

Introduction to Programming by MPI for Parallel FEM Report S1 & S2 in C (1/2)

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Motivation for Parallel Computing (and this class)

- Large-scale parallel computer enables fast computing in large-scale scientific simulations with detailed models. Computational science develops new frontiers of science and engineering.
- Why parallel computing ?
 - faster & larger
 - “larger” is more important from the view point of “new frontiers of science & engineering”, but “faster” is also important.
 - + more complicated
 - Ideal: Scalable
 - Solving N^x scale problem using N^x computational resources during same computation time.

Scalable, Scaling, Scalability

- Solving N^x scale problem using N^x computational resources during same computation time
 - for large-scale problems: Weak Scaling, Weak Scalability
 - e.g. CG solver: more iterations needed for larger problems
- Solving a problem using N^x computational resources during $1/N$ computation time
 - for faster computation: Strong Scaling, Strong Scalability

Overview

- What is MPI ?
- Your First MPI Program: Hello World
- Collective Communication
- Point-to-Point Communication

What is MPI ? (1/2)

- Message Passing Interface
- “Specification” of message passing API for distributed memory environment
 - Not a program, Not a library
 - <http://www.mcs.anl.gov/mpi/www/>
 - <https://www.mpi-forum.org/docs/>
- History
 - 1992 MPI Forum
 - <https://www.mpi-forum.org/>
 - 1994 MPI-1
 - 1997 MPI-2: MPI I/O
 - 2012 MPI-3: Fault Resilience, Asynchronous Collective
- Implementation
 - mpich: ANL (Argonne National Laboratory), OpenMPI, MVAPICH
 - H/W vendors
 - C/C++, FORTRAN, Java ; Unix, Linux, Windows, Mac OS

What is MPI ? (2/2)

- “mpich” (free) is widely used
 - supports MPI-2 spec. (partially)
 - MPICH2 after Nov. 2005.
 - <http://www.mcs.anl.gov/mpi/>
- Why MPI is widely used as *de facto standard* ?
 - Uniform interface through MPI forum
 - Portable, can work on any types of computers
 - Can be called from Fortran, C, etc.
 - mpich
 - free, supports every architecture
- PVM (Parallel Virtual Machine) was also proposed in early 90’s but not so widely used as MPI

References

- W.Gropp et al., Using MPI second edition, MIT Press, 1999.
- M.J.Quinn, Parallel Programming in C with MPI and OpenMP, McGrawhill, 2003.
- W.Gropp et al., MPI: The Complete Reference Vol.I, II, MIT Press, 1998.
- <http://www.mcs.anl.gov/mpi/www/>
 - API (Application Interface) of MPI

How to learn MPI (1/2)

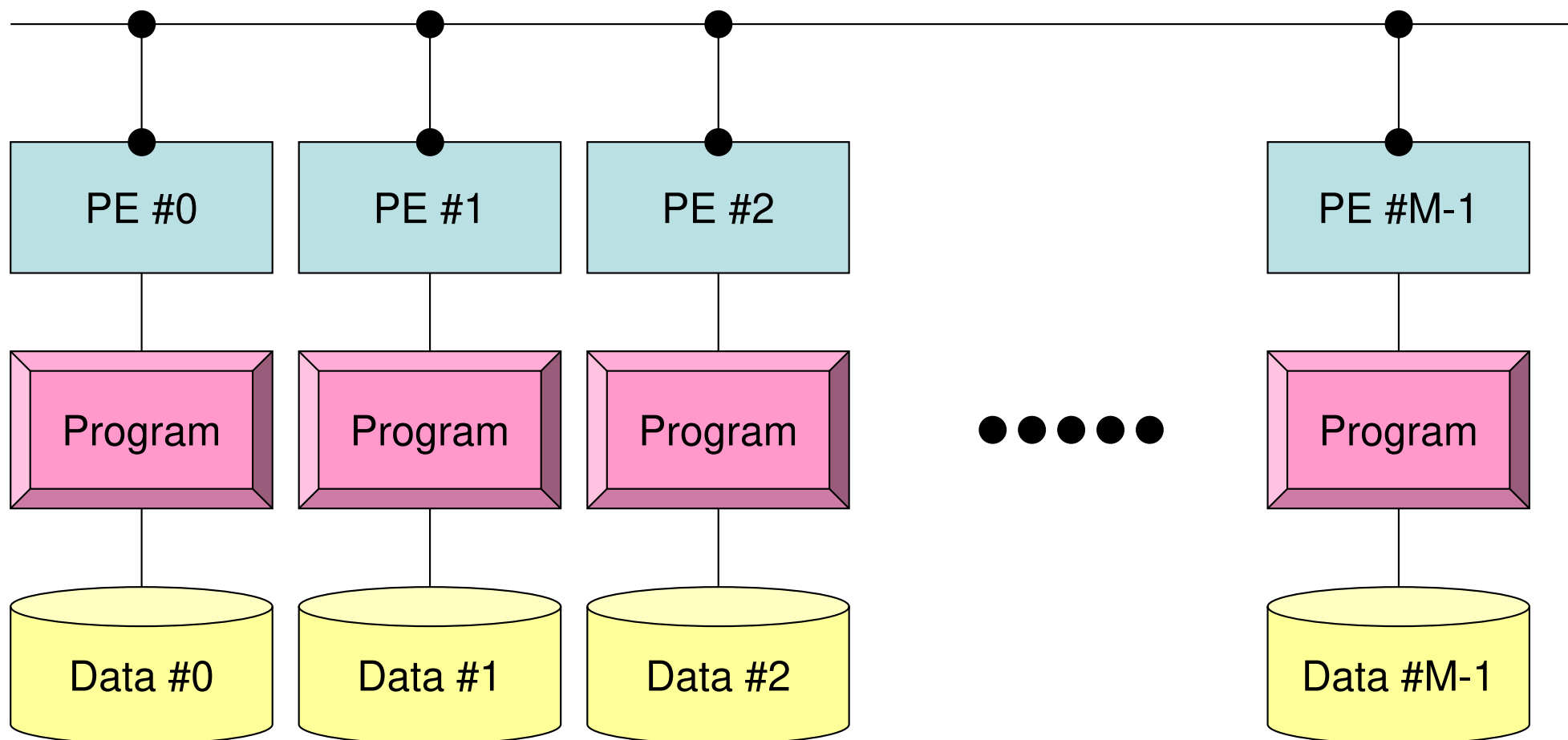
- Grammar
 - 10-20 functions of MPI-1 will be taught in the class
 - although there are many convenient capabilities in MPI-2
 - If you need further information, you can find information from web, books, and MPI experts.
- Practice is important
 - Programming
 - “Running the codes” is the most important
- Be familiar with or “grab” the idea of SPMD/SIMD op’s
 - Single Program/Instruction Multiple Data
 - Each process does same operation for different data
 - Large-scale data is decomposed, and each part is computed by each process
 - Global/Local Data, Global/Local Numbering

PE: Processing Element
Processor, Domain, Process

SPMD

You understand 90% MPI, if
you understand this figure.

```
mpirun -np M <Program>
```



Each process does same operation for different data

Large-scale data is decomposed, and each part is computed by each process

It is ideal that parallel program is not different from serial one except communication.

Some Technical Terms

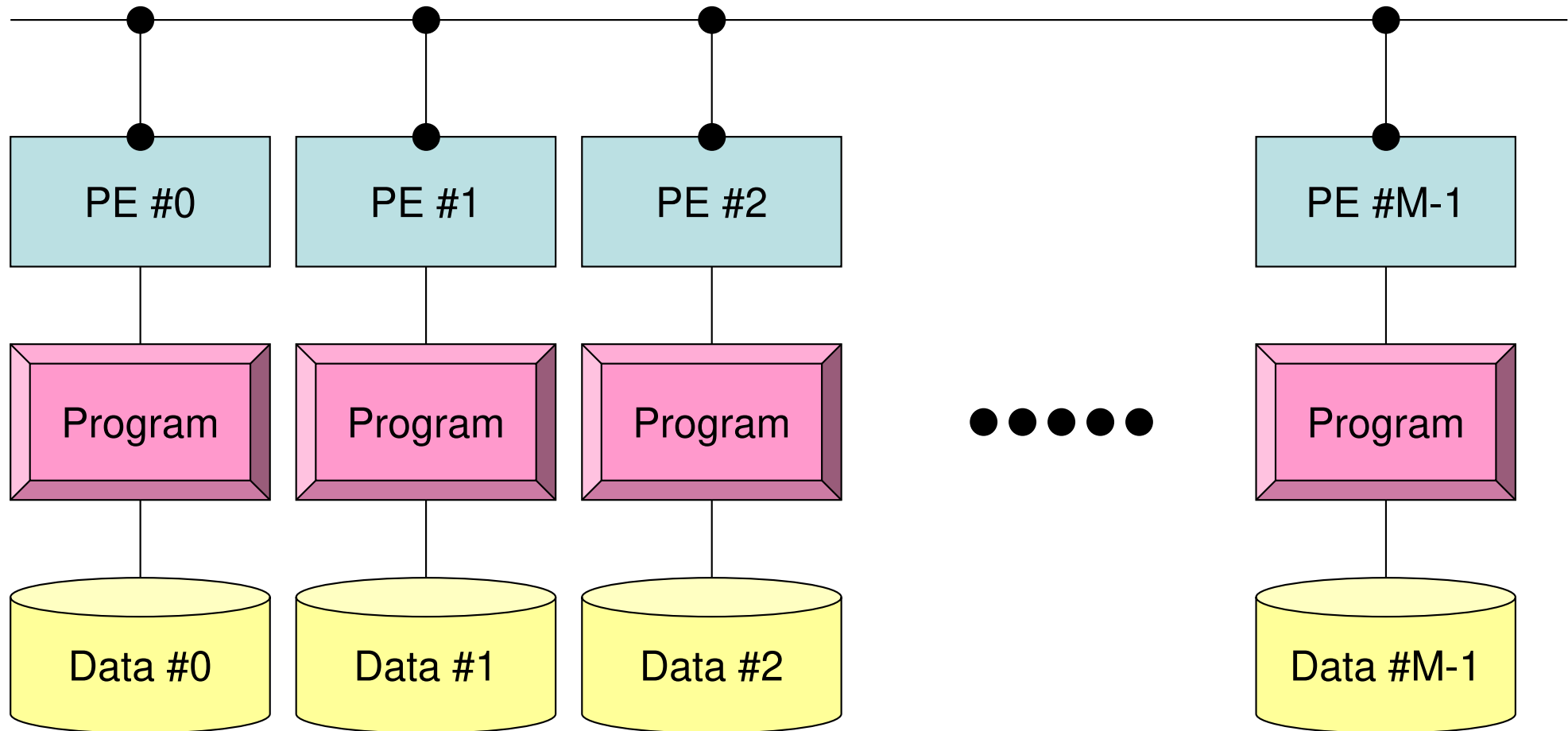
- Processor, Core
 - Processing Unit (H/W), Processor=Core for single-core proc's
- Process
 - Unit for MPI computation, nearly equal to “core”
 - Each core (or processor) can host multiple processes (but not efficient)
- PE (Processing Element)
 - PE originally mean “processor”, but it is sometimes used as “process” in this class. Moreover it means “domain” (next)
 - In multicore proc's: PE generally means “core”
- Domain
 - domain=process (=PE), each of “MD” in “SPMD”, each data set
- Process ID of MPI (ID of PE, ID of domain) starts from “0”
 - if you have 8 processes (PE's, domains), ID is 0~7

PE: Processing Element
Processor, Domain, Process

SPMD

You understand 90% MPI, if
you understand this figure.

```
mpirun -np M <Program>
```



Each process does same operation for different data

Large-scale data is decomposed, and each part is computed by each process

It is ideal that parallel program is not different from serial one except communication.

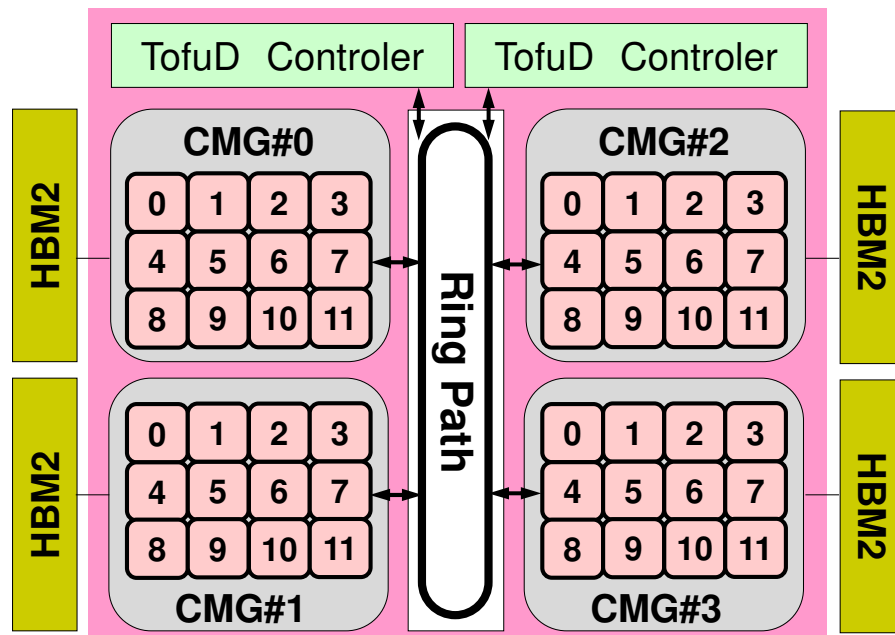
How to learn MPI (2/2)

- NOT so difficult.
- Therefore, 5-6-hour lectures are enough for just learning grammar of MPI.
- Grab the idea of SPMD !

Schedule

- MPI
 - Basic Functions
 - Collective Communication
 - Point-to-Point (or Peer-to-Peer) Communication
- 90 min. x 4-5 lectures
 - Collective Communication
 - Report S1
 - Point-to-Point Communication
 - Report S2: Parallelization of 1D code
 - At this point, you are almost an expert of MPI programming.

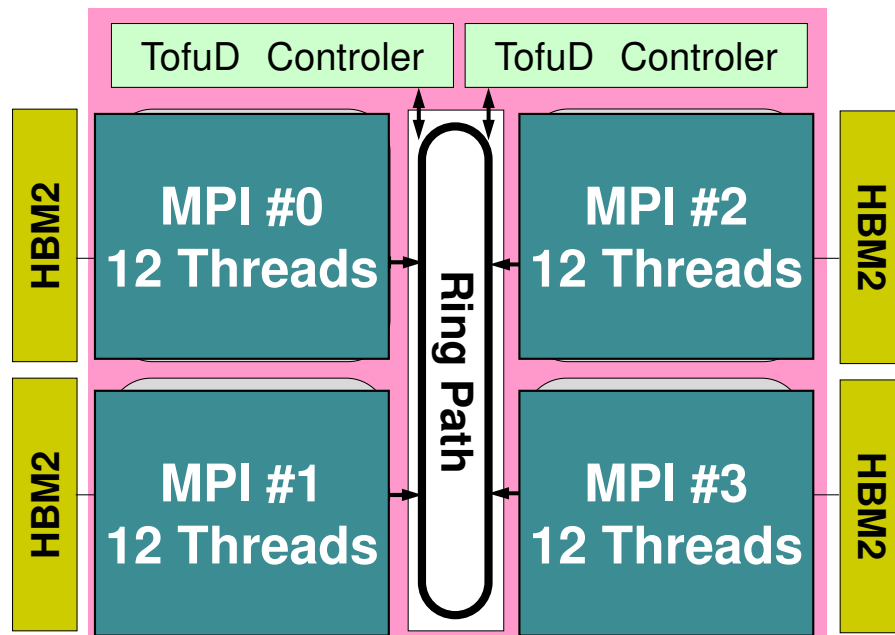
A64FX Processor on Odyssey



Name	A64FX
Processor # (Core #)	1 (48+ 2or4 Assistant Cores)
Frequency	2.2 GHz
Peak Performance	3.3792 TFLOPS
Memory Size	32 GiB
Memory Bandwidth	1,024 GB/s
L1 Cache	64 KiB/core (Inst/Data)
L2 Cache	8 MiB/CMG

- 4 CMG's (Core Memory Group), 12 cores/CMG
 - 48 Cores/Node (Processor)
 - $2.2\text{GHz} \times 32\text{DP} \times 48 = 3379.2 \text{ GFLOPS} = 3.3792 \text{ TFLOPS}$
- NUMA Architecture (Non-Uniform Memory Access)
 - Each core of a CMG can access to the memory on other CMG's
 - Utilization of the local memory is more efficient
- Multiple Nodes: Recommended Programming Model
 - 1-MPI Process for each CMG with 12 OpenMP threads
 - 4-MPI processes for each node

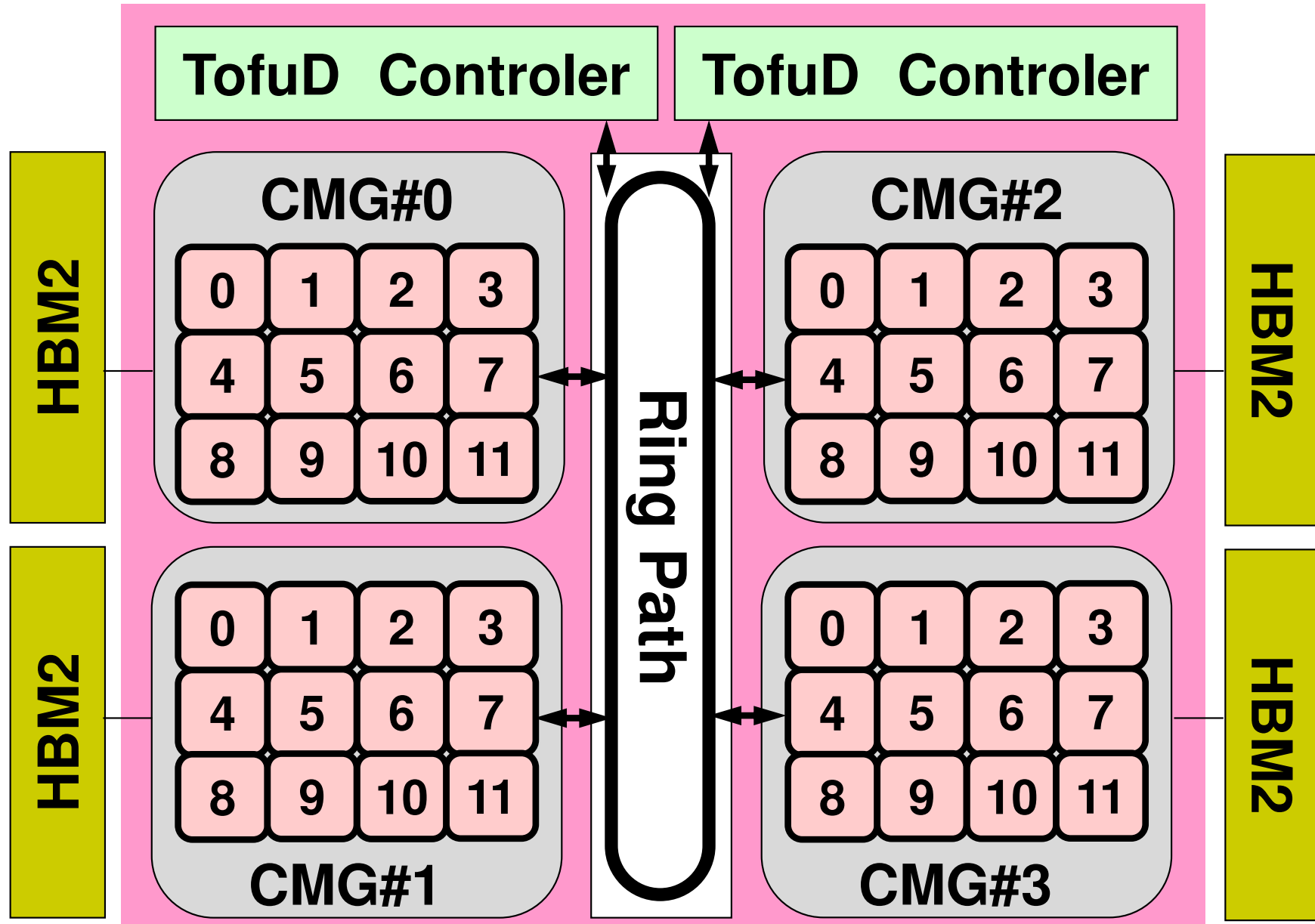
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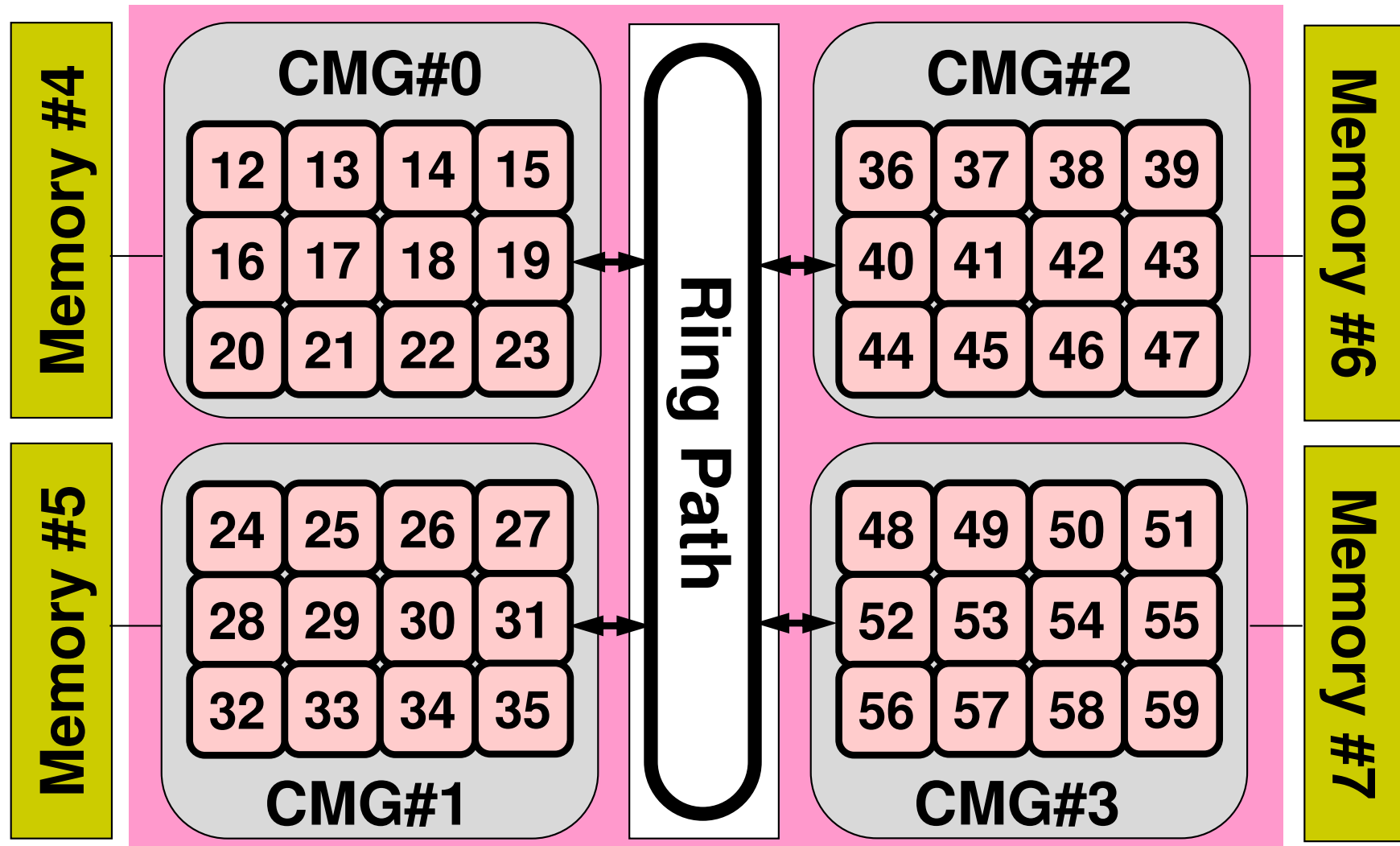
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A64FX: CMG (Core Memory Group)



