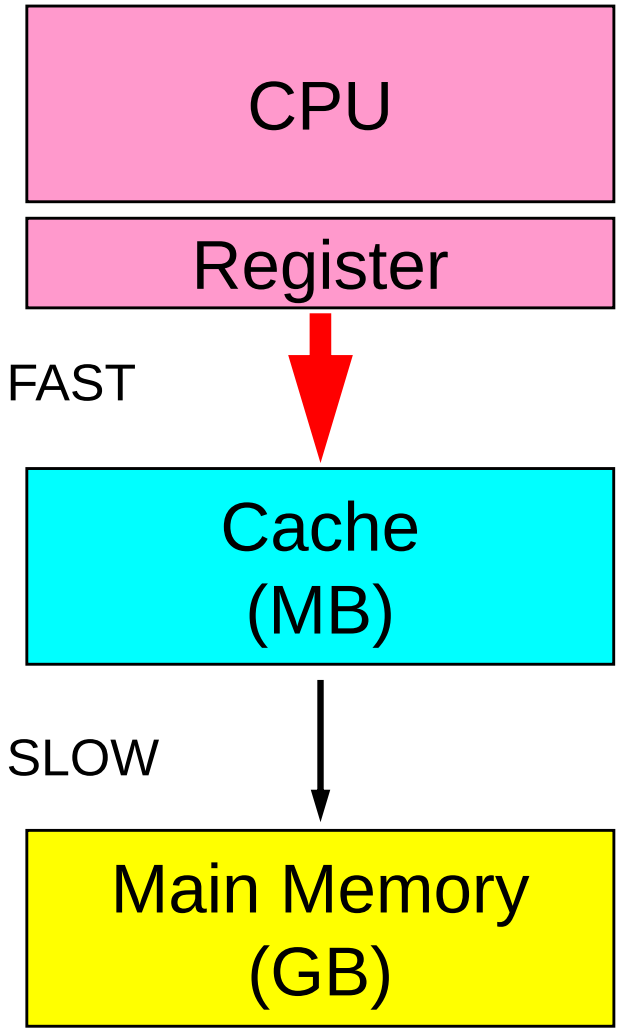


- In strong scaling case where entire problem size is fixed, performance is generally lower than the ideal one due to communication overhead.
- But sometimes, actual performance may be better than the ideal one. This is called “super-linear”

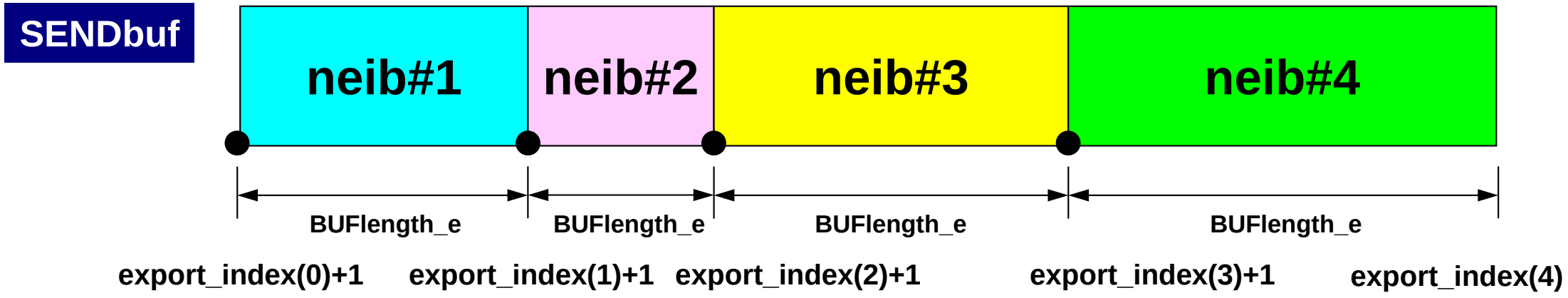
Why does “Super-Linear” happen ?

- Effect of Cache
- In scalar processors, performance for smaller problem is generally better.
  - Cache is well-utilized.





Memory Copy is expensive (1/2)



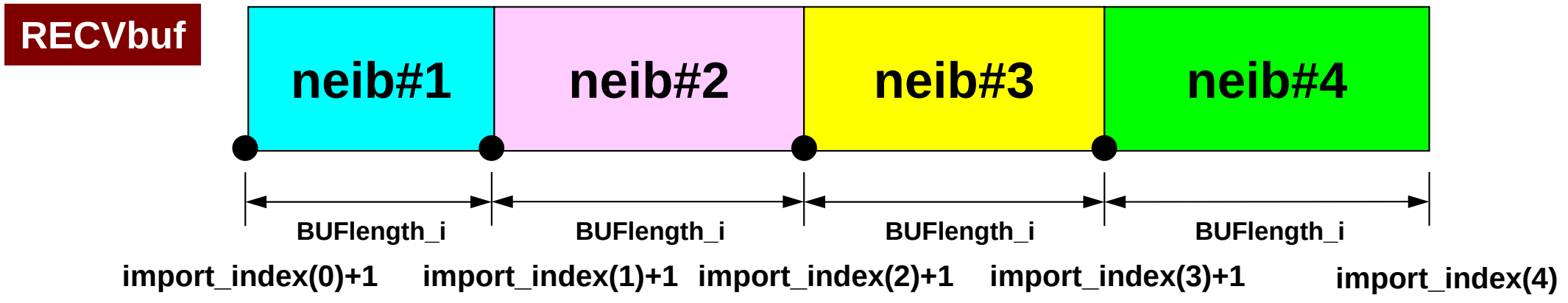

```
 do neib= 1, NEIBPETOT
   do k= export_index(neib-1)+1, export_index(neib)
     kk= export_item(k)
     SENDbuf(k)= VAL(kk)
   enddo
 enddo

 do neib= 1, NEIBPETOT
   iS_e= export_index(neib-1) + 1
   iE_e= export_index(neib  )
   BUFlength_e= iE_e + 1 - iS_e

   call MPI_ISEND                                                  &
&          (SENDbuf(iS_e), BUFlength_e, MPI_INTEGER, NEIBPE(neib), 0,&
&           MPI_COMM_WORLD, request_send(neib), ierr)
 enddo

 call MPI_WAITALL (NEIBPETOT, request_send, stat_send, ierr)
```