ID of CMGs, Cores, Memory's (1/2)

CMG:#0-#3, Core:#12-59, Memory:#4-#7



- What is MPI ?
- **Your First MPI Program: Hello World**
- Collective Communication
- Point-to-Point Communication

Login to Odyssey

```
ssh t18\*\*\*@wisteria.cc.u-tokyo.ac.jp
```

Create directory

```
>$ cd /work/gt18/t18***
```

```
>$ mkdir pFEM
```

```
>$ cd pFEM
```

```
>$ module load fj
```

Please type this every time you log-in !!

In this class this top-directory is called **<\$O-TOP>**.
Files are copied to this directory.

Under this directory, **S1**, **S2**, **S1-ref** are created:

```
<$O-S1> = <$O-TOP>/mpi/S1
```

```
<$O-S2> = <$O-TOP>/mpi/S2
```

Odyssey

PC

Copying files on Odyssey

FORTRAN

```
>$ cd /work/gt18/t18XXX/pFEM
>$ module load fj
>$ cp /work/gt00/z30088/pFEM/F/s1-f.tar ○
>$ tar xvf s1-f.tar
```

C

```
>$ cd /work/gt18/t18XXX/pFEM
>$ module load fj
>$ cp /work/gt00/z30088/pFEM/C/s1-c.tar ○
>$ tar xvf s1-c.tar
```

Confirmation

```
>$ ls
  mpi

>$ cd mpi/S1
```

This directory is called as $\langle \$O-S1 \rangle$.

$\langle \$O-S1 \rangle = \langle \$O-TOP \rangle / \text{mpi} / S1$

First Example

hello.f

```
implicit REAL*8 (A-H,O-Z)
include 'mpif.h'
integer :: PETOT, my_rank, ierr

call MPI_INIT      (ierr)
call MPI_COMM_SIZE (MPI_COMM_WORLD, PETOT, ierr )
call MPI_COMM_RANK (MPI_COMM_WORLD, my_rank, ierr )

write (*,'(a,2i8)') 'Hello World FORTRAN', my_rank, PETOT

call MPI_FINALIZE (ierr)

stop
end
```

hello.c

```
#include "mpi.h"
#include <stdio.h>
int main(int argc, char **argv)
{
    int n, myid, numprocs, i;

    MPI_Init(&argc,&argv);
    MPI_Comm_size(MPI_COMM_WORLD,&numprocs);
    MPI_Comm_rank(MPI_COMM_WORLD,&myid);
    printf ("Hello World %d\n", myid);
    MPI_Finalize();
}
```

Compiling hello.f/c

```
>$ cd /work/gt18/t18XXX/pFEM/mpi/S1  
>$ module load fj  
>$ mpifrtpx -Kfast hello.f  
>$ mpifccpx -Nclang -Kfast hello.c
```

FORTTRAN

\$> "mpifrtpx":
required compiler & libraries are included for Fujitsu's FORTRAN90+MPI

C

\$> "mpifccpx":
required compiler & libraries are included for Fujitsu's C+MPI

C-Compiler : 2-modes

<p>trad (-Nnoclang) (default)</p>	<ul style="list-style-type: none"> • Based on Fujitsu's compiler developed for K and PRIMEHPC FX100 or older • Compatible with Fujitsu's Traditional Compilers • C89/C99/C11, OpenMP 3.1/OpenMP 4.5 (partially) • Default (-Nnoclang) • Generally slow for the materials in this class • make -f make-org (make-o)
<p>clang (-Nclang)</p>	<ul style="list-style-type: none"> • Based on Clang/LLVM Compilers (Open Source) • Suitable for using Most Updated Capability's, and for using OSS (Open Source Software) • C89/C99/C11, OpenMP 4.5/OpenMP 5.0 (partially) • Generally faster than "trad" modes, difference between "trad" and "clang" is smaller for optimized codes • In this class, default is "clang" mode • make -f makeec (Makefile)

Running Job

- Batch Jobs
 - Only batch jobs are allowed.
 - Interactive executions of jobs are not allowed.
- How to run
 - writing job script
 - submitting job
 - checking job status
 - checking results
- Utilization of computational resources
 - 1-node (48 cores) is occupied by each job.
 - Your node is not shared by other jobs.

Job Script

- `<$O-S1>/hello.sh`
- Scheduling + Shell Script

```
#!/bin/sh
#PJM -N "hello"
#PJM -L rscgrp=lecture8-o
#PJM -L node=1
#PJM -mpi proc=4
#PJM -L elapse=00:15:00
#PJM -g gt18
#PJM -j
#PJM -e err
#PJM -o hello.lst

module load fj
module load fjmpi

mpexec ./a.out
```

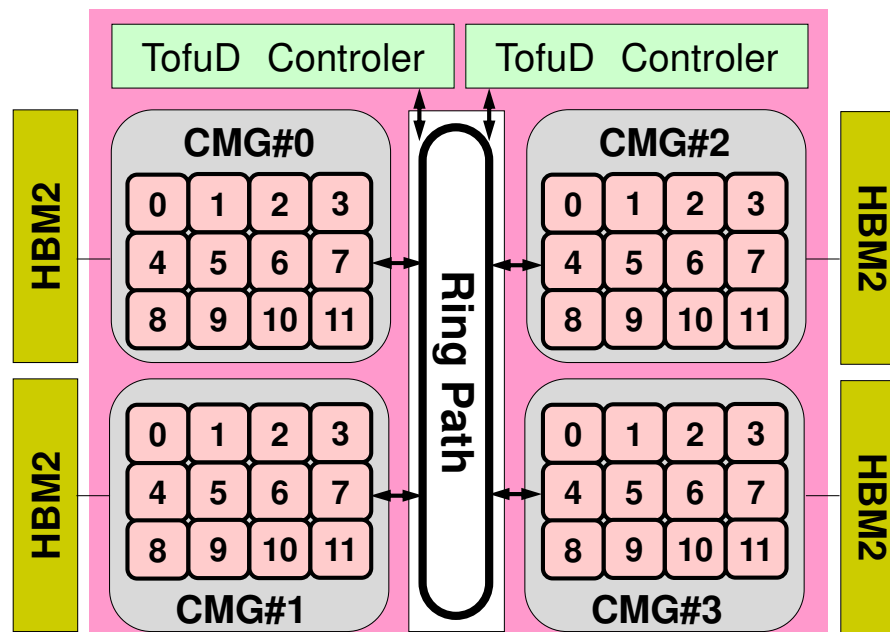
Job Name
Name of "Resource Group"
Node#
Total MPI Process#
Computation Time
Group Name (Wallet)
Standard Error
Standard Output

必須

Number of Processes

```
#PJM -L node=1; #PJM --mpi proc= 1      1-node, 1-proc, 1-proc/n
#PJM -L node=1; #PJM --mpi proc= 4      1-node, 4-proc, 4-proc/n
#PJM -L node=1; #PJM --mpi proc=12     1-node, 12-proc, 12-proc/n
#PJM -L node=1; #PJM --mpi proc=24     1-node, 24-proc, 24-proc/n
#PJM -L node=1; #PJM --mpi proc=48     1-node, 48-proc, 48-proc/n
```

```
#PJM -L node= 4; #PJM --mpi proc=192   4-node, 192-proc, 48-proc/n
#PJM -L node= 8; #PJM --mpi proc=384   8-node, 384-proc, 48-proc/n
#PJM -L node=12; #PJM --mpi proc=576  12-node, 576-proc, 48-proc/n
```



Job Submission

```
>$ cd /work/gt18/t18XXX/pFEM/mpi/S1
```

```
>$ module load fj
```

```
>$ pjsub hello.sh
```

```
>$ cat hello.lst
```

```
Hello World 0
```

```
Hello World 3
```

```
Hello World 2
```

```
Hello World 1
```

Available Resource Groups (Queue's)

- Following 2 resource groups are available
- Up to 12 nodes are available
 - **lecture-o**
 - 12 nodes (576 cores), 15 min., valid until October 15, 2023
 - Shared by all “educational” users
 - **lecture8-o**
 - 12 nodes (576 cores), 15 min., active during class time
 - More jobs (compared to **lecture-o**) can be processed up on availability.

Submitting & Checking Jobs

- Submitting Jobs `pjsub SCRIPT NAME`
- Checking status of jobs `pjstat`
- Deleting/aborting `pjdel JOB ID`
- Checking status of queues `pjstat --rsc`
- Detailed info. of queues `pjstat --rsc -x`
- Number of running jobs `pjstat -a`
- History of Submission `pjstat -H`
- Limitation of submission `pjstat --limit`

```
[t00470@wisteria01 run]$ pjsub f2_48.sh
[INFO] PJM 0000 pjsub Job 15713 submitted.
```

```
[t00470@wisteria01 run]$ pjsub f3_48.sh
[INFO] PJM 0000 pjsub Job 15714 submitted.
```

```
[t00470@wisteria01 run]$ pjstat
Wisteria/BDEC-01 scheduled stop time: 2021/05/28(Fri) 09:00:00 (Remain: 4days 1:25:56)
```

JOB_ID	JOB_NAME	STATUS	PROJECT	RSCGROUP	START_DATE	ELAPSE	TOKEN	NODE	GPU
15713	f2_48	RUNNING	gt00	lecture-o	05/24 07:34:03	00:00:02	-	1	-
15714	f3_48	QUEUED	gt00	lecture-o	--/-- --:--:--	00:00:00	-	1	-

```
[t00470@wisteria01 run]$ pjstat
Wisteria/BDEC-01 scheduled stop time: 2021/05/28(Fri) 09:00:00 (Remain: 4days 1:25:56)
```

JOB_ID	JOB_NAME	STATUS	PROJECT	RSCGROUP	START_DATE	ELAPSE	TOKEN	NODE	GPU
15713	f2_48	RUNNING	gt00	lecture-o	05/24 07:34:03	00:00:02	-	1	-
15714	f3_48	RUNNING	gt00	lecture-o	(05/24 07:34)	00:00:00	-	1	-

```
[t00XYZ@wisteria01 ~]$ pjdel 15714
[INFO] PJM 0100 pjdel Accepted Job 15714
```

```
[t00XYZ@wisteria01 ~]$ pjstat
Wisteria/BDEC-01 scheduled stop time: 2021/05/28(Fri) 09:00:00 (Remain: 4days 1:25:56)
```

JOB_ID	JOB_NAME	STATUS	PROJECT	RSCGROUP	START_DATE	ELAPSE	TOKEN	NODE	GPU
15713	f2_48	RUNNING	gt00	lecture-o	05/24 07:34:03	00:00:02	-	1	-

```
[t00XYZ@wisteria01 ~]$ pjstat
Wisteria/BDEC-01 scheduled stop time: 2021/05/28(Fri) 09:00:00 (Remain: 4days 1:21:45)
```

```
No unfinished job found.
```

```
[t00XYZ@wisteria01 ~]$ pjstat --rsc
```

```
SYSTEM: Odyssey
```

RSCGRP	STATUS	NODE
lecture-o	[ENABLE, START]	96
lecture8-o	[DISABLE, STOP]	2x12x16

```
[t00XYZ@wisteria01 ~]$ pjstat --rsc -x
```

```
SYSTEM: Odyssey
```

RSCGRP	STATUS	MIN_NODE	MAX_NODE	MAX_ELAPSE	REMAIN_ELAPSE	MEM(GiB)	PROJECT
lecture-o	[ENABLE, START]	1	12	00:15:00	00:15:00	28	gt00
lecture8-o	[DISABLE, STOP]	1	12	00:15:00	--:--:--	28	gt00

```
[t00XYZ@wisteria01 ~]$ pjstat --limit
```

```
SYSTEM: Odyssey
```

PROJECT	ACCEPT	RUN	BULK_ACCEPT	BULK_RUN	NODE
gt18	0/ 128	0/ 16	0/ 8	0/ 16	0/ 2304

```
SYSTEM: Aquarius
```

PROJECT	ACCEPT	RUN	BULK_ACCEPT	BULK_RUN	GPU
gt18	0/ 4	0/ 2	0/ 0	0/ 0	0/ 64

Basic/Essential Functions

```
implicit REAL*8 (A-H,O-Z)
include 'mpif.h'
integer :: PETOT, my_rank, ierr

call MPI_INIT      (ierr)
call MPI_COMM_SIZE (MPI_COMM_WORLD, PETOT, ierr )
call MPI_COMM_RANK (MPI_COMM_WORLD, my_rank, ierr )

write (*,'(a,2i8)') 'Hello World FORTRAN', my_rank, PETOT

call MPI_FINALIZE (ierr)

stop
end
```

```
#include "mpi.h"
#include <stdio.h>
int main(int argc, char **argv)
{
    int n, myid, numprocs, i;

    MPI_Init (&argc, &argv);
    MPI_Comm_size (MPI_COMM_WORLD, &numprocs);
    MPI_Comm_rank (MPI_COMM_WORLD, &myid);

    printf ("Hello World %d¥n", myid);
    MPI_Finalize ();
}
```

'mpif.h', "mpi.h"
Essential Include file
"use mpi" is possible in F90

MPI_Init
Initialization

MPI_Comm_size
Number of MPI Processes
mpirun -np XX <prog>

MPI_Comm_rank
Process ID starting from 0

MPI_Finalize
Termination of MPI processes

Difference between FORTRAN/C

- (Basically) same interface
 - In C, UPPER/lower cases are considered as different
 - e.g.: **MPI_Comm_size**
 - MPI: UPPER case
 - First character of the function except “MPI_” is in UPPER case.
 - Other characters are in lower case.
- In Fortran, return value `ierr` has to be added at the end of the argument list.
- C needs special types for variables:
 - `MPI_Comm`, `MPI_Datatype`, `MPI_Op` etc.
- **MPI_INIT** is different:
 - call `MPI_INIT (ierr)`
 - `MPI_Init (int *argc, char ***argv)`

What's are going on ?

```
#include "mpi.h"
#include <stdio.h>
int main(int argc, char **argv)
{
    int n, myid, numprocs, i;
    MPI_Init (&argc, &argv);
    MPI_Comm_size (MPI_COMM_WORLD, &numprocs);
    MPI_Comm_rank (MPI_COMM_WORLD, &myid);
    printf ("Hello World %d¥n", myid);
    MPI_Finalize ();
}
```

```
#!/bin/sh
#PJM -N "hello"
#PJM -L rscgrp=lecture7-o
#PJM -L node=1
#PJM -mpi proc=4
#PJM -L elapse=00:15:00
#PJM -g gt87
#PJM -j
#PJM -e err
#PJM -o hello.lst

module load fj
module load fjmpi

mpiexec ./a.out
```

Job Name	
Name of "Resource Group"	
Node#	
Total MPI Process#	
Computation Time	
Group Name (Wallet)	
Standard Error	
Standard Output	
必須	

- **mpiexec** starts up 4 MPI processes ("proc=4")
 - A single program runs on four processes.
 - each process writes a value of `myid`
- Four processes do same operations, but values of `myid` are different.
- Output of each process is different.
- That is SPMD !

mpi.h, mpif.h

```
implicit REAL*8 (A-H,O-Z)
include 'mpif.h'
integer :: PETOT, my_rank, ierr

call MPI_INIT      (ierr)
call MPI_COMM_SIZE (MPI_COMM_WORLD, PETOT, ierr )
call MPI_COMM_RANK (MPI_COMM_WORLD, my_rank, ierr )

write (*,'(a,2i8)') 'Hello World FORTRAN', my_rank, PETOT

call MPI_FINALIZE (ierr)

stop
end
```

```
#include "mpi.h"
#include <stdio.h>
int main(int argc, char **argv)
{
    int n, myid, numprocs, i;

    MPI_Init (&argc, &argv);
    MPI_Comm_size (MPI_COMM_WORLD, &numprocs);
    MPI_Comm_rank (MPI_COMM_WORLD, &myid);

    printf ("Hello World %d\n", myid);
    MPI_Finalize ();
}
```

- Various types of parameters and variables for MPI & their initial values.
- Name of each var. starts from “MPI_”
- Values of these parameters and variables cannot be changed by users.
- Users do not specify variables starting from “MPI_” in users’ programs.

MPI_Init

- Initialize the MPI execution environment (required)
- It is recommended to put this BEFORE all statements in the program.
- **MPI_Init (argc, argv)**

```
#include "mpi.h"
#include <stdio.h>
int main(int argc, char **argv)
{
    int n, myid, numprocs, i;

    MPI_Init (&argc, &argv);
    MPI_Comm_size (MPI_COMM_WORLD, &numprocs);
    MPI_Comm_rank (MPI_COMM_WORLD, &myid);

    printf ("Hello World %d\n", myid);
    MPI_Finalize ();
}
```

MPI_Finalize

- Terminates MPI execution environment (required)
- It is recommended to put this AFTER all statements in the program.
- **Please do not forget this.**
- **MPI_Finalize ()**

```
#include "mpi.h"
#include <stdio.h>
int main(int argc, char **argv)
{
    int n, myid, numprocs, i;

    MPI_Init (&argc, &argv);
    MPI_Comm_size (MPI_COMM_WORLD, &numprocs);
    MPI_Comm_rank (MPI_COMM_WORLD, &myid);

    printf ("Hello World %d\n", myid);
    MPI_Finalize ();
}
```

MPI_Comm_size

- Determines the size of the group associated with a communicator
- not required, but very convenient function
- **MPI_Comm_size (comm, size)**
 - **comm** MPI_Comm I communicator
 - **size** int O number of processes in the group of communicator

```
#include "mpi.h"
#include <stdio.h>
int main(int argc, char **argv)
{
    int n, myid, numprocs, i;

    MPI_Init (&argc, &argv);
    MPI_Comm_size (MPI_COMM_WORLD, &numprocs);
    MPI_Comm_rank (MPI_COMM_WORLD, &myid);

    printf ("Hello World %d\n", myid);
    MPI_Finalize ();
}
```

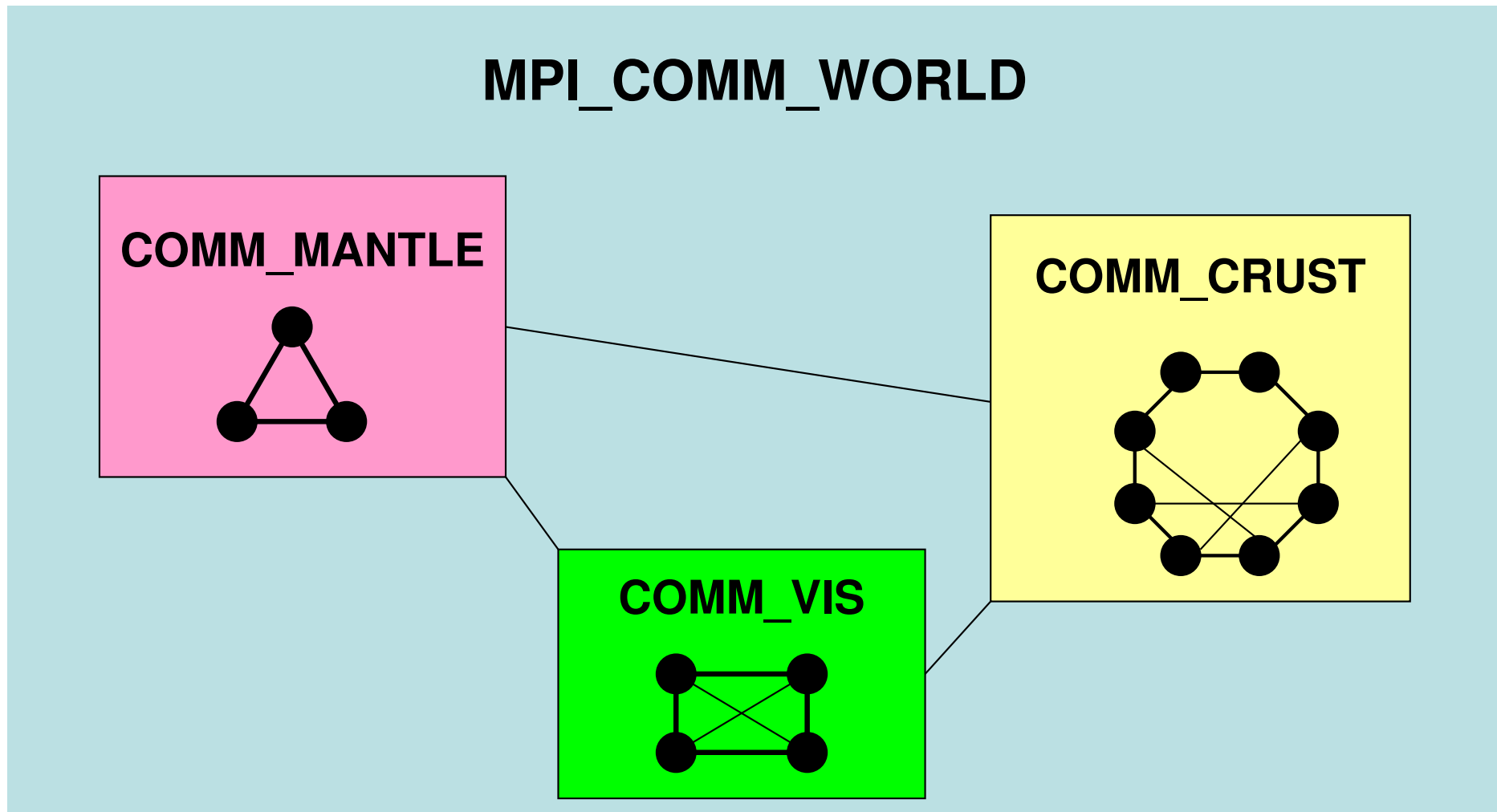
What is Communicator ?

```
MPI_Comm_Size (MPI_COMM_WORLD, PETOT)
```

- Group of processes for communication
- Communicator must be specified in MPI program as a unit of communication
- All processes belong to a group, named “**MPI_COMM_WORLD**” (default)
- Multiple communicators can be created, and complicated operations are possible.
 - Computation, Visualization
- Only “**MPI_COMM_WORLD**” is needed in this class.

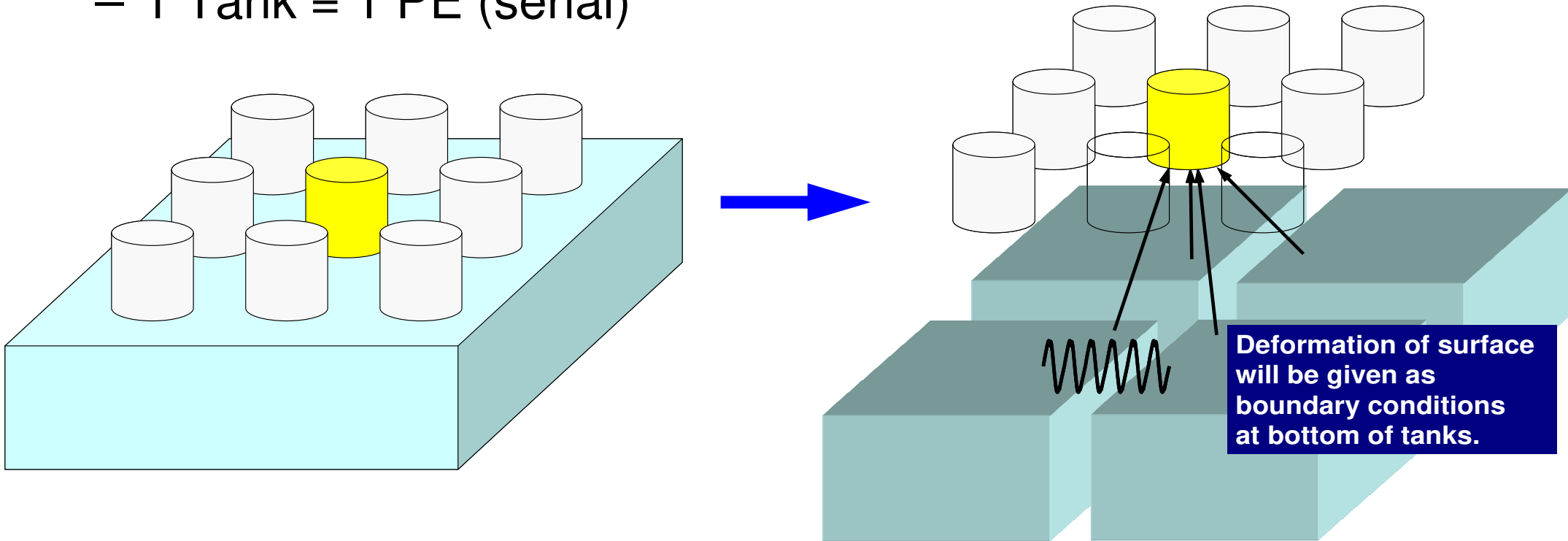
Communicator in MPI

One process can belong to multiple communicators



Target Application

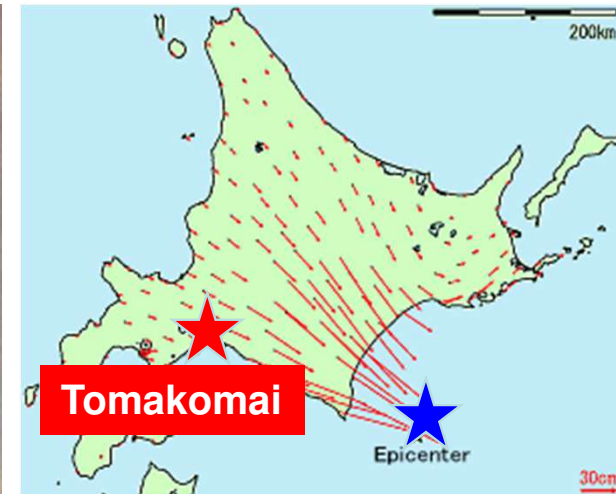
- Coupling between “Ground Motion” and “Sloshing of Tanks for Oil-Storage”
 - “One-way” coupling from “Ground Motion” to “Tanks”.
 - Displacement of ground surface is given as forced displacement of bottom surface of tanks.
 - 1 Tank = 1 PE (serial)



2003 Tokachi Earthquake (M8.0)

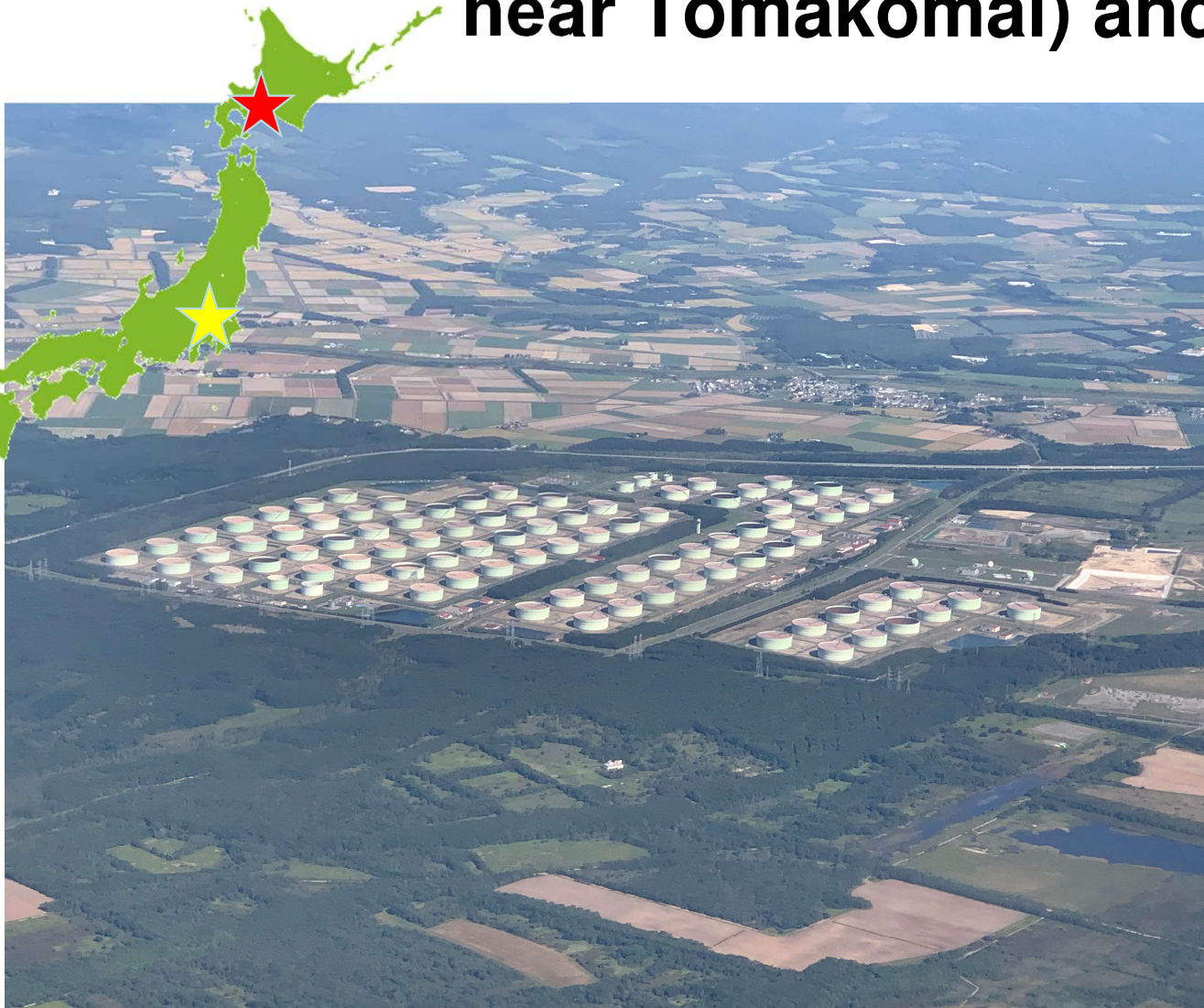
Fire accident by sloshing due to long-period ground motion

- Oil tanks in Tomakomai shook violently by the earthquake
- Sparks with rubbed metal fittings ignited the oil that swung on the liquid level
- The fire was so large that it was impossible to extinguish it
- We needed to wait a few days for all the oil to burn out.



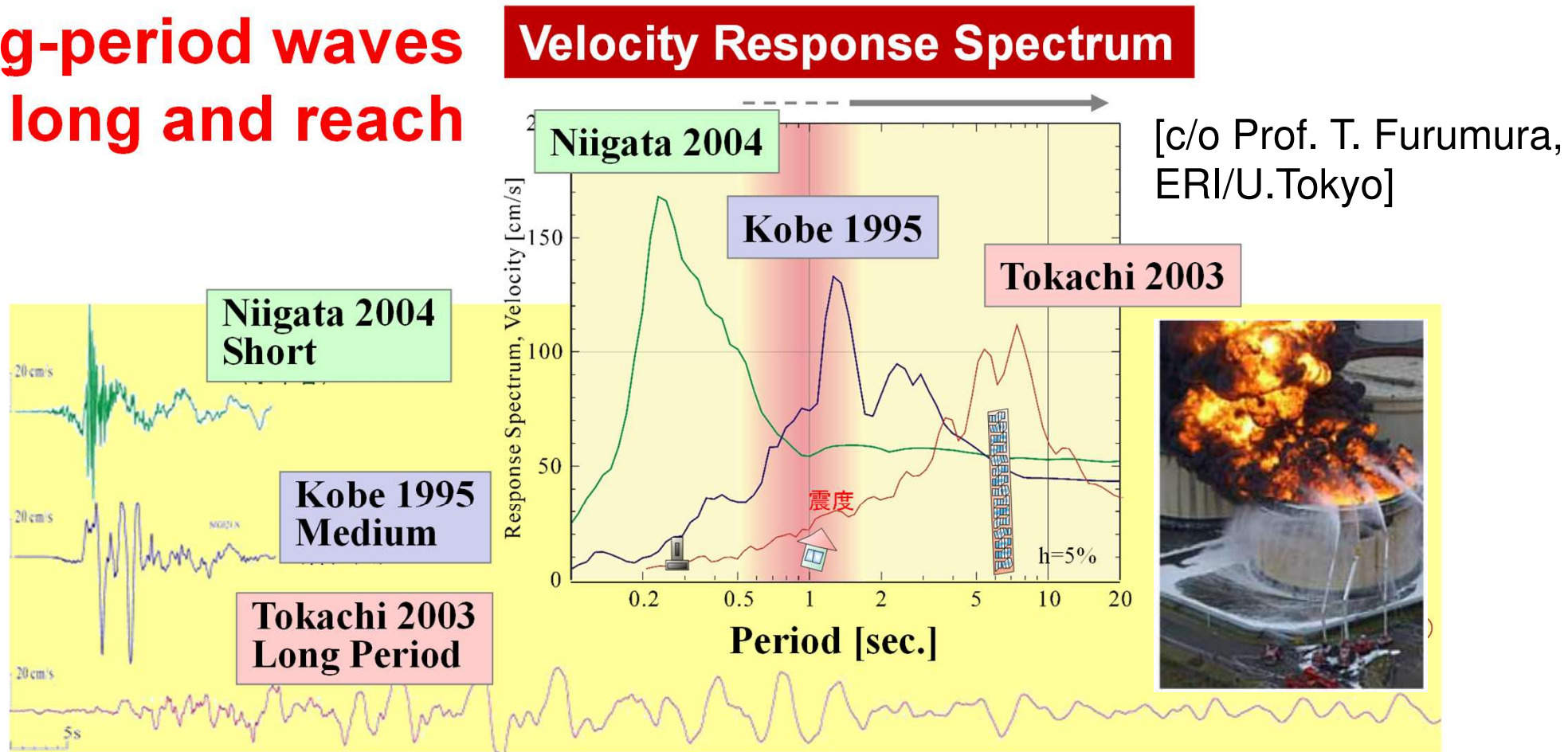
The Tanks on September 10, 2022

View from flight between Sapporo (Capital of Hokkaido, near Tomakomai) and Tokyo

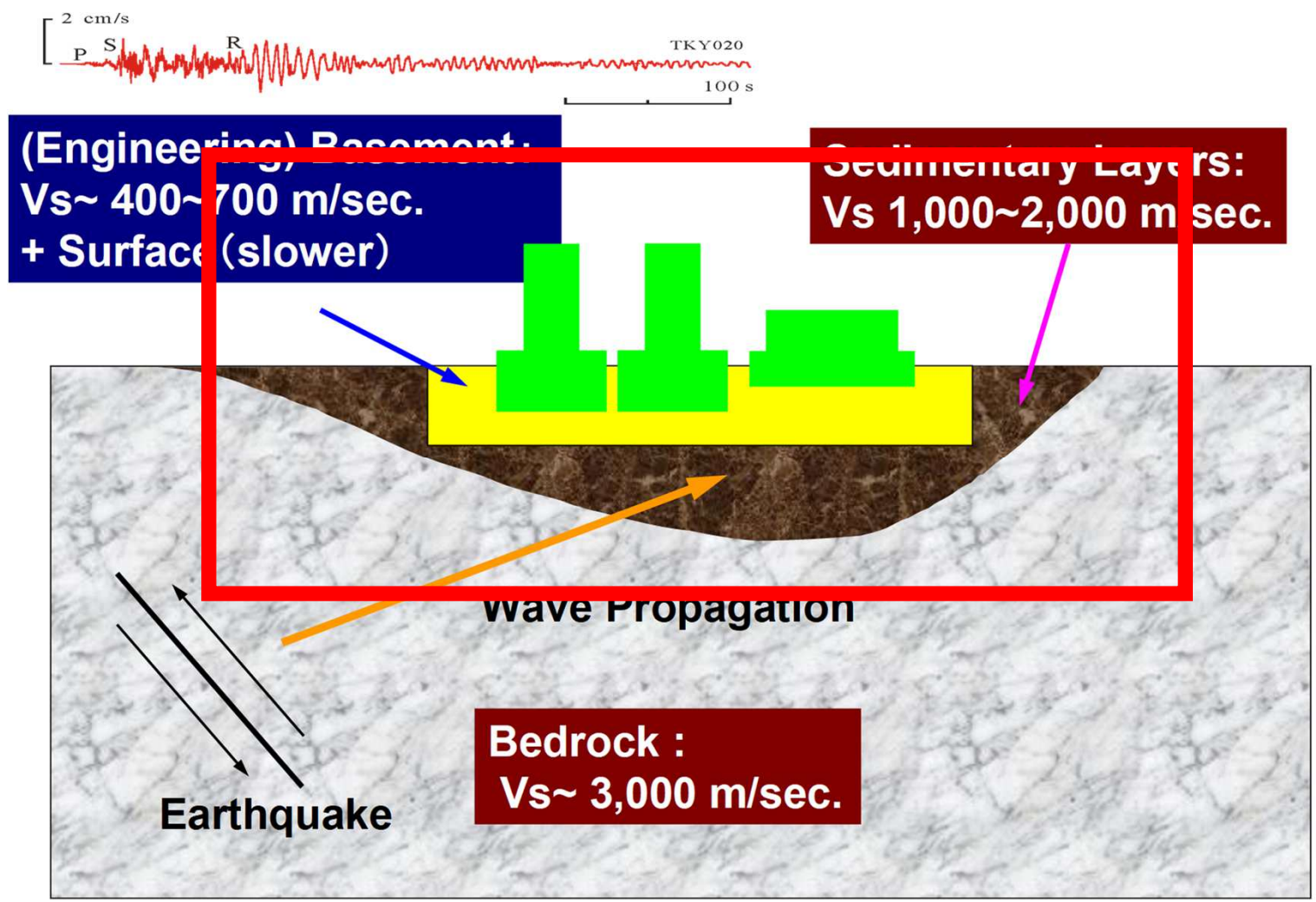


Seismic Wave

- Various Components of Wavelength
- Buildings with the same natural period as the predominant component of seismic waves shake most violently (0.1-10 sec.): a kind of "resonance"
- **Long-period waves last long and reach far**

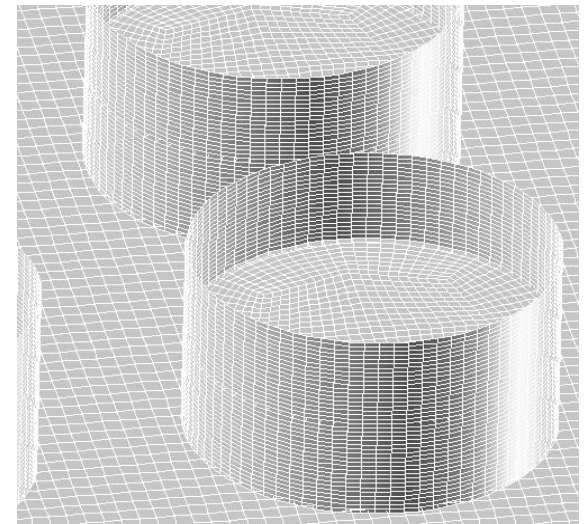
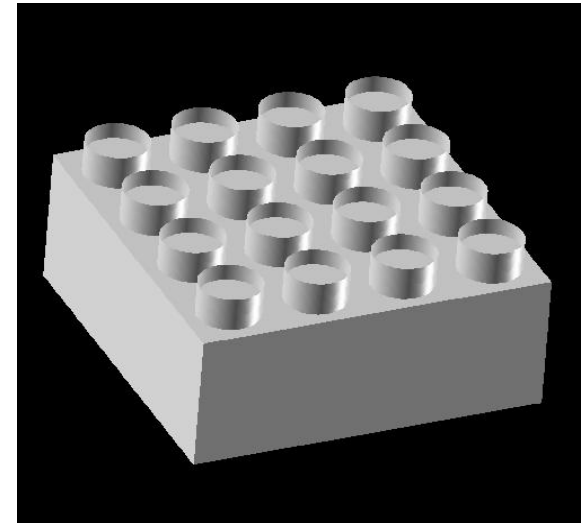