Copied to sending buffers

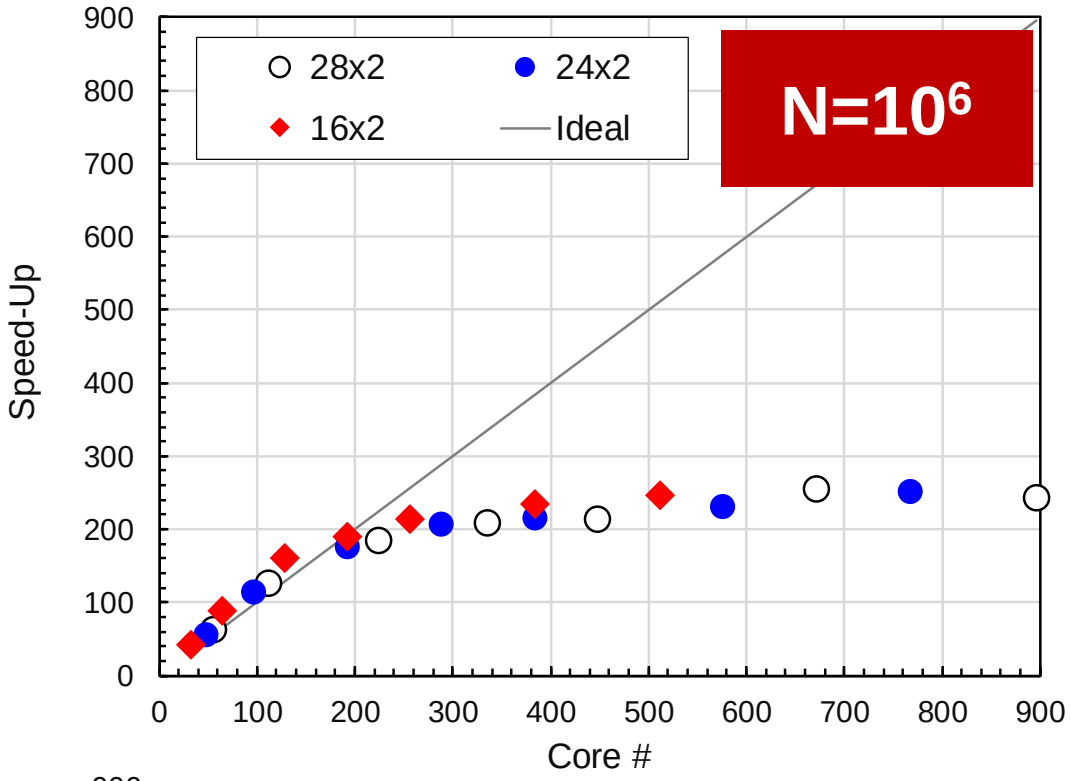
Memory Copy is expensive (2/2)

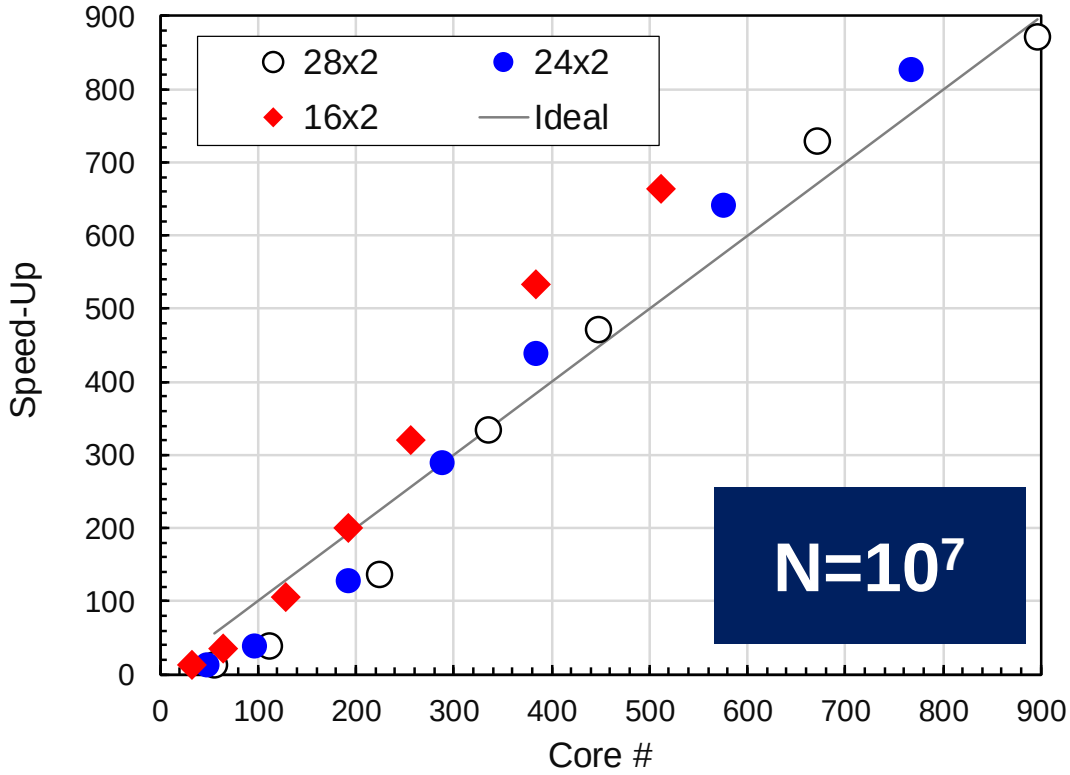
```
 do neib= 1, NEIBPETOT
   iS_i= import_index(neib-1) + 1
   iE_i= import_index(neib  )
   BUFlength_i= iE_i + 1 - iS_i

   call MPI_IRECV                                                    &
&          (RECVbuf(iS_i), BUFlength_i, MPI_INTEGER, NEIBPE(neib), 0,&
&           MPI_COMM_WORLD, request_recv(neib), ierr)
  enddo

  call MPI_WAITALL (NEIBPETOT, request_recv, stat_recv, ierr)

  do neib= 1, NEIBPETOT
     do k= import_index(neib-1)+1, import_index(neib)
       kk= import_item(k)
       VAL(kk)= RECVbuf(k)
     enddo
  enddo
```