Seismic Wave Propagation, Underground Structure

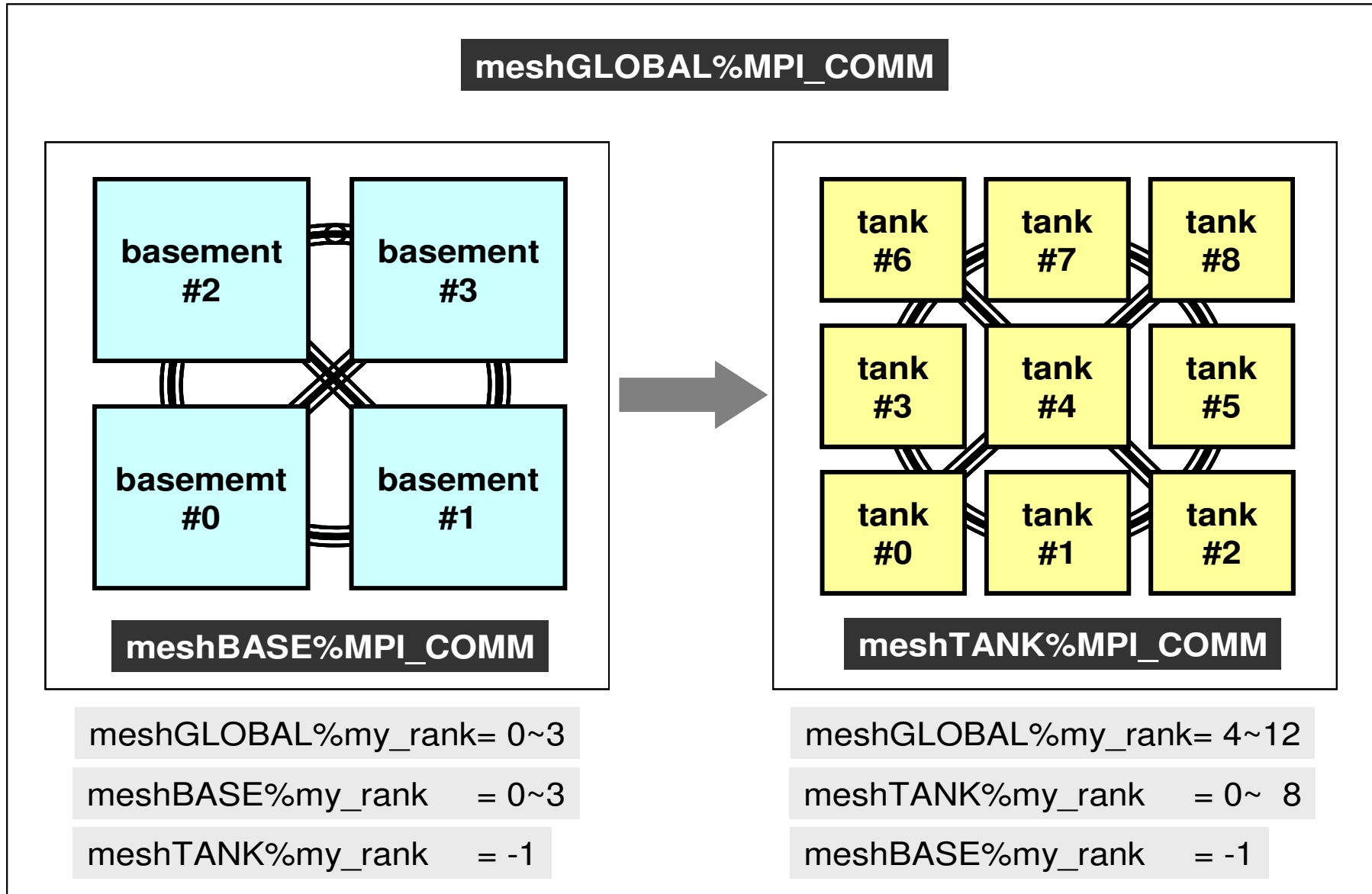


Simulation Codes

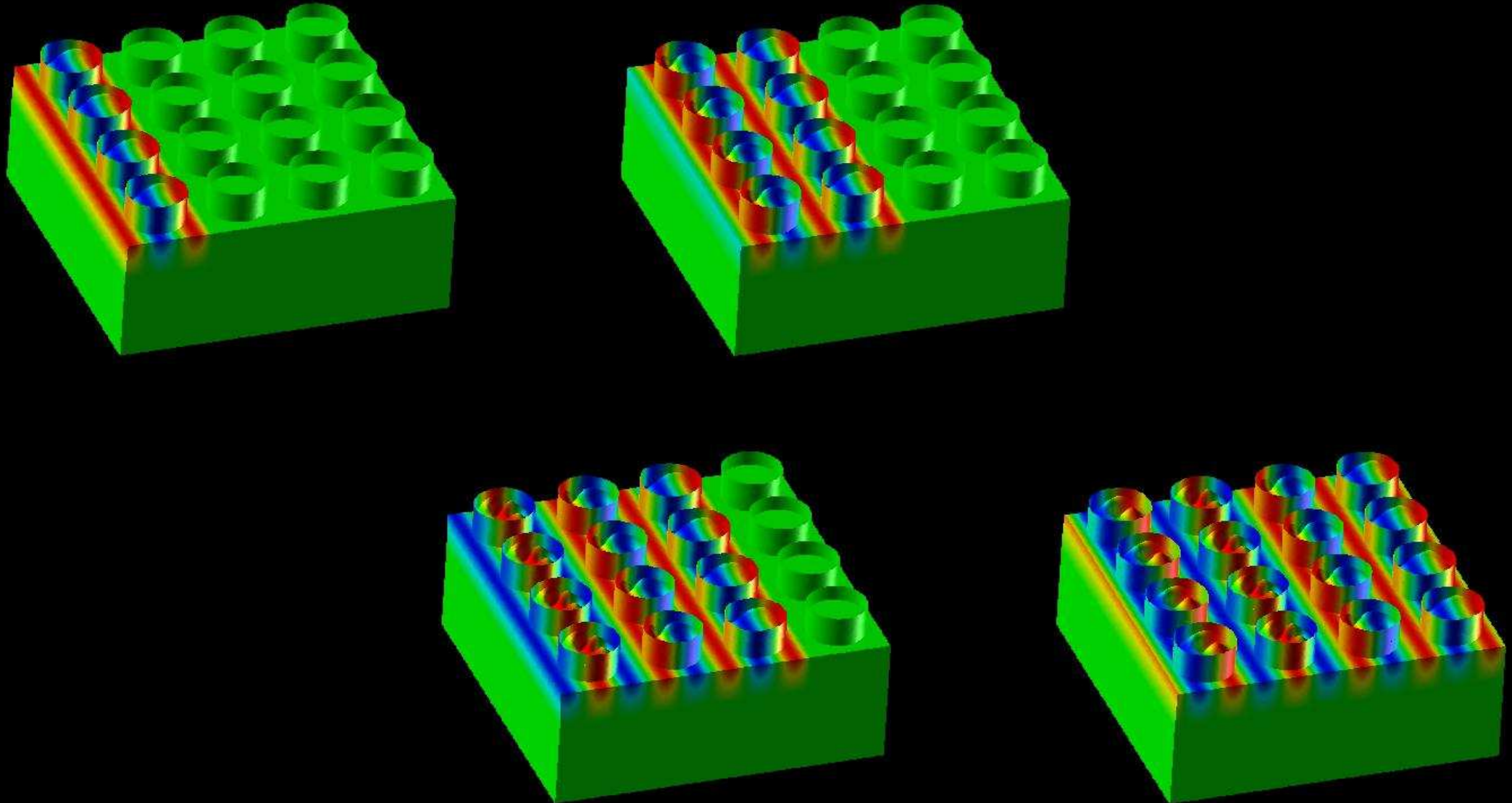
- Ground Motion (Ichimura): Fortran
 - Parallel FEM, 3D Elastic/Dynamic
 - Explicit forward Euler scheme
 - Each element: $2\text{m} \times 2\text{m} \times 2\text{m}$ cube
 - $240\text{m} \times 240\text{m} \times 100\text{m}$ region
- Sloshing of Tanks (Nagashima): C
 - Serial FEM (Embarrassingly Parallel)
 - Implicit backward Euler, Skyline method
 - Shell elements + Inviscid potential flow
 - D: 42.7m, H: 24.9m, T: 20mm,
 - Frequency: 7.6sec.
 - 80 elements in circ., 0.6m mesh in height
 - Tank-to-Tank: 60m, 4×4
- Total number of unknowns: 2,918,169

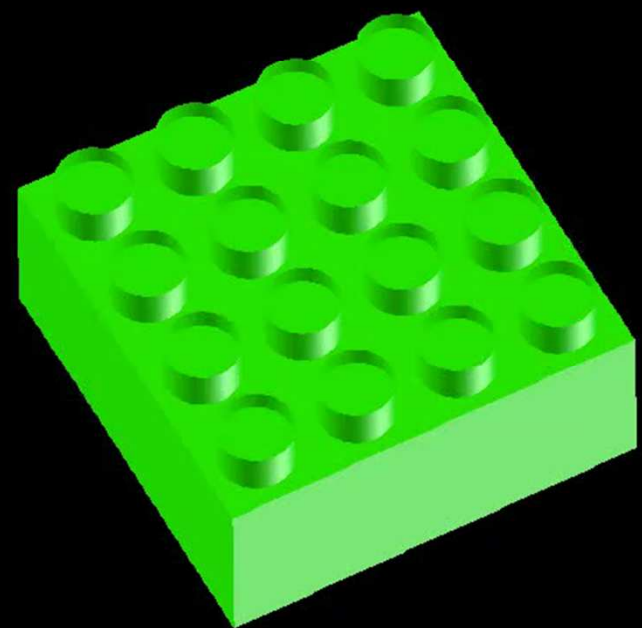
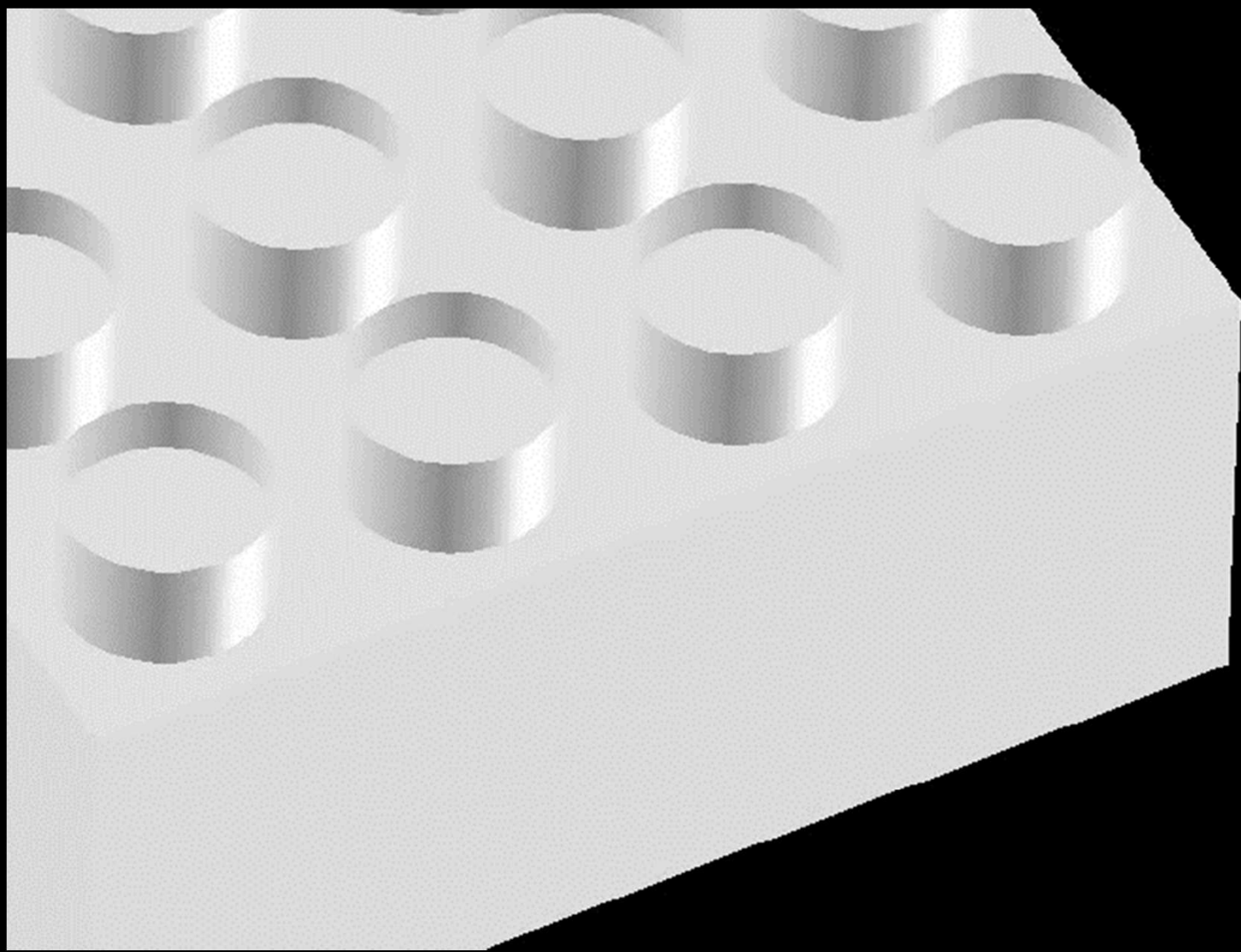


Three Communicators



Coupling between “Ground Motion” and “Sloshing of Tanks for Oil-Storage”





MPI_Comm_rank

- Determines the rank of the calling process in the communicator
 - “ID of MPI process” is sometimes called “rank”
- **MPI_Comm_rank (comm, rank)**
 - **comm** MPI_Comm I communicator
 - **rank** int 0 rank of the calling process in the group of comm

Starting from “0”

```
#include "mpi.h"
#include <stdio.h>
int main(int argc, char **argv)
{
    int n, myid, numprocs, i;

    MPI_Init (&argc, &argv);
    MPI_Comm_size (MPI_COMM_WORLD, &numprocs);
    MPI_Comm_rank (MPI_COMM_WORLD, &myid);

    printf ("Hello World %d\n", myid);
    MPI_Finalize ();
}
```



MPI_Abort

- Aborts MPI execution environment
- **MPI_Abort (comm, errcode)**
 - comm MPI_Comm I communicator
 - errcode int 0 error code

MPI_Wtime

- Returns an elapsed time on the calling processor

- **time= MPI_Wtime ()**

– time double 0

Time in seconds since an arbitrary time in the past.

```
...  
double Stime, Etime;  
  
Stime= MPI_Wtime ();  
  
(...)  
  
Etime= MPI_Wtime ();
```

Example of MPI_Wtime

```
$> cd /work/gt18/t18XXX/pFEM/mpi/S1
```

```
$> module load fj
```

```
$> mpifccpx -Nclang -O1 time.c
```

```
$> mpifrtpx -O1 time.f
```

(modify go4.sh, 4 processes)

```
$> pjsub go4.sh
```

```
0      1.113281E+00
3      1.113281E+00
2      1.117188E+00
1      1.117188E+00
```

**Process
ID**

Time

go4.sh

```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture8-o
#PJM -L node=1
#PJM --mpi proc=4
#PJM -L elapse=00:15:00
#PJM -g gt18
#PJM -j
#PJM -e err
#PJM -o test.lst

module load fj
module load fjmpi

mpiexec ./a.out
```

MPI_Wtick

- Returns the resolution of MPI_Wtime
- **depends on hardware, and compiler**

- **time= MPI_Wtick ()**

– time double 0

Time in seconds of resolution of MPI_Wtime

```
implicit REAL*8 (A-H,O-Z)
include 'mpif.h'
```

```
...
TM= MPI_WTICK ()
write (*,*) TM
...
```

```
double Time;
```

```
...
Time = MPI_Wtick();
printf("%5d%16.6E\n", MyRank, Time);
...
```

Example of MPI_Wtick

```
$> cd /work/gt18/t18XXX/pFEM/mpi/S1
```

```
$> module load fj
```

```
$> mpifccpx -Nclang -O1 wtick.c
```

```
$> mpifrtpx -O1 wtick.f
```

```
(modify go1.sh, 1 process)
```

```
$> pjsub go1.sh
```


go1.sh

```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture8-o
#PJM -L node=1
#PJM --mpi proc=1
#PJM -L elapse=00:15:00
#PJM -g gt18
#PJM -j
#PJM -e err
#PJM -o test.lst

module load fj
module load fjmpi

mpiexec ./a.out
```



MPI_Barrier

- Blocks until all processes in the communicator have reached this routine.
- Mainly for debugging, huge overhead, not recommended for real code.
- **MPI_Barrier (comm)**
 - comm MPI_Comm I communicator

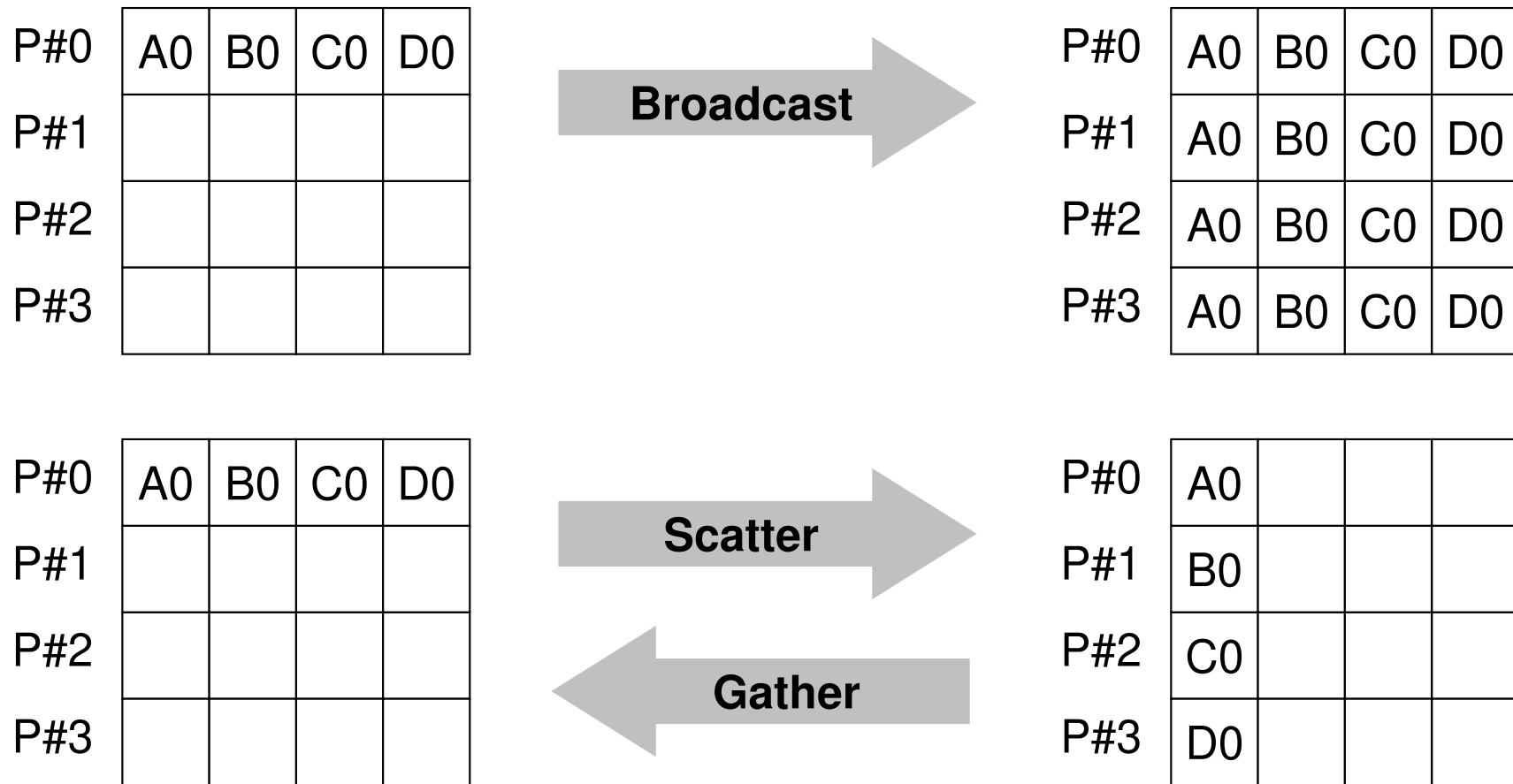
- What is MPI ?
- Your First MPI Program: Hello World
- **Collective Communication**
- Point-to-Point Communication

What is Collective Communication ?

集団通信, グループ通信

- Collective communication is the process of exchanging information between multiple MPI processes in the communicator: one-to-all or all-to-all communications.
- Examples
 - Broadcasting control data
 - Max, Min
 - Summation
 - Dot products of vectors
 - Transformation of dense matrices

Example of Collective Communications (1/4)



Example of Collective Communications (2/4)

P#0	A0			
P#1	B0			
P#2	C0			
P#3	D0			

All gather

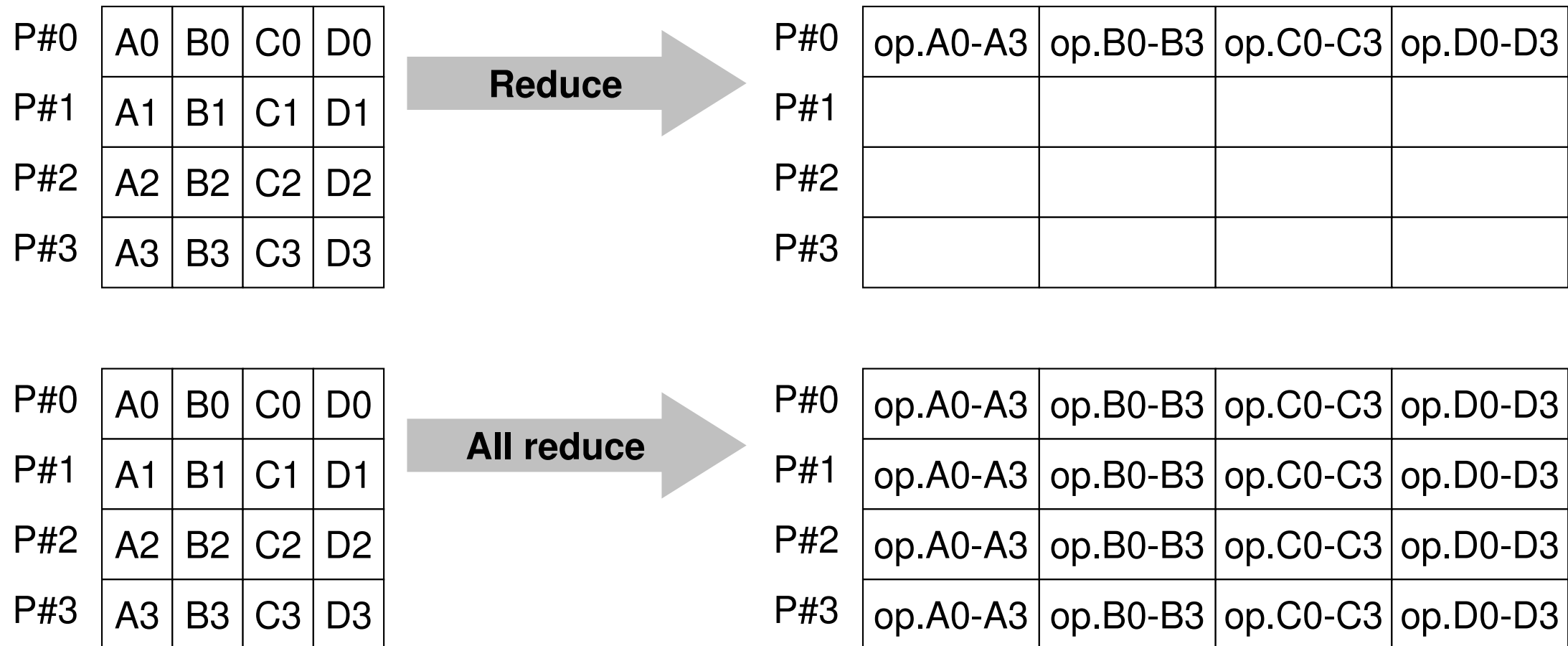
P#0	A0	B0	C0	D0
P#1	A0	B0	C0	D0
P#2	A0	B0	C0	D0
P#3	A0	B0	C0	D0

P#0	A0	A1	A2	A3
P#1	B0	B1	B2	B3
P#2	C0	C1	C2	C3
P#3	D0	D1	D2	D3

All-to-All

P#0	A0	B0	C0	D0
P#1	A1	B1	C1	D1
P#2	A2	B2	C2	D2
P#3	A3	B3	C3	D3

Example of Collective Communications (3/4)



Example of Collective Communications (4/4)

P#0	A0	B0	C0	D0
P#1	A1	B1	C1	D1
P#2	A2	B2	C2	D2
P#3	A3	B3	C3	D3

Reduce scatter



P#0	op.A0-A3			
P#1	op.B0-B3			
P#2	op.C0-C3			
P#3	op.D0-D3			

Examples by Collective Comm.

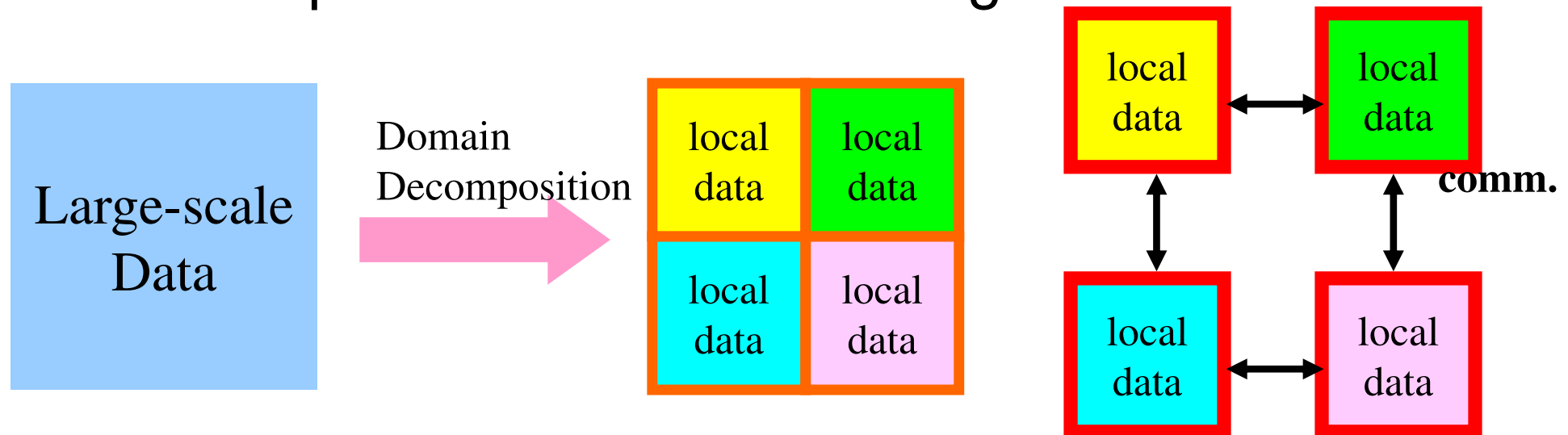
- **Dot Products of Vectors**
- Scatter/Gather
- Reading Distributed Files
- MPI_Allgatherv

Global/Local Data

- Data structure of parallel computing based on SPMD, where large scale “global data” is decomposed to small pieces of “local data”.

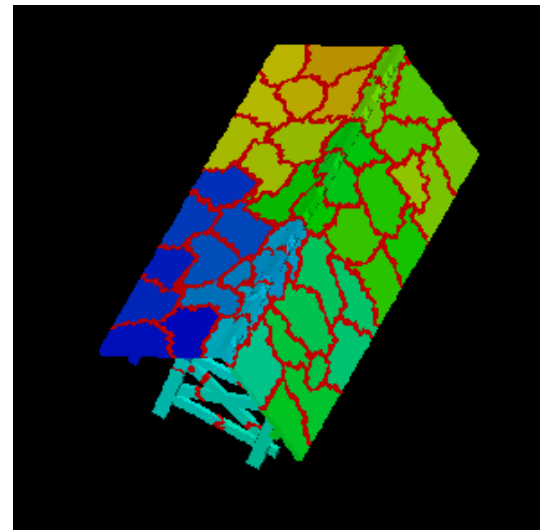
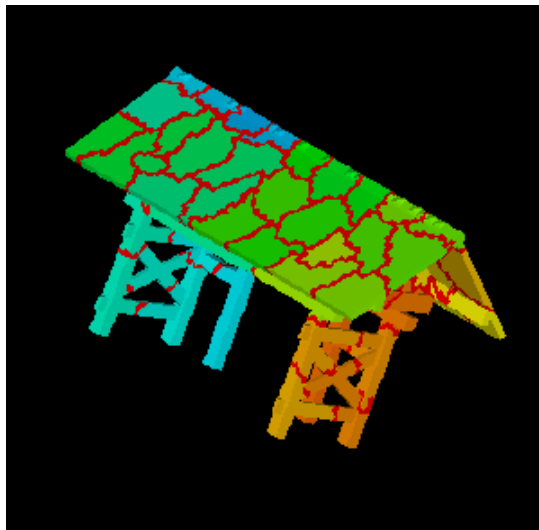
Domain Decomposition/Partitioning

- PC with 1GB RAM: can execute FEM application with up to 10^6 meshes
 - $10^3\text{km} \times 10^3\text{ km} \times 10^2\text{ km}$ (SW Japan): 10^8 meshes by 1km cubes
- Large-scale Data: Domain decomposition, parallel & local operations
- Global Computation: Comm. among domains needed



Local Data Structure

- It is important to define proper local data structure for target computation (and its algorithm)
 - Algorithms= Data Structures
- Main objective of this class !



Global/Local Data

- Data structure of parallel computing based on SPMD, where large scale “global data” is decomposed to small pieces of “local data”.
- Consider the dot product of following VEC_p and VEC_s with length=20 by parallel computation using 4 processors

```
VECp [ 0 ] =  2  
        [ 1 ] =  2  
        [ 2 ] =  2  
...  
        [17 ] =  2  
        [18 ] =  2  
        [19 ] =  2
```

```
VECs [ 0 ] =  3  
        [ 1 ] =  3  
        [ 2 ] =  3  
...  
        [17 ] =  3  
        [18 ] =  3  
        [19 ] =  3
```

<\$O-S1>/dot.f, dot.c

```
implicit REAL*8 (A-H,O-Z)
real(kind=8),dimension(20):: &
    VECp,  VECs

do i= 1, 20
    VECp(i)= 2.0d0
    VECs(i)= 3.0d0
enddo

sum= 0.d0
do ii= 1, 20
    sum= sum + VECp(ii)*VECs(ii)
enddo

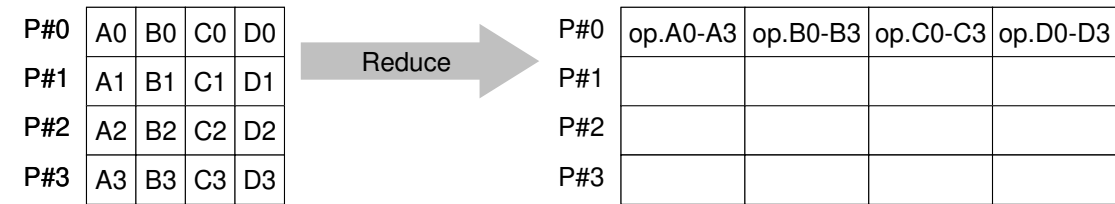
stop
end
```

```
#include <stdio.h>
int main(){
    int i;
    double VECp[20], VECs[20]
    double sum;

    for(i=0;i<20;i++){
        VECp[i]= 2.0;
        VECs[i]= 3.0;
    }

    sum = 0.0;
    for(i=0;i<20;i++){
        sum += VECp[i] * VECs[i];
    }
    return 0;
}
```

MPI_Reduce



- Reduces values on all processes to a single value
 - Summation, Product, Max, Min etc.

- **MPI_Reduce (sendbuf, recvbuf, count, datatype, op, root, comm)**
 - **sendbuf** choice I starting address of send buffer
 - **recvbuf** choice O starting address receive buffer
 - type is defined by "datatype"**
 - **count** int I number of elements in send/receive buffer
 - **datatype** MPI_Datatype I data type of elements of send/recive buffer
 - FORTRAN MPI_INTEGER, MPI_REAL, MPI_DOUBLE_PRECISION, MPI_CHARACTER etc.
 - C MPI_INT, MPI_FLOAT, MPI_DOUBLE, MPI_CHAR etc
 - **op** MPI_Op I reduce operation
 - MPI_MAX, MPI_MIN, MPI_SUM, MPI_PROD, MPI_LAND, MPI_BAND etc
 - Users can define operations by MPI_OP_CREATE**
 - **root** int I rank of root process
 - **comm** MPI_Comm I communicator

Send/Receive Buffer (Sending/Receiving)

- Arrays of “send (sending) buffer” and “receive (receiving) buffer” often appear in MPI.
- Addresses of “send (sending) buffer” and “receive (receiving) buffer” must be different.

Send/Receive Buffer (1/3)

A: Scalar

```
call MPI_REDUCE  
(A, recvbuf, 1, datatype, op, root, comm, ierr)
```

```
MPI_Reduce  
(A, recvbuf, 1, datatype, op, root, comm)
```

Send/Receive Buffer (2/3)

A: Array

```
call MPI_REDUCE
(A, recvbuf, 3, datatype, op, root, comm, ierr)
```

```
MPI_Reduce
(A, recvbuf, 3, datatype, op, root, comm)
```

- Starting Address of Send Buffer
 - $A(1)$: Fortran, $A[0]$: C
 - 3 (continuous) components of A ($A(1) - A(3)$, $A[0] - A[2]$) are sent

A(:)	1	2	3	4	5	6	7	8	9	10
-------------	---	---	---	---	---	---	---	---	---	----

A[:]	0	1	2	3	4	5	6	7	8	9
-------------	---	---	---	---	---	---	---	---	---	---

Send/Receive Buffer (3/3)

A: Array

```
call MPI_REDUCE
(A(4), recvbuf, 3, datatype, op, root, comm, ierr)
```

```
MPI_Reduce
(A[3], recvbuf, 3, datatype, op, root, comm)
```

- Starting Address of Send Buffer
 - $A(4)$: Fortran, $A[3]$: C
 - 3 (continuous) components of A ($A(4) - A(6)$, $A[3] - A[5]$) are sent

A(:)	1	2	3	4	5	6	7	8	9	10
A[:]	0	1	2	3	4	5	6	7	8	9

Example of MPI_Reduce (1/2)

MPI_Reduce

(sendbuf, recvbuf, count, datatype, op, root, comm)

```
double X0, X1;
```

```
MPI_Reduce
```

```
(&X0, &X1, 1, MPI_DOUBLE, MPI_MAX, 0, <comm>);
```

```
double X0[4], XMAX[4];
```

```
MPI_Reduce
```

```
(X0, XMAX, 4, MPI_DOUBLE, MPI_MAX, 0, <comm>);
```

Global Max values of X0[i] go to XMAX[i] on #0 process (i=0~3)

Example of MPI_Reduce (2/2)

MPI_Reduce

(sendbuf, recvbuf, count, datatype, op, root, comm)

```
double X0, XSUM;
```

```
MPI_Reduce
```

```
(&X0, &XSUM, 1, MPI_DOUBLE, MPI_SUM, 0, <comm>)
```

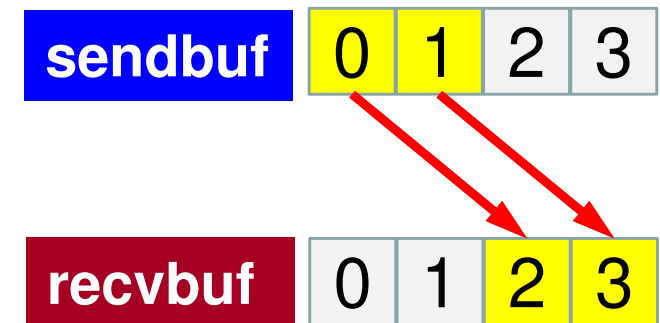
Global summation of X0 goes to XSUM on #0 process.

```
double X0[4];
```

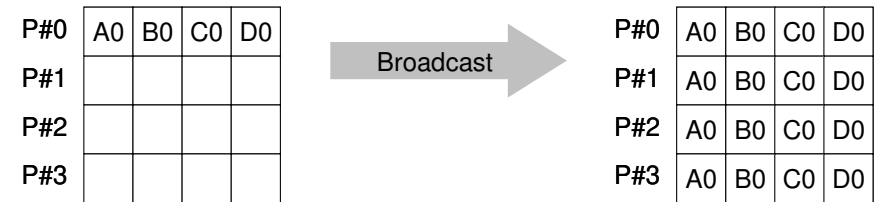
```
MPI_Reduce
```

```
(&X0[0], &X0[2], 2, MPI_DOUBLE_PRECISION, MPI_SUM, 0, <comm>)
```

- Global summation of X0[0] goes to X0[2] on #0 process.
- Global summation of X0[1] goes to X0[3] on #0 process.

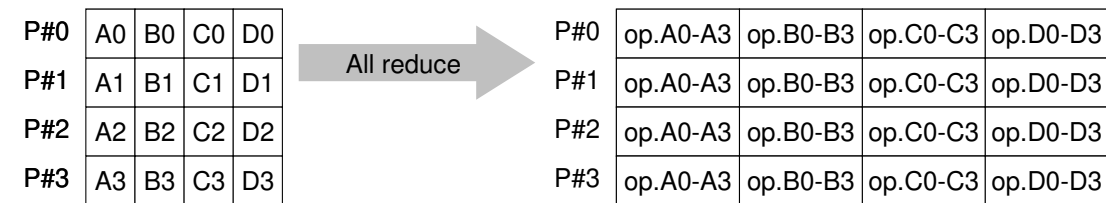


MPI_Bcast



- Broadcasts a message from the process with rank "root" to all other processes of the communicator
- **MPI_Bcast (buffer, count, datatype, root, comm)**
 - **buffer** choice I/O starting address of buffer
type is defined by "**datatype**"
 - **count** int I number of elements in send/recv buffer
 - **datatype** MPI_Datatype I data type of elements of send/recv buffer
 - FORTTRAN MPI_INTEGER, MPI_REAL, MPI_DOUBLE_PRECISION, MPI_CHARACTER etc.
 - C MPI_INT, MPI_FLOAT, MPI_DOUBLE, MPI_CHAR etc.
 - **root** int I rank of root process
 - **comm** MPI_Comm I communicator

MPI_Allreduce



- MPI_Reduce + MPI_Bcast
- Summation (of dot products) and MAX/MIN values are likely to be utilized in each process

- call MPI_Allreduce

(**sendbuf**, **recvbuf**, **count**, **datatype**, **op**, **comm**)

- **sendbuf** choice I starting address of send buffer
- **recvbuf** choice O starting address receive buffer
type is defined by "**datatype**"
- **count** int I number of elements in send/recv buffer
- **datatype** MPI_Datatype I data type of elements of send/recv buffer
- **op** MPI_Op I reduce operation
- **comm** MPI_Comm I communicator

“op” of MPI_Reduce/Allreduce

MPI_Reduce

(sendbuf, recvbuf, count, datatype, op, root, comm)

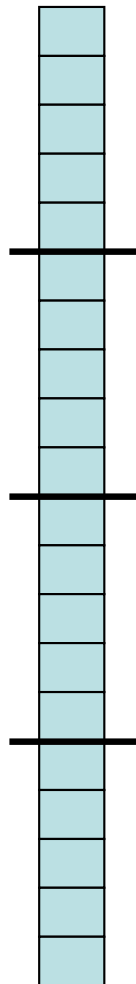
- `MPI_MAX`, `MPI_MIN` Max, Min
- `MPI_SUM`, `MPI_PROD` Summation, Product
- `MPI_LAND` Logical AND

Local Data (1/2)

- Decompose vector with length=20 into 4 domains (processes)
- Each process handles a vector with length= 5

```
VECp [ 0] = 2  
     [ 1] = 2  
     [ 2] = 2  
     ...  
     [17] = 2  
     [18] = 2  
     [19] = 2
```

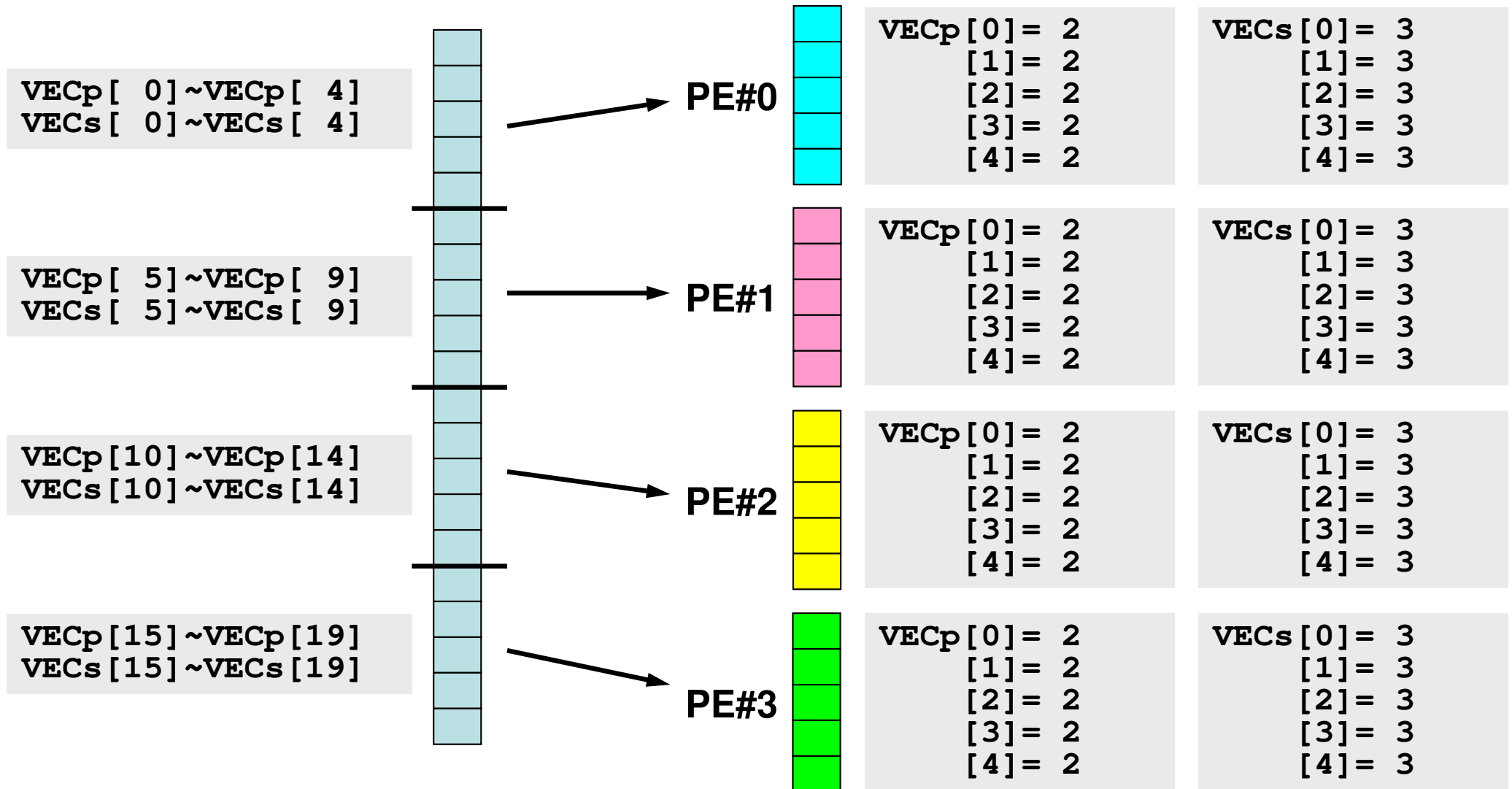
```
VECs [ 0] = 3  
     [ 1] = 3  
     [ 2] = 3  
     ...  
     [17] = 3  
     [18] = 3  
     [19] = 3
```