Copied from receiving buffer

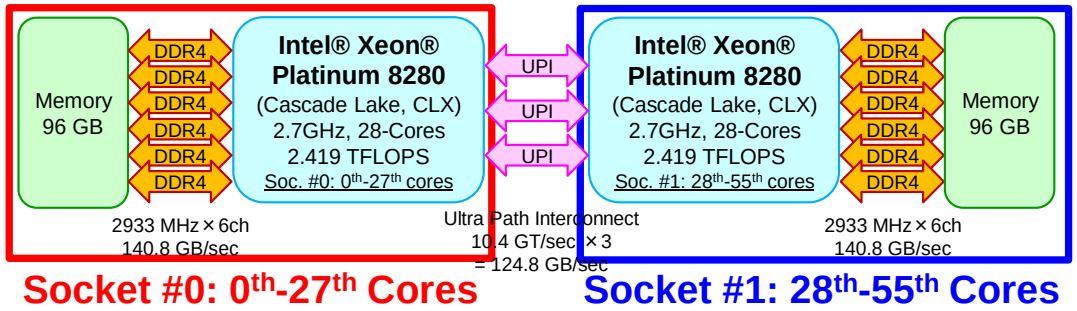


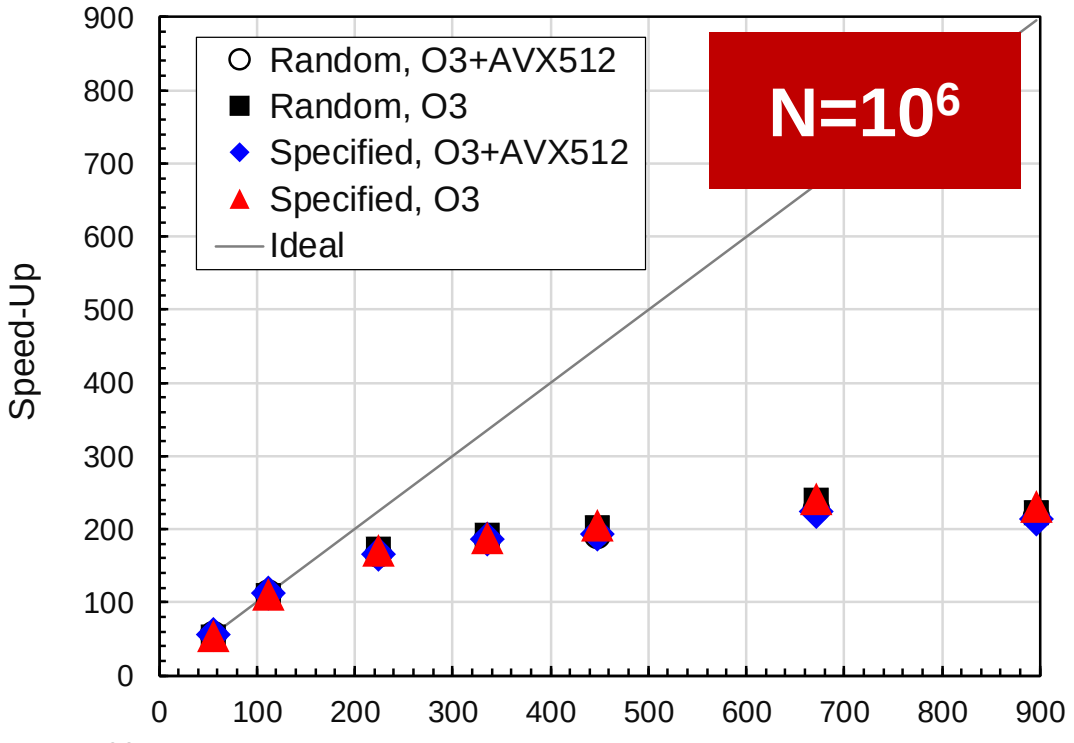
Up to 16-nodes
Random, O3+AVX512

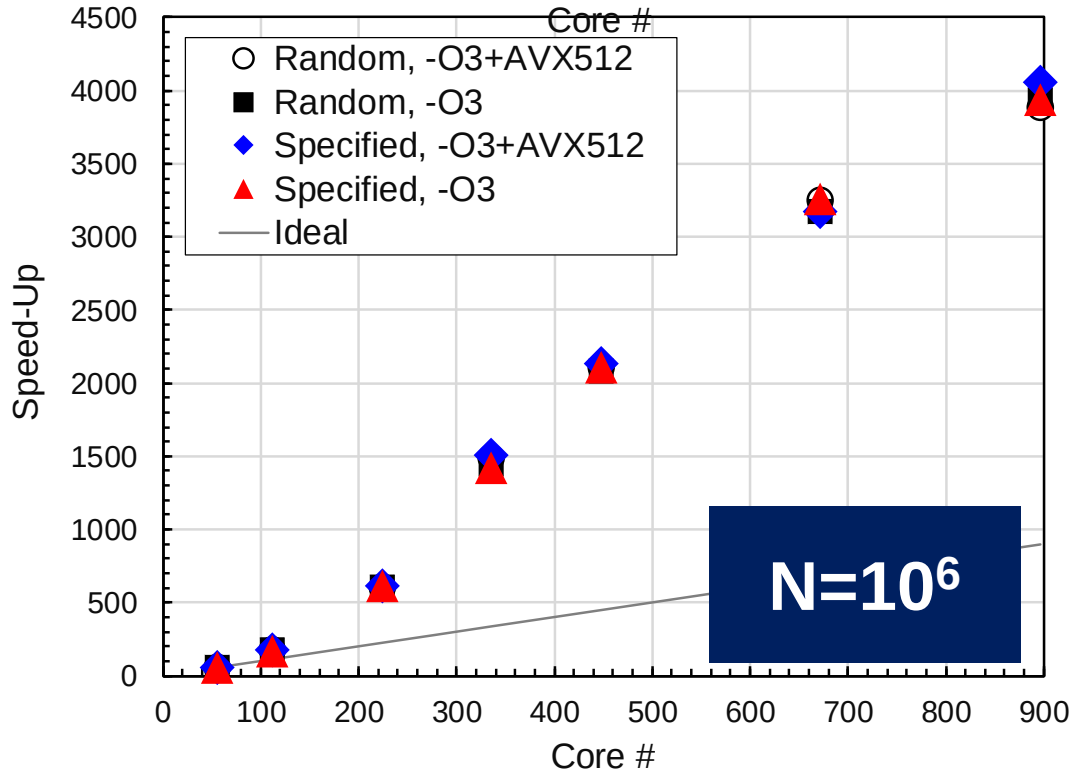




- MPI Processes per Node
  - 16x2, 24x2, 28x2
- N=10^6
  - No diff. if core# is larger
- N=10^7
  - At same core#, 16x2 is the best: Memory is more efficiently utilized



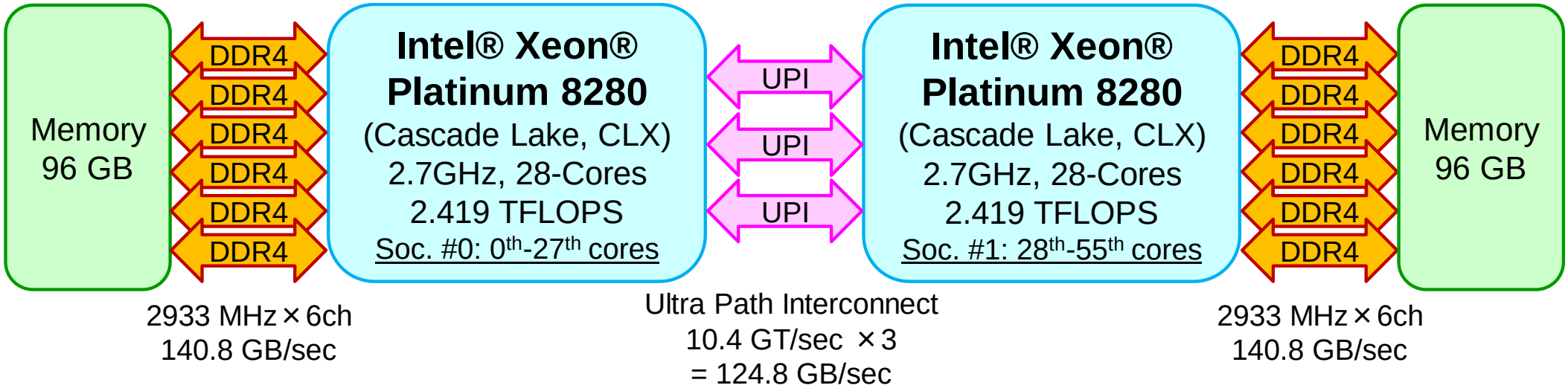




Up to 896 cores (16-nodes)

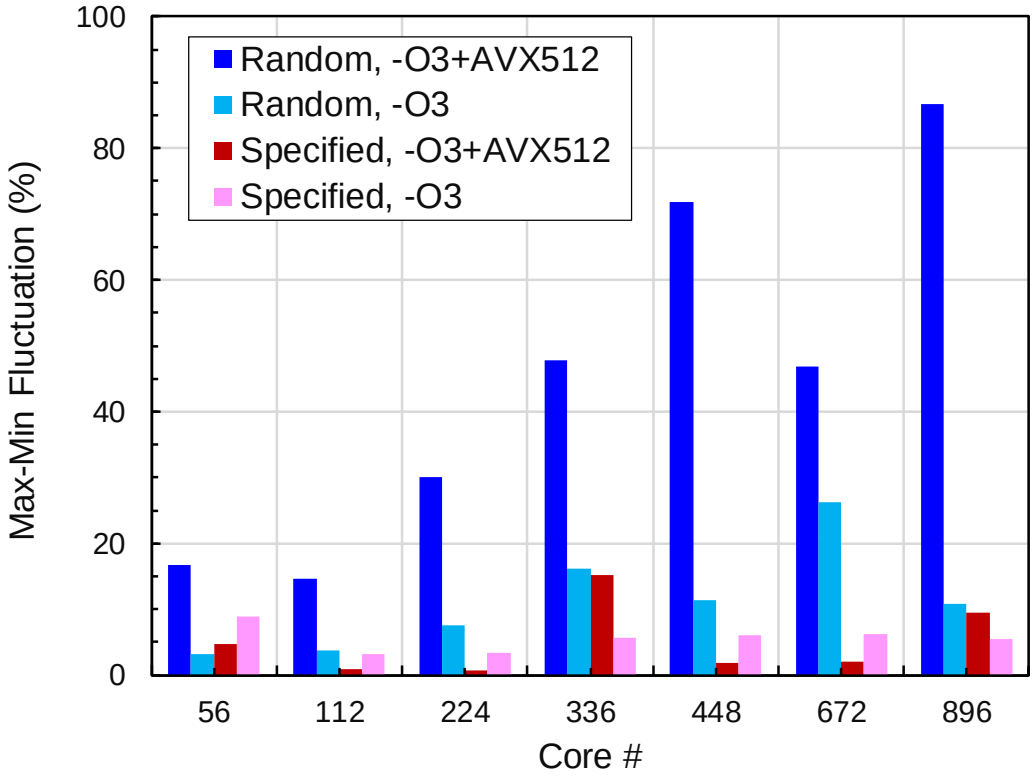
- Performance at 56 cores (= single node) = 56.0
- It is reasonable to evaluate strong scalability for multiple nodes based on the performance by a node with 56 cores
- L2 cache: 1MB/core
- L3: 38.5MB/socket (shared)
- $N=10^6$
  - Significant Comm. Overhead
- $N=10^7$
  - Very Superlinear

Category	Capacity	X-Way Set Associative	Cache Line
L1$Data	32 KB/core	8-Way	64B
L1$Instruction	32 KB/core	8-Way	64B
L2	1.00 MB/core	16-Way	64B
L3	38.5 MB/socket	11-Way	64B