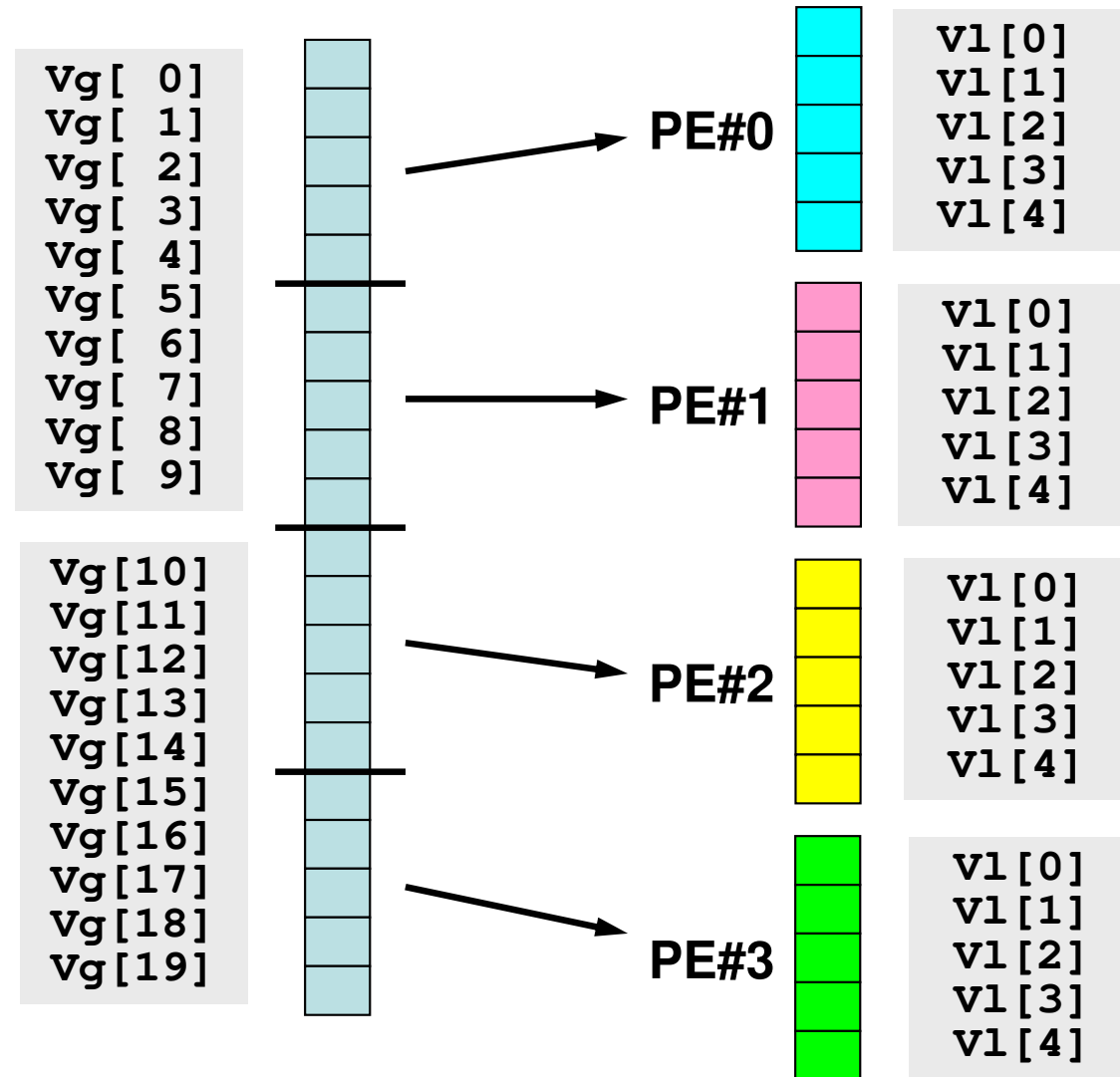
Local Data (2/2)

- 1th-5th components of original global vector go to 1th-5th components of PE#0, 6th-10th -> PE#1, 11th-15th -> PE#2, 16th-20th -> PE#3.



But ...

- It is too easy !! Just decomposing and renumbering from 1 (or 0).
- Of course, this is not enough. Further examples will be shown in the latter part.



Example: Dot Product (1/3)

<\$O-S1>/allreduce.c

```
#include <stdio.h>
#include <stdlib.h>
#include "mpi.h"

int main(int argc, char **argv){
    int i,N;
    int PeTot, MyRank;
    double VECp[5], VECs[5];
    double sumA, sumR, sum0;

    MPI_Init(&argc, &argv);
    MPI_Comm_size(MPI_COMM_WORLD, &PeTot);
    MPI_Comm_rank(MPI_COMM_WORLD, &MyRank);

    sumA= 0.0;
    sumR= 0.0;

    N=5;
    for(i=0;i<N;i++){
        VECp[i] = 2.0;
        VECs[i] = 3.0;
    }

    sum0 = 0.0;
    for(i=0;i<N;i++){
        sum0 += VECp[i] * VECs[i];
    }
}
```

Local vector is generated
at each local process.

Example: Dot Product (2/3)

<\$0-S1>/allreduce.c

```
MPI_Reduce(&sum0, &sumR, 1, MPI_DOUBLE, MPI_SUM, 0, MPI_COMM_WORLD);
MPI_Allreduce(&sum0, &sumA, 1, MPI_DOUBLE, MPI_SUM, MPI_COMM_WORLD);
printf("before BCAST %5d %15.0F %15.0F\n", MyRank, sumA, sumR);

MPI_Bcast(&sumR, 1, MPI_DOUBLE, 0, MPI_COMM_WORLD);
printf("after BCAST %5d %15.0F %15.0F\n", MyRank, sumA, sumR);

MPI_Finalize();

return 0;
}
```

Example: Dot Product (3/3)

```
<$O-S1>/allreduce.c
```

```
MPI_Reduce(&sum0, &sumR, 1, MPI_DOUBLE, MPI_SUM, 0, MPI_COMM_WORLD);  
MPI_Allreduce(&sum0, &sumA, 1, MPI_DOUBLE, MPI_SUM, MPI_COMM_WORLD);
```

Dot Product

Summation of results of each process (sum0)

“sumR” has value only on PE#0.

“sumA” has value on all processes by MPI_Allreduce

```
MPI_Bcast(&sumR, 1, MPI_DOUBLE, 0, MPI_COMM_WORLD);
```

“sumR” has value on PE#1-#3 by MPI_Bcast

Execute <\$0-S1>/allreduce.f/c

```
$> cd /work/gt18/t18XXX/pFEM/mpi/S1
$> module load fj
$> mpifrtpx -Kfast allreduce.f
$> mpifccpx -Nclang -Kfast allreduce.c
(modify go4.sh, 4-processes)
$> pjsub go4.sh
```

```
(my_rank, sumALLREDUCE, sumREDUCE)
before BCAST      0      1.2000000E+02      1.2000000E+02
after  BCAST      0      1.2000000E+02      1.2000000E+02

before BCAST      1      1.2000000E+02      0.0000000E+00
after  BCAST      1      1.2000000E+02      1.2000000E+02

before BCAST      3      1.2000000E+02      0.0000000E+00
after  BCAST      3      1.2000000E+02      1.2000000E+02

before BCAST      2      1.2000000E+02      0.0000000E+00
after  BCAST      2      1.2000000E+02      1.2000000E+02
```

Examples by Collective Comm.

- Dot Products of Vectors
- **Scatter/Gather**
- Reading Distributed Files
- MPI_Allgatherv

Global/Local Data (1/3)

- Parallelization of an easy process where a real number α is added to each component of real vector **VECg**:

```
do i= 1, NG
  VECg(i)= VECg(i) + ALPHA
enddo
```

```
for (i=0; i<NG; i++){
  VECg[i]= VECg[i] + ALPHA
}
```

Global/Local Data (2/3)

- Configurationa
 - **NG= 32 (length of the vector)**
 - **ALPHA=1000.**
 - Process # of MPI= 4
- Vector VECg has following 32 components
(`<$O-S1>/a1x.all`):

101.0,	103.0,	105.0,	106.0,	109.0,	111.0,	121.0,	151.0,
201.0,	203.0,	205.0,	206.0,	209.0,	211.0,	221.0,	251.0,
301.0,	303.0,	305.0,	306.0,	309.0,	311.0,	321.0,	351.0,
401.0,	403.0,	405.0,	406.0,	409.0,	411.0,	421.0,	451.0)

Global/Local Data (3/3)

- Procedure
 - ① Reading vector **VECg** with length=32 from one process (*e.g.* 0th process)
 - Global Data
 - ② Distributing vector components to 4 MPI processes equally (*i.e.* length= 8 for each processes)
 - Local Data, Local ID/Numbering
 - ③ Adding **ALPHA** to each component of the local vector (with length= 8) on each process.
 - ④ Merging the results to global vector with length= 32.
- Actually, we do not need parallel computers for such a kind of small computation.

Operations of Scatter/Gather (1/8)

Reading **VECg** (length=32) from a process (*e.g.* #0)

- Reading global data from #0 process

```
include 'mpif.h'
integer, parameter :: NG= 32
real(kind=8), dimension(NG):: VECg

call MPI_INIT (ierr)
call MPI_COMM_SIZE (<comm>, PETOT , ierr)
call MPI_COMM_RANK (<comm>, my_rank, ierr)

if (my_rank.eq.0) then
  open (21, file= 'a1x.all', status= 'unknown')
  do i= 1, NG
    read (21,*) VECg(i)
  enddo
  close (21)
endif
```

```
#include <mpi.h>
#include <stdio.h>
#include <math.h>
#include <assert.h>

int main(int argc, char **argv) {
  int i, NG=32;
  int PeTot, MyRank, MPI_Comm;
  double VECg[32];
  char filename[80];
  FILE *fp;

  MPI_Init(&argc, &argv);
  MPI_Comm_size(<comm>, &PeTot);
  MPI_Comm_rank(<comm>, &MyRank);

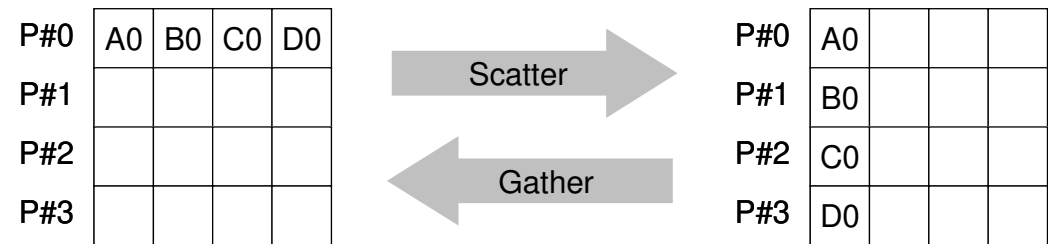
  fp = fopen("a1x.all", "r");
  if(!MyRank) for(i=0;i<NG;i++) {
    fscanf(fp, "%lf", &VECg[i]);
  }
}
```

Operations of Scatter/Gather (2/8)

Distributing global data to 4 process equally (*i.e.* length=8 for each process)

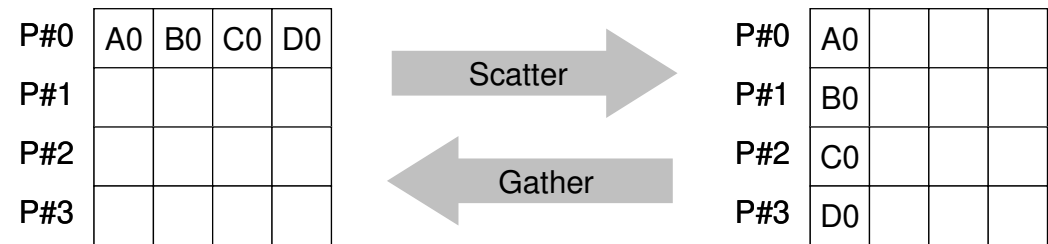
- MPI_Scatter

MPI_Scatter



- Sends data from one process to all other processes in a communicator
 - `scount`-size messages are sent to each process
- **MPI_Scatter** (`sendbuf`, `scount`, `sendtype`, `recvbuf`, `rcount`, `recvtype`, `root`, `comm`)
 - **`sendbuf`** choice I starting address of sending buffer
type is defined by "datatype"
 - **`scount`** int I number of elements sent to each process
 - **`sendtype`** MPI_Datatype I data type of elements of sending buffer
 - FORTRAN MPI_INTEGER, MPI_REAL, MPI_DOUBLE_PRECISION, MPI_CHARACTER etc.
 - C MPI_INT, MPI_FLOAT, MPI_DOUBLE, MPI_CHAR etc.
 - **`recvbuf`** choice O starting address of receiving buffer
 - **`rcount`** int I number of elements received from the root process
 - **`recvtype`** MPI_Datatype I data type of elements of receiving buffer
 - **`root`** int I *rank of root process*
 - **`comm`** MPI_Comm I communicator

MPI_Scatter (cont.)



- **MPI_Scatter** (**sendbuf**, **scount**, **sendtype**, **recvbuf**, **rcount**, **recvtype**, **root**, **comm**)
 - **sendbuf** choice I starting address of sending buffer
 - **scount** int I number of elements sent to each process
 - **sendtype** MPI_Datatype I data type of elements of sending buffer
 - **recvbuf** choice O starting address of receiving buffer
 - **rcount** int I number of elements received from the root process
 - **recvtype** MPI_Datatype I data type of elements of receiving buffer
 - **root** int I **rank of root process**
 - **comm** MPI_Comm I communicator

- Usually
 - **scount = rcount**
 - **sendtype= recvtype**
- This function sends **scount** components starting from **sendbuf** (sending buffer) at process **#root** to each process in **comm**. Each process receives **rcount** components starting from **recvbuf** (receiving buffer).

Operations of Scatter/Gather (3/8)

Distributing global data to 4 processes equally

- Allocating receiving buffer **VEC** (length=8) at each process.
- 8 components sent from sending buffer **VECg** of process #0 are received at each process #0-#3 as 1st-8th components of receiving buffer **VEC**.

```
integer, parameter :: N = 8
real(kind=8), dimension(N) :: VEC
...
call MPI_Scatter
      (VECg, N, MPI_DOUBLE_PRECISION, &
       VEC, N, MPI_DOUBLE_PRECISION, &
       0, <comm>, ierr)
```

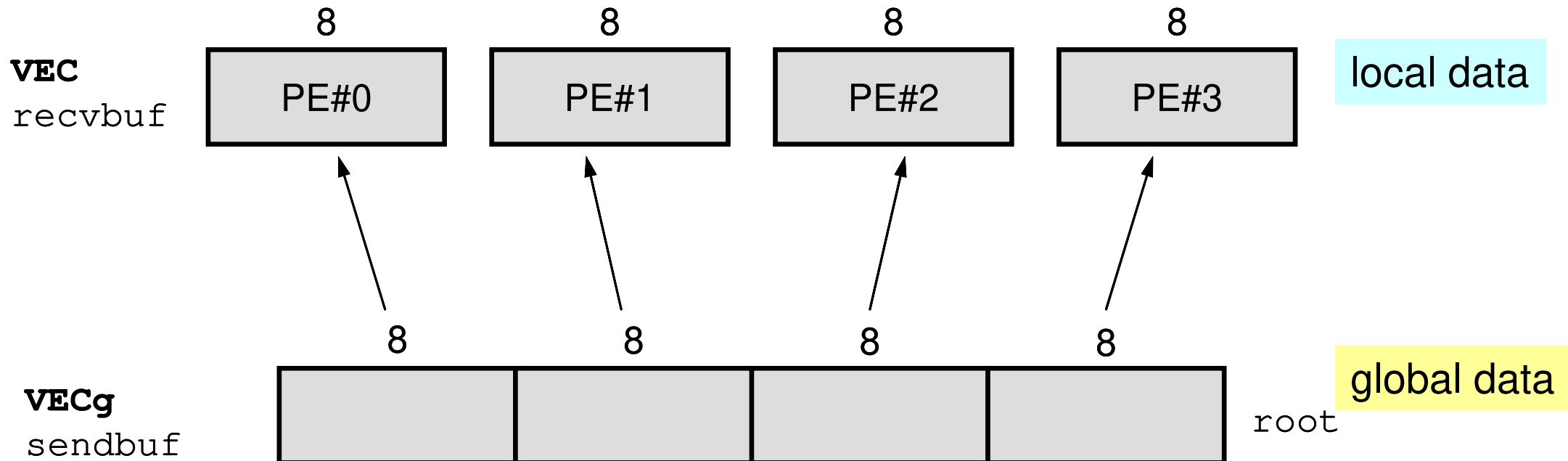
```
int N=8;
double VEC [8];
...
MPI_Scatter (VECg, N, MPI_DOUBLE, VEC, N,
            MPI_DOUBLE, 0, <comm>);
```

```
call MPI_SCATTER
(sendbuf, scount, sendtype, recvbuf, rcount,
recvtype, root, comm, ierr)
```


Operations of Scatter/Gather (4/8)

Distributing global data to 4 processes equally

- 8 components are *scattered* to each process from root (#0)
- 1st-8th components of **VECg** are stored as 1st-8th ones of **VEC** at **#0**,
9th-16th components of **VECg** are stored as 1st-8th ones of **VEC** at **#1**,
etc.
 - **VECg**: Global Data, **VEC**: Local Data



Operations of Scatter/Gather (5/8)

Distributing global data to 4 processes equally

- Global Data: 1st-32nd components of **VECg** at **#0**
- Local Data: 1st-8th components of **VEC** at each process
- Each component of **VEC** can be written from each process in the following way:

```
do i= 1, N
  write (*, '(a, 2i8, f10.0)') 'before', my_rank, i, VEC(i)
enddo
```

```
for(i=0; i<N; i++) {
  printf("before %5d %5d %10.0F\n", MyRank, i+1, VEC[i]);
}
```

Operations of Scatter/Gather (5/8)

Distributing global data to 4 processes equally

- Global Data: 1st-32nd components of **VECg** at **#0**
- Local Data: 1st-8th components of **VEC** at each process
- Each component of **VEC** can be written from each process in the following way:

PE#0

```
before 0 1 101.  
before 0 2 103.  
before 0 3 105.  
before 0 4 106.  
before 0 5 109.  
before 0 6 111.  
before 0 7 121.  
before 0 8 151.
```

PE#1

```
before 1 1 201.  
before 1 2 203.  
before 1 3 205.  
before 1 4 206.  
before 1 5 209.  
before 1 6 211.  
before 1 7 221.  
before 1 8 251.
```

PE#2

```
before 2 1 301.  
before 2 2 303.  
before 2 3 305.  
before 2 4 306.  
before 2 5 309.  
before 2 6 311.  
before 2 7 321.  
before 2 8 351.
```

PE#3

```
before 3 1 401.  
before 3 2 403.  
before 3 3 405.  
before 3 4 406.  
before 3 5 409.  
before 3 6 411.  
before 3 7 421.  
before 3 8 451.
```

Operations of Scatter/Gather (6/8)

On each process, **ALPHA** is added to each of 8 components of **VEC**

- On each process, computation is in the following way

```
real(kind=8), parameter :: ALPHA= 1000.
do i= 1, N
  VEC(i)= VEC(i) + ALPHA
enddo
```

```
double ALPHA=1000. ;
...
for (i=0; i<N; i++) {
  VEC[i]= VEC[i] + ALPHA;}

```

- Results:

<u>PE#0</u>			
after	0	1	1101.
after	0	2	1103.
after	0	3	1105.
after	0	4	1106.
after	0	5	1109.
after	0	6	1111.
after	0	7	1121.
after	0	8	1151.

<u>PE#1</u>			
after	1	1	1201.
after	1	2	1203.
after	1	3	1205.
after	1	4	1206.
after	1	5	1209.
after	1	6	1211.
after	1	7	1221.
after	1	8	1251.

<u>PE#2</u>			
after	2	1	1301.
after	2	2	1303.
after	2	3	1305.
after	2	4	1306.
after	2	5	1309.
after	2	6	1311.
after	2	7	1321.
after	2	8	1351.

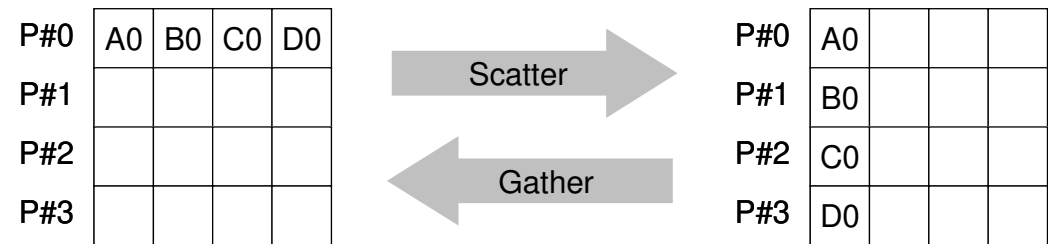
<u>PE#3</u>			
after	3	1	1401.
after	3	2	1403.
after	3	3	1405.
after	3	4	1406.
after	3	5	1409.
after	3	6	1411.
after	3	7	1421.
after	3	8	1451.