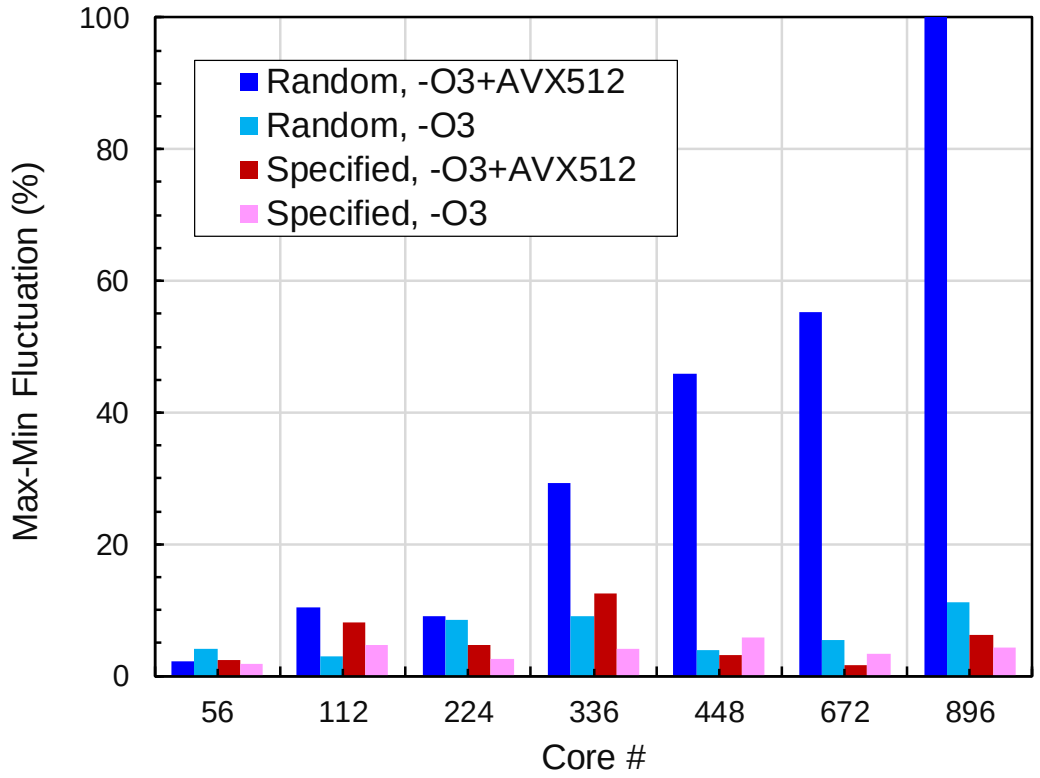


Fluctuation of Computation Time for 5 measurements

N=10^6



N=10^6

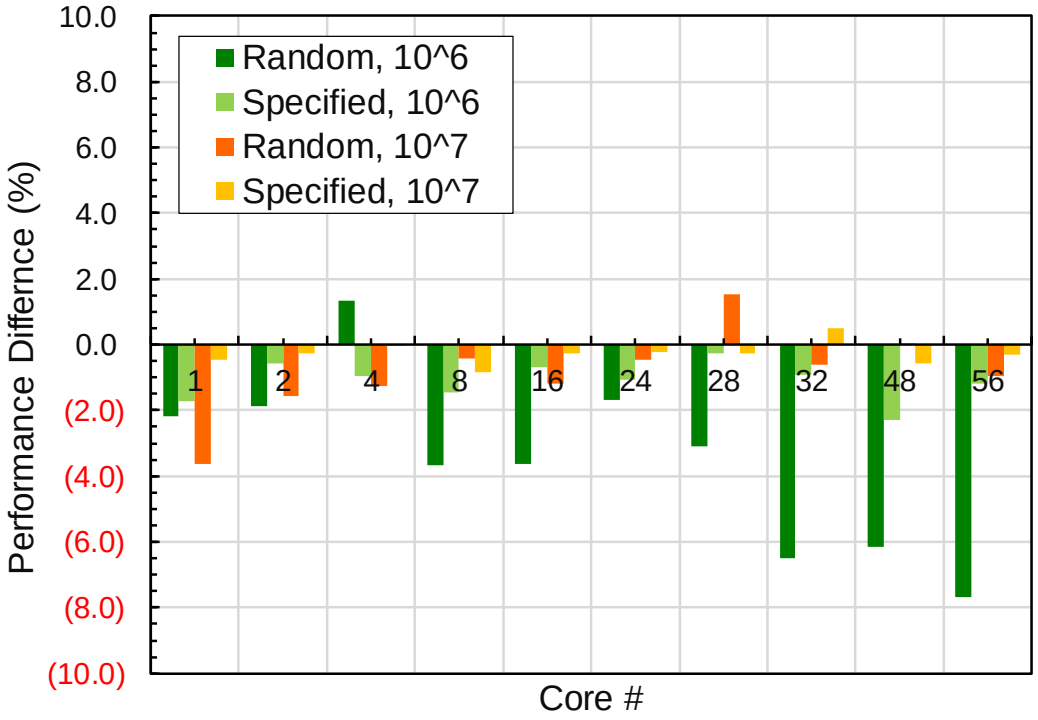


“-O3+AVX512” and “-O3” Best Case of 5 Measurements

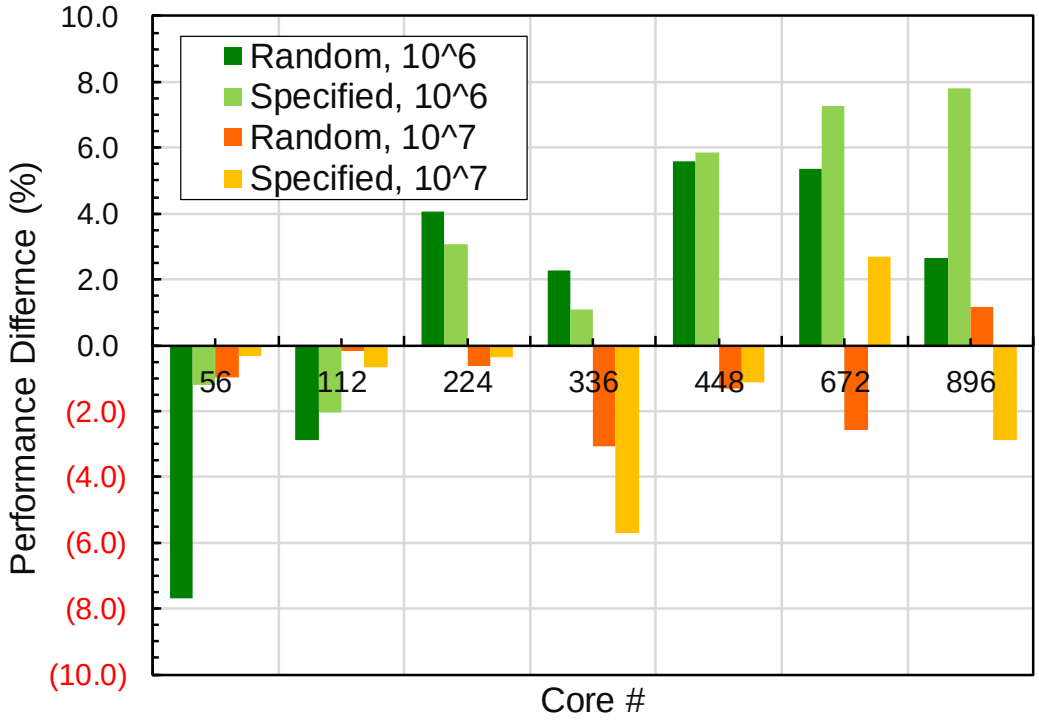
+: -O3 is better, -: -O3 is worse

“-O3” is rather faster in 10^6 cases with many cores

up to 1 node, 56cores



up to 16 nodes, 896 cores



Summary: Parallel FEM

- Proper design of data structure of distributed local meshes.

- Open Technical Issues
  - Parallel Mesh Generation, Parallel Visualization
  - Parallel Preconditioner for Ill-Conditioned Problems
  - Large-Scale I/O

Distributed Local Data Structure for Parallel Computation

- Distributed local data structure for domain-to-doain communications has been introduced, which is appropriate for such applications with sparse coefficient matrices (e.g. FDM, FEM, FVM etc.).
  - SPMD
  - Local Numbering: Internal pts to External pts
  - Generalized communication table

- Everything is easy, if proper data structure is defined:
  - Values at boundary pts are copied into sending buffers
  - Send/Recv
  - Values at external pts are updated through receiving buffers

If numbering of external nodes is continuous in each neighboring process ...

	84	81	85	82	83	86	88	87	
96	57	58	59	60	61	62	63	64	73
95	49	50	51	52	53	54	55	56	74
94	41	42	43	44	45	46	47	48	80
93	33	34	35	36	37	38	39	40	79
92	25	26	27	28	29	30	31	32	78
91	17	18	19	20	21	22	23	24	77
90	9	10	11	12	13	14	15	16	76
89	1	2	3	4	5	6	7	8	75
	65	66	67	68	69	70	71	72	

[A]{p}= {q} (Original), 1d.f

```
 allocate (stat_send(MPI_STATUS_SIZE,NEIBPETOT))
 allocate (stat_recv(MPI_STATUS_SIZE,NEIBPETOT))
 allocate (request_send(NEIBPETOT)); allocate (request_recv(NEIBPETOT))

 do neib= 1, NEIBPETOT
   do k= export_index(neib-1)+1, export_index(neib)
     kk= export_item(k)
     SENDbuf(k)= W(kk,P)
   enddo
 enddo

 do neib= 1, NEIBPETOT