Operations of Scatter/Gather (7/8)

Merging the results to global vector with length= 32

- Using MPI_Gather (inverse operation to MPI_Scatter)

MPI_Gather



- Gathers together values from a group of processes, inverse operation to MPI_Scatter
- **MPI_Gather (sendbuf, scount, sendtype, recvbuf, rcount, recvtype, root, comm)**
 - **sendbuf** choice I starting address of sending buffer
 - **scount** int I number of elements sent to each process
 - **sendtype** MPI_Datatype I data type of elements of sending buffer
 - **recvbuf** choice O starting address of receiving buffer
 - **rcount** int I number of elements received from the root process
 - **recvtype** MPI_Datatype I data type of elements of receiving buffer
 - **root** int I **rank of root process**
 - **comm** MPI_Comm I communicator
- **recvbuf** is on **root** process.

Operations of Scatter/Gather (8/8)

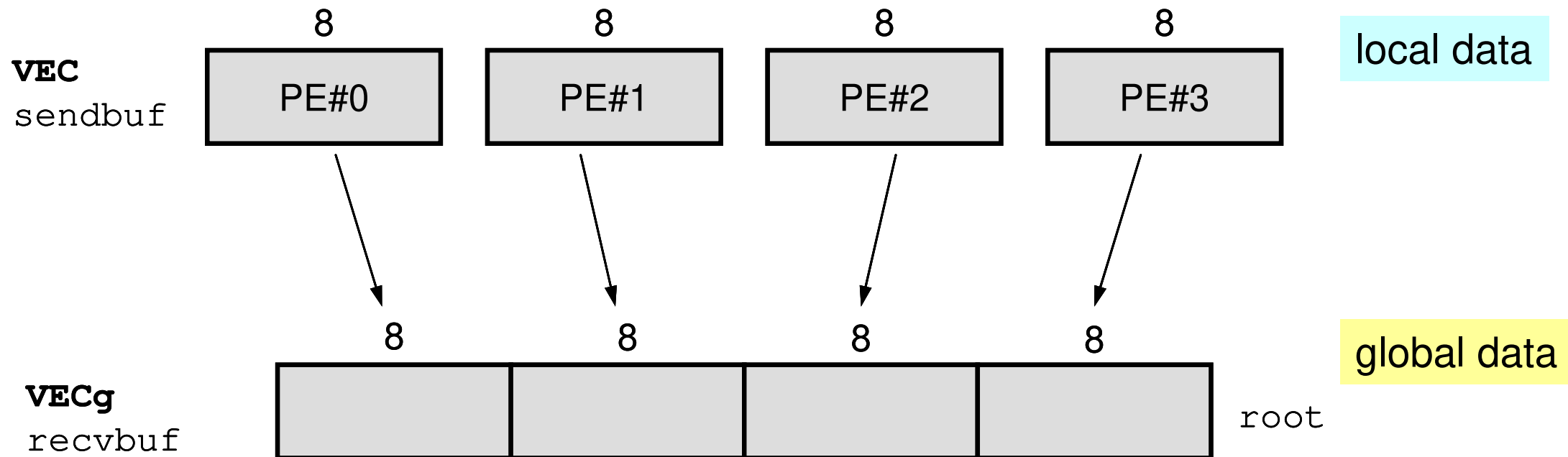
Merging the results to global vector with length= 32

- Each process components of **VEC** to **VECg** on root (#0 in this case).

```
call MPI_Gather
      (VEC , N, MPI_DOUBLE_PRECISION, &
      VECg, N, MPI_DOUBLE_PRECISION, &
      0, <comm>, ierr)
```

```
MPI_Gather (VEC, N, MPI_DOUBLE, VECg, N,
           MPI_DOUBLE, 0, <comm>);
```

- 8 components are gathered from each process to the root process.



<\$0-S1>/scatter-gather.f/c

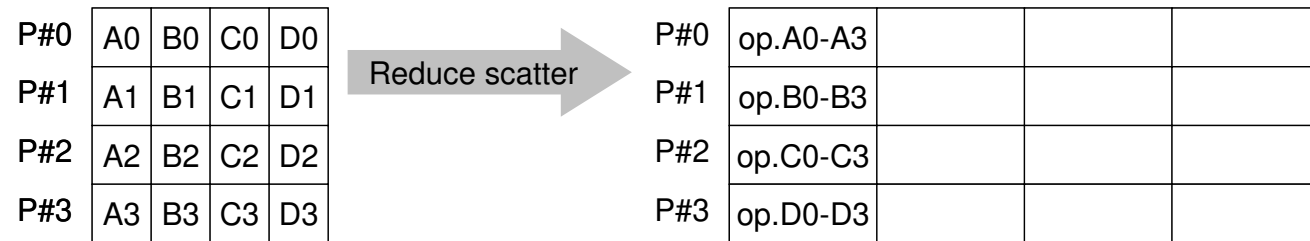
```
$> cd /work/gt18/t18XXX/pFEM/mpi/S1
$> module load fj
$> mpifccpx -Nclang -Kfast scatter-gather.c
$> mpifrtpx -Kfast scatter-gather.f
```

(modify go4.sh, 4-processes)

```
$> pjsub go4.sh
```

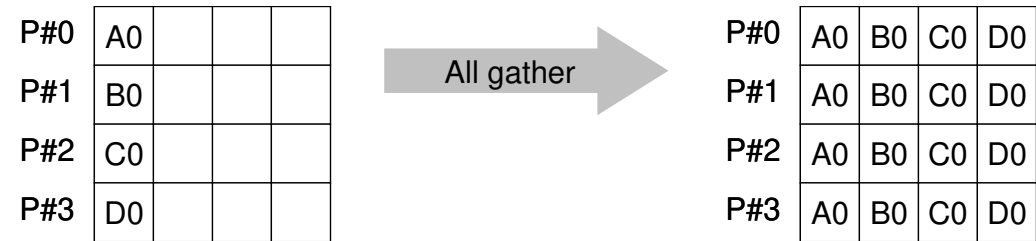
PE#0	PE#1	PE#2	PE#3
before 0 1 101.	before 1 1 201.	before 2 1 301.	before 3 1 401.
before 0 2 103.	before 1 2 203.	before 2 2 303.	before 3 2 403.
before 0 3 105.	before 1 3 205.	before 2 3 305.	before 3 3 405.
before 0 4 106.	before 1 4 206.	before 2 4 306.	before 3 4 406.
before 0 5 109.	before 1 5 209.	before 2 5 309.	before 3 5 409.
before 0 6 111.	before 1 6 211.	before 2 6 311.	before 3 6 411.
before 0 7 121.	before 1 7 221.	before 2 7 321.	before 3 7 421.
before 0 8 151.	before 1 8 251.	before 2 8 351.	before 3 8 451.
after 0 1 1101.	after 1 1 1201.	after 2 1 1301.	after 3 1 1401.
after 0 2 1103.	after 1 2 1203.	after 2 2 1303.	after 3 2 1403.
after 0 3 1105.	after 1 3 1205.	after 2 3 1305.	after 3 3 1405.
after 0 4 1106.	after 1 4 1206.	after 2 4 1306.	after 3 4 1406.
after 0 5 1109.	after 1 5 1209.	after 2 5 1309.	after 3 5 1409.
after 0 6 1111.	after 1 6 1211.	after 2 6 1311.	after 3 6 1411.
after 0 7 1121.	after 1 7 1221.	after 2 7 1321.	after 3 7 1421.
after 0 8 1151.	after 1 8 1251.	after 2 8 1351.	after 3 8 1451.

MPI_Reduce_scatter



- MPI_Reduce + MPI_Scatter
- **MPI_Reduce_Scatter** (**sendbuf**, **recvbuf**, **rcount**, **datatype**, **op**, **comm**)
 - **sendbuf** choice I starting address of sending buffer
 - **recvbuf** choice O starting address of receiving buffer
 - **rcount** int I integer array specifying the number of elements in result distributed to each process. Array must be identical on all calling processes.
 - **datatype** MPI_Datatype I data type of elements of sending/receiving buffer
 - **op** MPI_Op I reduce operation
MPI_MAX, MPI_MIN, MPI_SUM, MPI_PROD, MPI_LAND, MPI_BAND etc
 - **comm** MPI_Comm I communicator

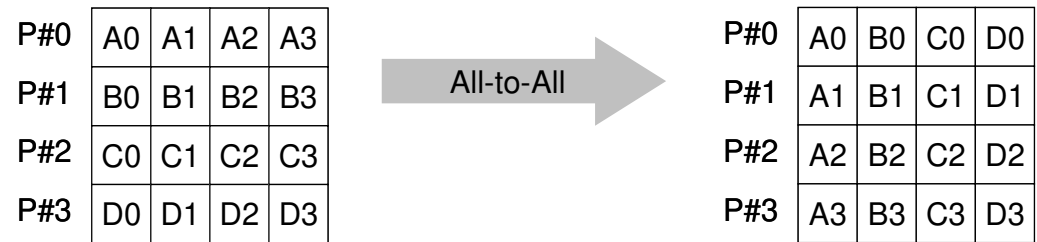
MPI_Allgather



- MPI_Gather + MPI_Bcast
 - Gathers data from all tasks and distribute the combined data to all tasks

- **MPI_Allgather (sendbuf, scount, sendtype, recvbuf, rcount, recvtype, comm)**
 - **sendbuf** choice I starting address of sending buffer
 - **scount** int I number of elements sent to each process
 - **sendtype** MPI_Datatype I data type of elements of sending buffer
 - **recvbuf** choice O starting address of receiving buffer
 - **rcount** int I number of elements received from each process
 - **recvtype** MPI_Datatype I data type of elements of receiving buffer
 - **comm** MPI_Comm I communicator

MPI_Alltoall



- Sends data from all to all processes: transformation of dense matrix
- **MPI_Alltoall** (**sendbuf**, **scount**, **sendtype**, **recvbuf**, **rcount**, **recvtype**, **comm**)
 - **sendbuf** choice I starting address of sending buffer
 - **scount** int I number of elements sent to each process
 - **sendtype** MPI_Datatype I data type of elements of sending buffer
 - **recvbuf** choice O starting address of receiving buffer
 - **rcount** int I number of elements received from the root process
 - **recvtype** MPI_Datatype I data type of elements of receiving buffer
 - **comm** MPI_Comm I communicator

Examples by Collective Comm.

- Dot Products of Vectors
- Scatter/Gather
- **Reading Distributed Files**
- MPI_Allgatherv

Operations of Distributed Local Files

- In Scatter/Gather example, PE#0 reads global data, that is *scattered* to each processor, then parallel operations are done.
- If the problem size is very large, a single processor may not read entire global data.
 - If the entire global data is decomposed to distributed local data sets, each process can read the local data.
 - If global operations are needed to a certain sets of vectors, MPI functions, such as MPI_Gather etc. are available.

Reading Distributed Local Files: Uniform Vec. Length (1/2)

```
>$ cd /work/gt18/t18XXX/pFEM/mpi/S1
>$ module load fj
>$ ls a1.*
a1.0 a1.1 a1.2 a1.3    a1x.all is decomposed to
                        4 files.
>$ mpifccpx -Nclang -Kfast file.c
>$ mpifrtpx -Kfast file.f
(modify go4.sh for 4 processes)
>$ pjsub go4.sh
```

a1.0

101.0
103.0
105.0
106.0
109.0
111.0
121.0
151.0

a1.1

201.0
203.0
205.0
206.0
209.0
211.0
221.0
251.0

a1.2

301.0
303.0
305.0
306.0
309.0
311.0
321.0
351.0

a1.3

401.0
403.0
405.0
406.0
409.0
411.0
421.0
451.0

go4.sh

```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture8-o
#PJM -L node=1
#PJM --mpi proc=4
#PJM -L elapse=00:15:00
#PJM -g gt18
#PJM -j
#PJM -e err
#PJM -o test.lst

module load fj
module load fjmpi

mpiexec ./a.out
```

Job Name
Name of "QUEUE"
Node#
Total MPI Process#
Computation Time
Group Name (Wallet)
Standard Error
Standard Output

Reading Distributed Local Files: Uniform Vec. Length (2/2)

<\$O-S1>/file.c

```
int main(int argc, char **argv){
    int i;
    int PeTot, MyRank;
    MPI_Comm SolverComm;
    double vec[8];
    char FileName[80];
    FILE *fp;

    MPI_Init(&argc, &argv);
    MPI_Comm_size(MPI_COMM_WORLD, &PeTot);
    MPI_Comm_rank(MPI_COMM_WORLD, &MyRank);

    sprintf(FileName, "a1.%d", MyRank);

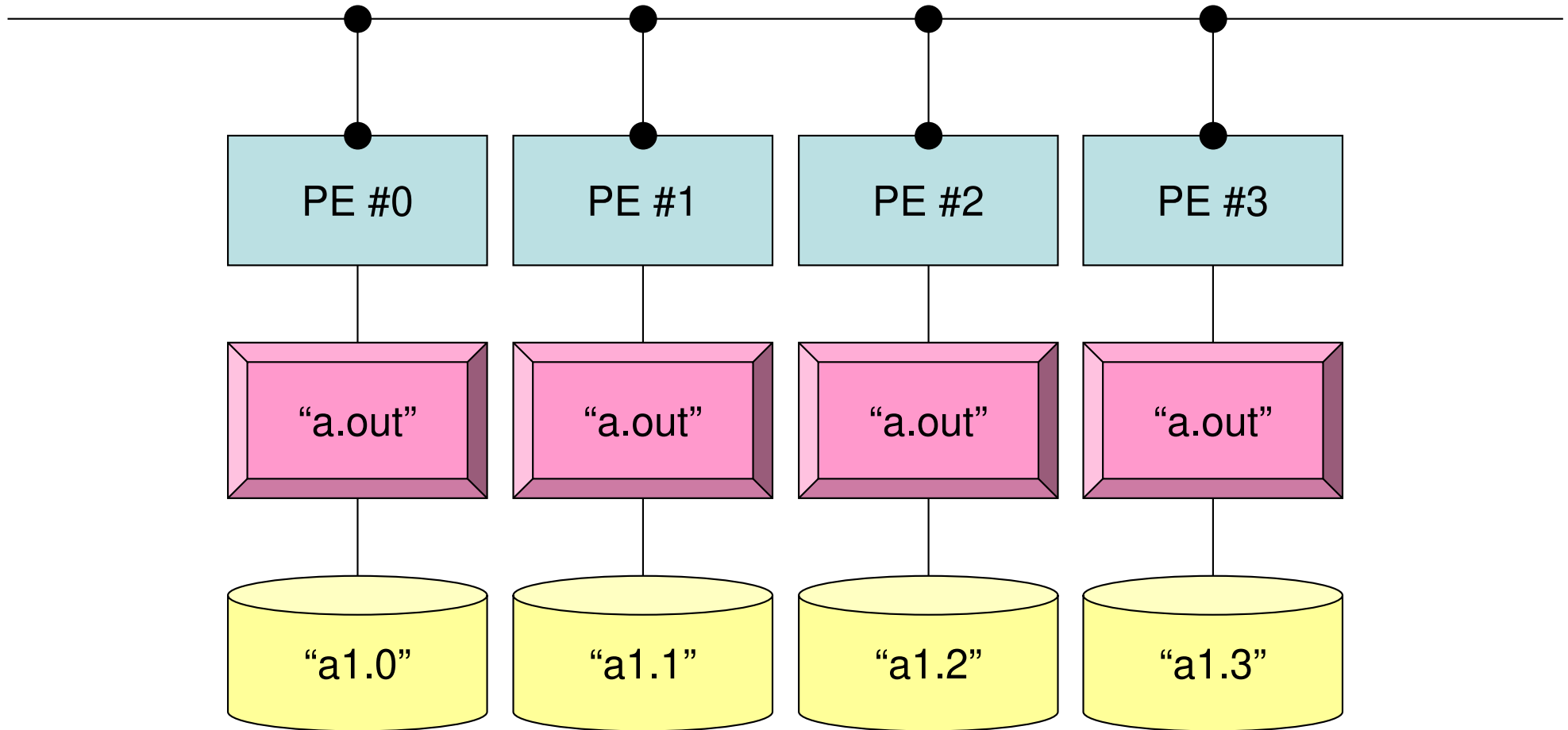
    fp = fopen(FileName, "r");
    if(fp == NULL) MPI_Abort(MPI_COMM_WORLD, -1)
    for (i=0; i<8; i++){
            fscanf(fp, "%lf", &vec[i]);
    }

    for(i=0; i<8; i++){
        printf("%5d%5d%10.0f¥n", MyRank, i+1, vec[i]);
    }
    MPI_Finalize();
    return 0;
}
```

Similar to
"Hello"

Local ID is 0-7

Typical SPMD Operation



```
mpirun -np 4 a.out
```

Non-Uniform Vector Length (1/2)

```
>$ cd /work/gt18/t18XXX/pFEM/mpi/S1
>$ module load fj
>$ ls a2.*
  a2.0 a2.1 a2.2 a2.3
>$ cat a2.1
  5          Number of Components at each Process
201.0       Components
203.0
205.0
206.0
209.0

>$ mpifccpx -Nclang -Kfast file2.c
>$ mpifrtpx -Kfast file2.f

(modify go4.sh for 4 processes)
>$ pjsub go4.sh
```

a2.0~a2.3

PE#0

8
101.0
103.0
105.0
106.0
109.0
111.0
121.0
151.0

PE#1

5
201.0
203.0
205.0
206.0
209.0

PE#2

7
301.0
303.0
305.0
306.0
311.0
321.0
351.0

PE#3

3
401.0
403.0
405.0

Non-Uniform Vector Length (2/2)

<\$O-S1>/file2.c

```
int main(int argc, char **argv){
    int i, int PeTot, MyRank;
    MPI_Comm SolverComm;
    double *vec, *vec2, *vecg;
    int num;
    double sum0, sum;
    char filename[80];
    FILE *fp;

    MPI_Init(&argc, &argv);
    MPI_Comm_size(MPI_COMM_WORLD, &PeTot);
    MPI_Comm_rank(MPI_COMM_WORLD, &MyRank);

    sprintf(filename, "a2.%d", MyRank);
    fp = fopen(filename, "r");
    assert(fp != NULL);

    fscanf(fp, "%d", &num);
    vec = malloc(num * sizeof(double));
    for(i=0;i<num;i++){fscanf(fp, "%lf", &vec[i]);}

    for(i=0;i<num;i++){
        printf(" %5d%5d%5d%10.0f¥n", MyRank, i+1, num, vec[i]);}

    MPI_Finalize();
}
```

“num” is different at each process

How to generate local data

- Reading global data ($N=NG$)
 - Scattering to each process
 - Parallel processing on each process
 - (If needed) reconstruction of global data by gathering local data
- Generating local data ($N=NL$), or reading distributed local data
 - Generating or reading local data on each process
 - Parallel processing on each process
 - (If needed) reconstruction of global data by gathering local data
- In future, latter case is more important, but former case is also introduced in this class for understanding of operations of global/local data.

Examples by Collective Comm.

- Dot Products of Vectors
- Scatter/Gather
- Reading Distributed Files
- **MPI_Allgatherv**

MPI_Gatherv, MPI_Scatterv

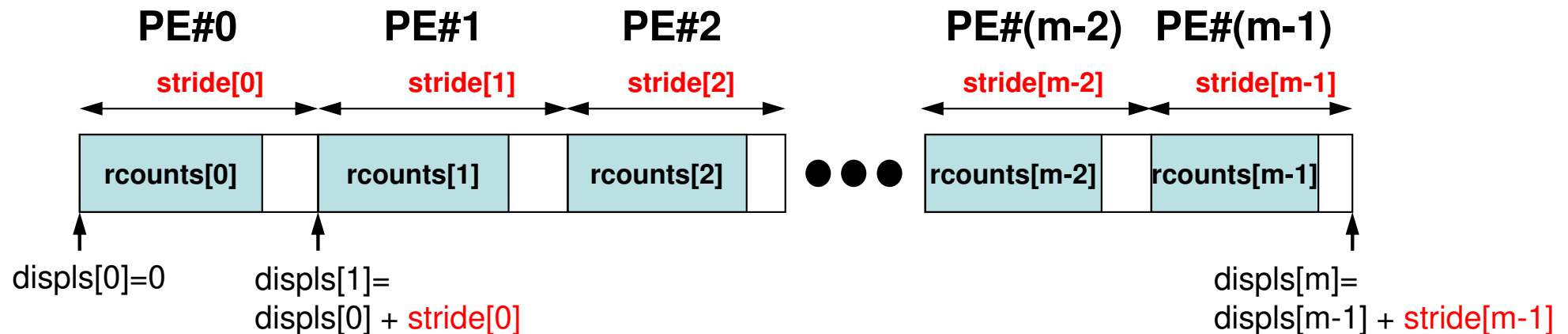
- MPI_Gather, MPI_Scatter
 - Length of message from/to each process is uniform
- MPI_XXXv extends functionality of MPI_XXX by allowing a varying count of data from each process:
 - MPI_Gatherv
 - MPI_Scatterv
 - MPI_Allgatherv
 - MPI_Alltoallv

MPI_Allgatherv

- Variable count version of MPI_Allgather
 - creates “global data” from “local data”
- **MPI_Allgatherv (sendbuf, scount, sendtype, recvbuf, rcounts, displs, recvtype, comm)**
 - **sendbuf** choice I starting address of sending buffer
 - **scount** int I number of elements sent to each process
 - **sendtype** MPI_Datatype I data type of elements of sending buffer
 - **recvbuf** choice O starting address of receiving buffer
 - **rcounts** int I integer array (of length *groupsize*) containing the number of elements that are to be received from each process (array: size= PETOT)
 - **displs** int I integer array (of length *groupsize*). Entry *i* specifies the displacement (relative to *recvbuf*) at which to place the incoming data from process *i* (array: size= PETOT+1)
 - **recvtype** MPI_Datatype I data type of elements of receiving buffer
 - **comm** MPI_Comm I communicator

MPI_Allgatherv (cont.)