   is   = export_index(neib-1) + 1
   len_s= export_index(neib) - export_index(neib-1)
   call MPI_Isend (SENDbuf(is), len_s, MPI_DOUBLE_PRECISION,        &
&                  NEIBPE(neib), 0, MPI_COMM_WORLD,                 &
&                  request_send(neib), ierr)
 enddo

 do neib= 1, NEIBPETOT
   ir   = import_index(neib-1) + 1
   len_r= import_index(neib) - import_index(neib-1)
   call MPI_Irecv (RECVbuf(ir), len_r, MPI_DOUBLE_PRECISION,        &
&                  NEIBPE(neib), 0, MPI_COMM_WORLD,                 &
&                  request_recv(neib), ierr)
 enddo
 call MPI_Waitall (NEIBPETOT, request_recv, stat_recv, ierr)

 do neib= 1, NEIBPETOT
   do k= import_index(neib-1)+1, import_index(neib)
     kk= import_item(k)
     W(kk,P)= RECVbuf(k)
   enddo
 enddo
 call MPI_Waitall (NEIBPETOT, request_send, stat_send, ierr)
```

[A]{p}= {q} (Mod.): No Copy for RECV 1d2.f

```
      allocate (stat_send(MPI_STATUS_SIZE,2*NEIBPETOT))
      allocate (request_send(2*NEIBPETOT))

      do neib= 1, NEIBPETOT
        do k= export_index(neib-1)+1, export_index(neib)
          kk= export_item(k)
          SENDbuf(k)= W(kk,P)
        enddo
      enddo

      do neib= 1, NEIBPETOT
        is   = export_index(neib-1) + 1
        len_s= export_index(neib) - export_index(neib-1)
        call MPI_Isend (SENDbuf(is), len_s, MPI_DOUBLE_PRECISION,        &
     &                  NEIBPE(neib), 0, MPI_COMM_WORLD,                 &
     &                  request_send(neib), ierr)
      enddo

      do neib= 1, NEIBPETOT
        ir   = import_index(neib-1) + 1
        len_r= import_index(neib) - import_index(neib-1)
        call MPI_Irecv (W(ir+N,P), len_r, MPI_DOUBLE_PRECISION,        &
     &                  NEIBPE(neib), 0, MPI_COMM_WORLD,                 &
     &                  request_send(neib+NEIBPETOT), ierr)
      enddo

      call MPI_Waitall (2*NEIBPETOT, request_send, stat_send, ierr)
```

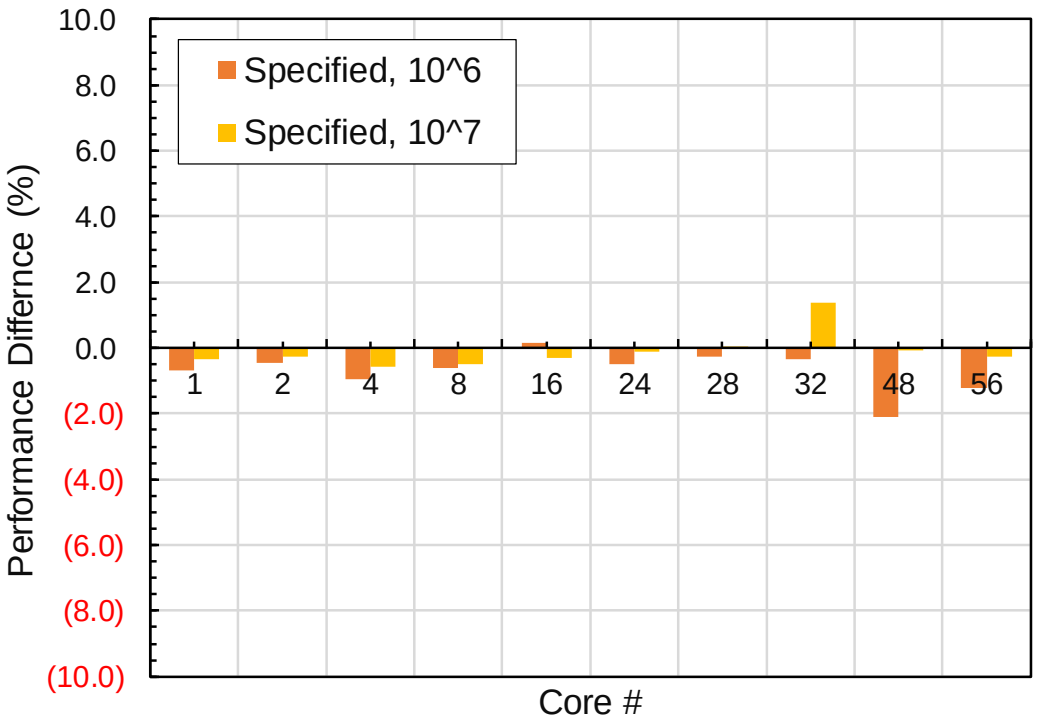
“Original” and “Modified” Best Case of 5 Measurements

+: “Modified” is better, -: “Modified” is worse

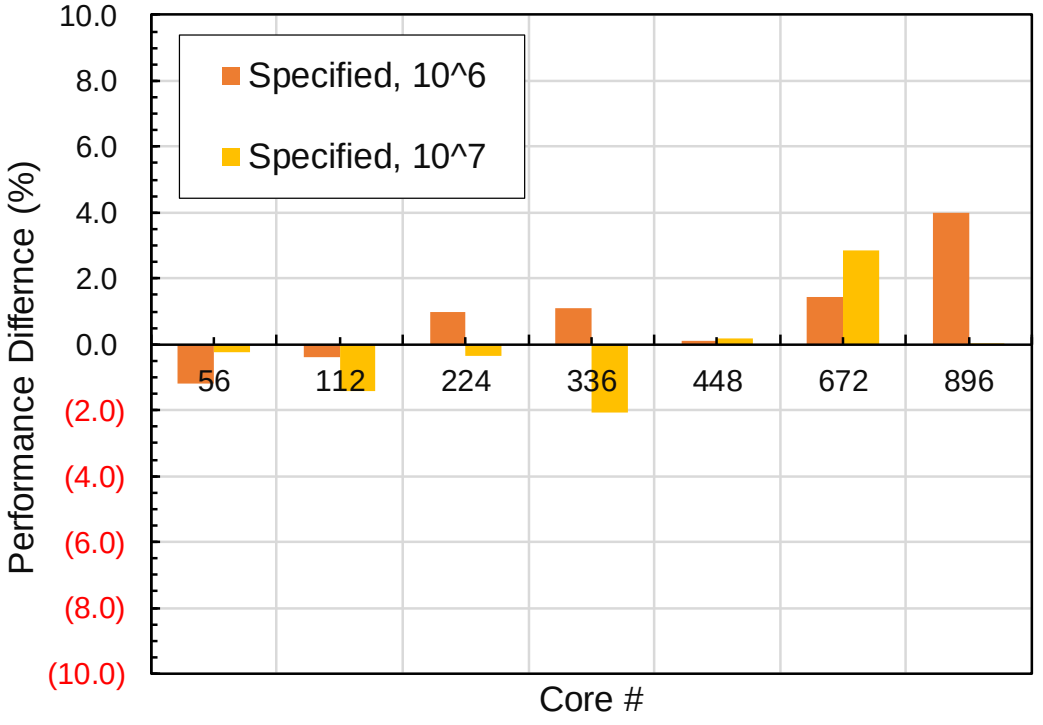
“Modified” is faster in 10^6 cases with many cores

But differences are small

up to 1 node, 56cores



up to 16 nodes, 896 cores



Further Investigations

- Compare with/without NUMA control

```
mpiexec.hydra -n ${PJM_MPI_PROC} numactl -l ./a.out
mpiexec.hydra -n ${PJM_MPI_PROC} ./a.out
```

- Explain why number of iterations does not change, as number of MPI processes changes.