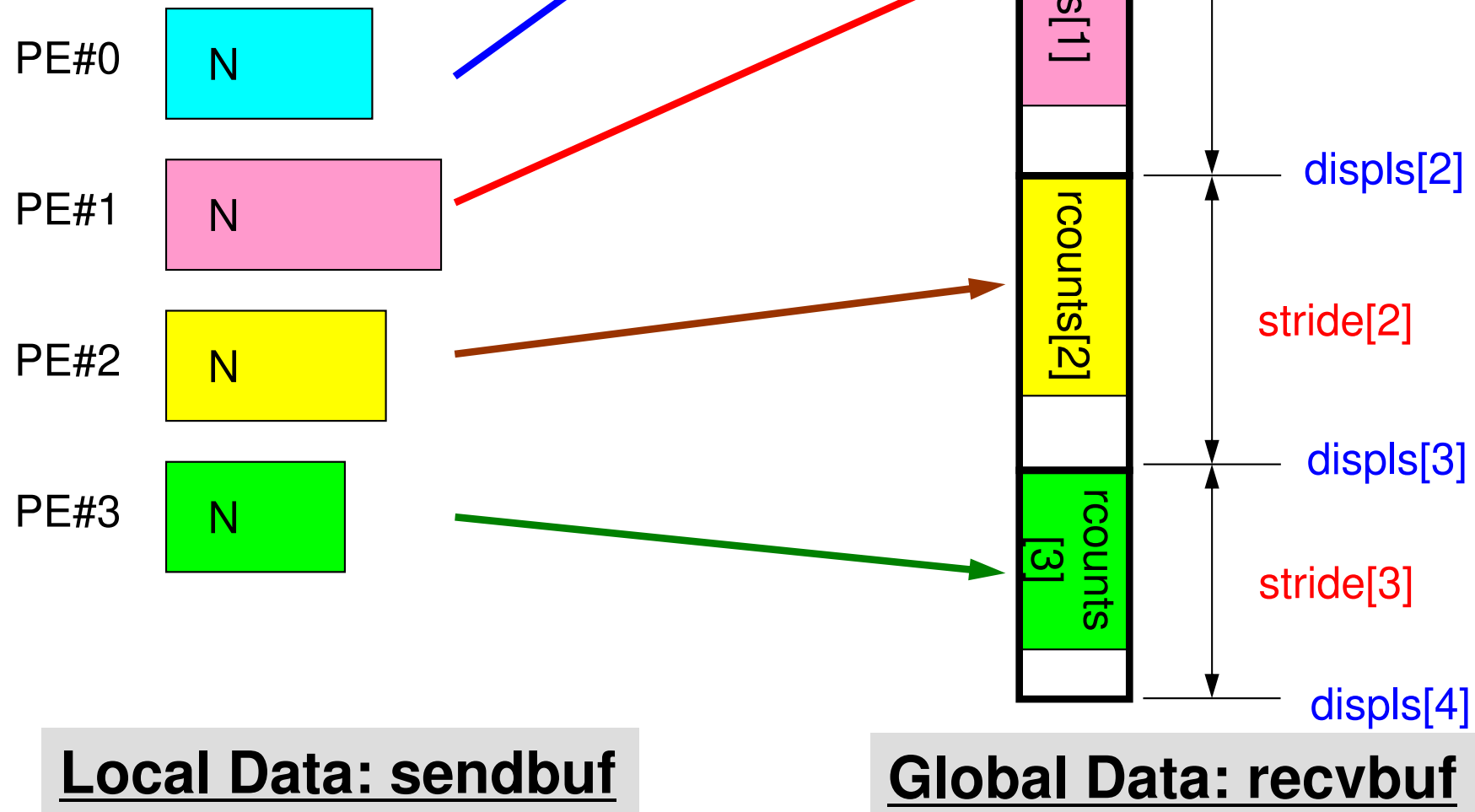
- `MPI_Allgatherv (sendbuf, scount, sendtype, recvbuf, rcounts, displs, recvtype, comm)`
 - **`rcounts`** `int` `I` integer array (of length *groupsize*) containing the number of elements that are to be received from each process (array: size= `PETOT`)
 - **`displs`** `int` `I` integer array (of length *groupsize*). Entry *i* specifies the displacement (relative to `recvbuf`) at which to place the incoming data from process *i* (array: size= `PETOT+1`)
 - These two arrays are related to size of final “global data”, therefore each process requires information of these arrays (`rcounts`, `displs`)
 - Each process must have same values for all components of both vectors
 - Usually, **`stride(i) = rcounts(i)`**



$$\text{size[recvbuf]} = \text{displs[PETOT]} = \text{sum[stride]}$$

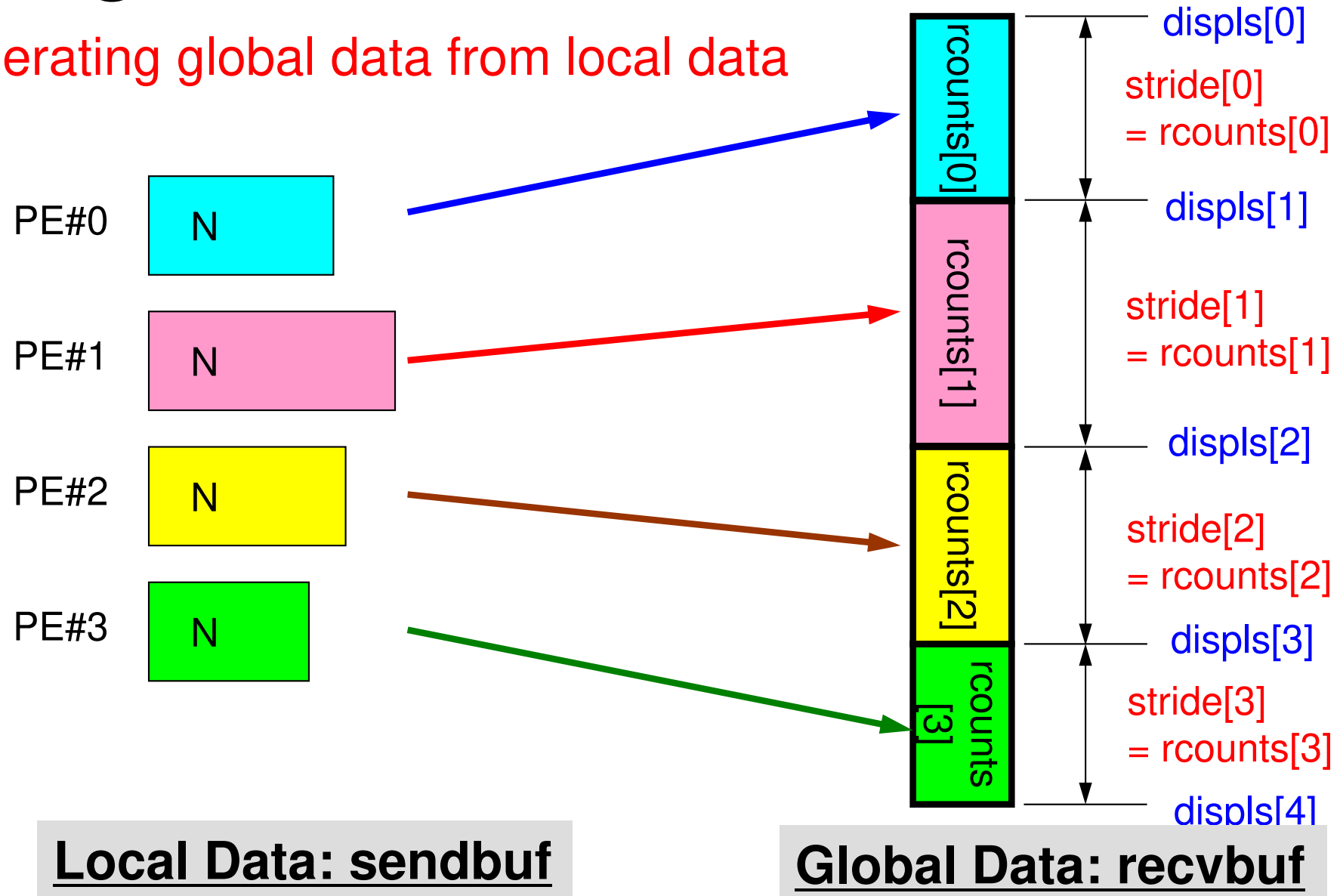
What MPI_Allgatherv is doing

Generating global data from local data



What MPI_Allgatherv is doing

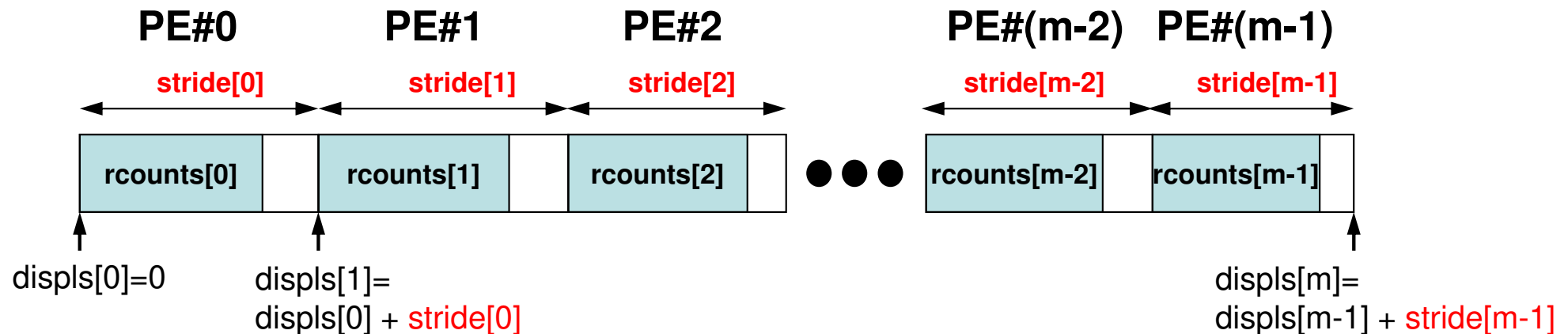
Generating global data from local data



MPI_Allgatherv in detail (1/2)

C

- `MPI_Allgatherv (sendbuf, scount, sendtype, recvbuf, rcounts, displs, recvtype, comm)`
- **rcounts**
 - Size of message from each PE: Size of Local Data (Length of Local Vector)
- **displs**
 - Address/index of each local data in the vector of global data
 - **displs (PETOT+1) = Size of Entire Global Data (Global Vector)**

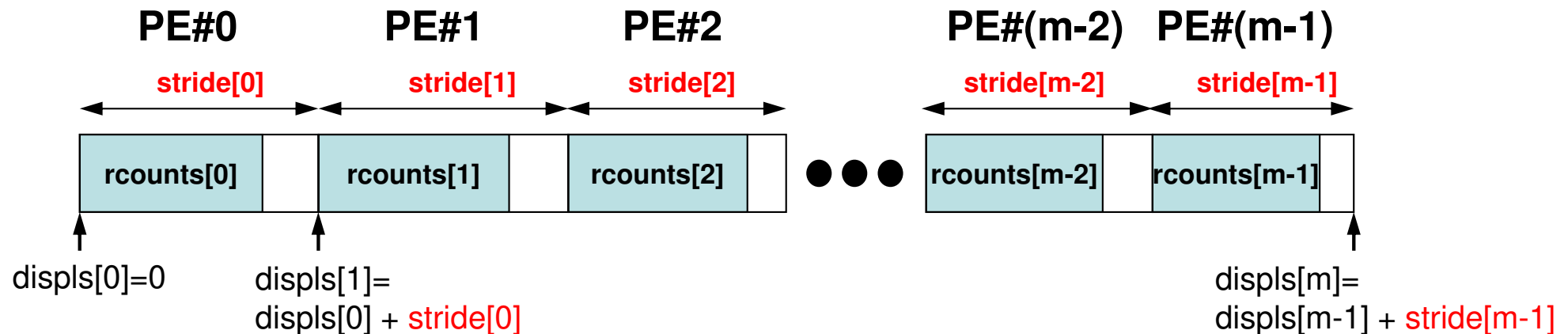


$$\text{size[recvbuf]} = \text{displs[PETOT]} = \text{sum[stride]}$$

MPI_Allgatherv in detail (2/2)



- Each process needs information of **rcounts** & **displs**
 - “**rcounts**” can be created by gathering local vector length “**N**” from each process.
 - On each process, “**displs**” can be generated from “**rcounts**” on each process.
 - $\text{stride}[i] = \text{rcounts}[i]$
 - Size of “**recvbuf**” is calculated by summation of “**rcounts**”.



$$\text{size}[\text{recvbuf}] = \text{displs}[\text{PETOT}] = \text{sum}[\text{stride}]$$

Preparation for MPI_Allgather

<O-S1>/agv.c

- Generating global vector from “a2.0”~”a2.3”.
- Length of the each vector is 8, 5, 7, and 3, respectively. Therefore, size of final global vector is 23 (= 8+5+7+3).

a2.0~a2.3

PE#0

8
101.0
103.0
105.0
106.0
109.0
111.0
121.0
151.0

PE#1

5
201.0
203.0
205.0
206.0
209.0

PE#2

7
301.0
303.0
305.0
306.0
311.0
321.0
351.0

PE#3

3
401.0
403.0
405.0

Preparation: MPI_Allgatherv (1/4)

C

```
<$O-S1>/agv.c
```

```
int main(int argc, char **argv) {
    int i;
    int PeTot, MyRank;
    MPI_Comm SolverComm;
    double *vec, *vec2, *vecg;
    int *Rcounts, *Displs;
    int n;
    double sum0, sum;
    char filename[80];
    FILE *fp;

    MPI_Init(&argc, &argv);
    MPI_Comm_size(MPI_COMM_WORLD, &PeTot);
    MPI_Comm_rank(MPI_COMM_WORLD, &MyRank);

    sprintf(filename, "a2.%d", MyRank);
    fp = fopen(filename, "r");
    assert(fp != NULL);

    fscanf(fp, "%d", &n);
    vec = malloc(n * sizeof(double));
    for(i=0; i<n; i++) {
        fscanf(fp, "%lf", &vec[i]);
    }
}
```

n (NL) is different at
each process

Preparation: MPI_Allgatherv (2/4)

C

<\$O-S1>/agv.c

```
Rcounts= calloc(PeTot, sizeof(int));
Displs = calloc(PeTot+1, sizeof(int));
```

```
printf("before %d %d", MyRank, n);
for(i=0;i<PeTot;i++){printf(" %d", Rcounts[i]);}
```

```
MPI_Allgather(&n, 1, MPI_INT, Rcounts, 1, MPI_INT, MPI_COMM_WORLD);
```

```
printf("after %d %d", MyRank, n);
for(i=0;i<PeTot;i++){printf(" %d", Rcounts[i]);}
```

Rcounts on each PE

```
Displs[0] = 0;
```

PE#0 N=8

PE#1 N=5

PE#2 N=7

PE#3 N=3



MPI_Allgather

Rcounts[0:3]= {8, 5, 7, 3}

Rcounts[0:3]={8, 5, 7, 3}

Rcounts[0:3]={8, 5, 7, 3}

Rcounts[0:3]={8, 5, 7, 3}

Preparation: MPI_Allgatherv (2/4)

C

<\$O-S1>/agv.c

```

Rcounts= calloc(PeTot, sizeof(int));
Displs = calloc(PeTot+1, sizeof(int));

printf("before %d %d", MyRank, n);
for(i=0;i<PeTot;i++){printf(" %d", Rcounts[i]);}

MPI_Allgather(&n, 1, MPI_INT, Rcounts, 1, MPI_INT, MPI_COMM_WORLD);

printf("after %d %d", MyRank, n);
for(i=0;i<PeTot;i++){printf(" %d", Rcounts[i]);}    Rcounts on each PE

Displs[0] = 0;
for(i=0;i<PeTot;i++){
    Displs[i+1] = Displs[i] + Rcounts[i];}

    printf("CoundIndex %d ", MyRank);                Displs on each PE
    for(i=0;i<PeTot+1;i++){
        printf(" %d", Displs[i]);
    }
    MPI_Finalize();
    return 0;
}

```

Preparation: MPI_Allgatherv (3/4)

```
> cd /work/gt18/t18XXX/pFEM/mpi/S1
> module load fj
> mpifccpx -Nclang -Kfast agv.c
```

(modify go4.sh for 4 processes)

```
> pjsub go4.sh
```

before	0	8	0	0	0	0
after	0	8	8	5	7	3
Displs	0	0	8	13	20	23

before	1	5	0	0	0	0
after	1	5	8	5	7	3
Displs	1	0	8	13	20	23

before	3	3	0	0	0	0
after	3	3	8	5	7	3
Displs	3	0	8	13	20	23

before	2	7	0	0	0	0
after	2	7	8	5	7	3
Displs	2	0	8	13	20	23

```
write (*, '(a,10i8)') "before", my_rank, N, rcounts
write (*, '(a,10i8)') "after ", my_rank, N, rcounts
write (*, '(a,10i8)') "displs", my_rank, displs
```

Preparation: MPI_Allgatherv (4/4)

- Only "recvbuf" is not defined yet.
- Size of "recvbuf" = "Displs[PETOT]"

```
MPI_Allgatherv  
  ( VEC , N, MPI_DOUBLE,  
    recvbuf, rcounts, displs, MPI_DOUBLE,  
    MPI_COMM_WORLD);
```

Report S1 (1/2)

- ~~• Deadline: January 25th (Wed), 2023, 17:00@ITC-LMS~~
- Problem S1-1
 - Read local files $\langle \$O-S1 \rangle/a1.0 \sim a1.3$, $\langle \$O-S1 \rangle/a2.0 \sim a2.3$.
 - Develop codes which calculate norm $\|x\|_2$ of global vector for each case.
 - $\langle \$O-S1 \rangle/file.c$, $\langle \$O-S1 \rangle/file2.c$
- Problem S1-2
 - Read local files $\langle \$O-S1 \rangle/a2.0 \sim a2.3$.
 - Develop a code which constructs “global vector” using `MPI_Allgatherv`.

Report S1 (2/2)

- Problem S1-3

- Develop parallel program which calculates the following numerical integration using “trapezoidal rule” by MPI_Reduce, MPI_Bcast etc.
- Measure computation time, and parallel performance

$$\int_0^1 \frac{4}{1+x^2} dx$$

- Report

- Cover Page: Name, ID, and Problem ID (S1) must be written.
- Less than two pages including figures and tables (A4) for each of three sub-problems
 - Strategy, Structure of the Program, Remarks
- Source list of the program (if you have bugs)
- Output list (as small as possible)

a012.sh

```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture8-o
#PJM -L node=1
#PJM --mpi proc=12
#PJM -L elapse=00:15:00
#PJM -g gt18
#PJM -j
#PJM -e err
#PJM -o test.lst

module load fj
module load fjmpi
mpiexec ./a.out
mpiexec numactl -l ./a.out
```

a048.sh

```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture8-o
#PJM -L node=1
#PJM --mpi proc=48
#PJM -L elapse=00:15:00
#PJM -g gt18
#PJM -j
#PJM -e err
#PJM -o test.lst

module load fj
module load fjmpi
mpiexec ./a.out
mpiexec numactl -l ./a.out
```

a384.sh

```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture8-o
#PJM -L node=8
#PJM --mpi proc=384
#PJM -L elapse=00:15:00
#PJM -g gt18
#PJM -j
#PJM -e err
#PJM -o test.lst

module load fj
module load fjmpi
mpiexec ./a.out
mpiexec numactl -l ./a.out
```

a576.sh

```
#!/bin/sh
#PJM -N "test"
#PJM -L rscgrp=lecture8-o
#PJM -L node=12
#PJM --mpi proc=576
#PJM -L elapse=00:15:00
#PJM -g gt18
#PJM -j
#PJM -e err
#PJM -o test.lst

module load fj
module load fjmpi
mpiexec ./a.out
mpiexec numactl -l ./a.out
```

numactl -l/--localalloc for utilizing local memory (no effects)

Number of Processes

```
#PJM -L node=1; #PJM --mpi proc= 1      1-node, 1-proc, 1-proc/n
#PJM -L node=1; #PJM --mpi proc= 4      1-node, 4-proc, 4-proc/n
#PJM -L node=1; #PJM --mpi proc=12     1-node, 12-proc, 12-proc/n
#PJM -L node=1; #PJM --mpi proc=24     1-node, 24-proc, 24-proc/n
#PJM -L node=1; #PJM --mpi proc=48     1-node, 48-proc, 48-proc/n
```

```
#PJM -L node= 4; #PJM --mpi proc=192   4-node, 192-proc, 48-proc/n
#PJM -L node= 8; #PJM --mpi proc=384   8-node, 384-proc, 48-proc/n
#PJM -L node=12; #PJM --mpi proc=576  12-node, 576-proc, 48-proc/n
```

