

Communication/Synchronization Avoiding/Reducing in Krylov Iterative Solvers on Multicore/Manycore Clusters Early Case Studies on Oakforest-PACS with Intel Xeon Phi (抜粋)

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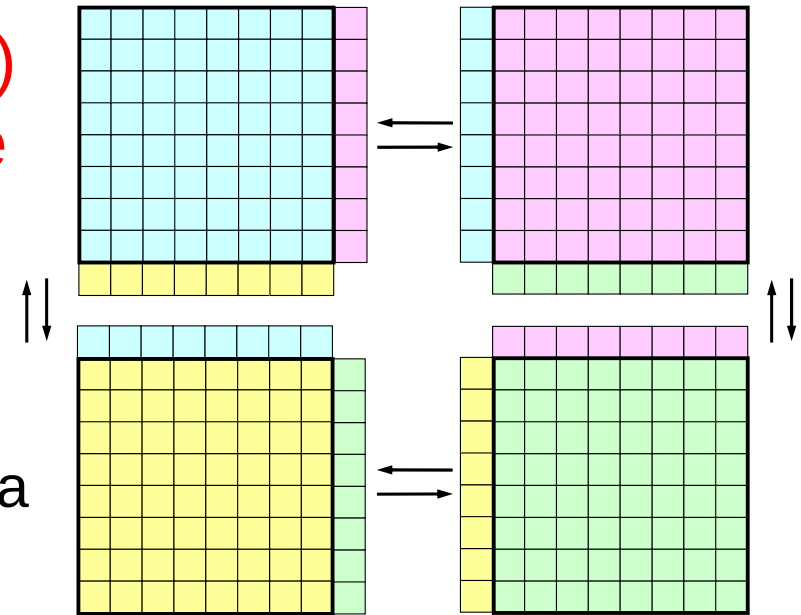
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**Berkeley Lab Computing Sciences Seminar
January 20, 2017**

- Communication/Synchronization
Avoiding/Reducing in Krylov Iterative
Solvers
- Pipelined CG: Background
- Pipelined CG: Results
- Communication-Computation
Overlapping
- Summary
- Preliminary Study on OFP

Parallel (Krylov) Iterative Solvers

- Both of convergence (robustness) and efficiency (single/parallel) are important
- Communications needed
 - SpMV (P2P communications, MPI_Isend/Irecv/Waitall): Local Data Structure with HALO
 - Dot-Products (MPI_Allreduce)
 - ✓ effect of latency
 - Preconditioning (up to algorithm)
- Remedy for Robust Parallel ILU Preconditioner
 - Additive Schwarz Domain Decomposition
 - HID (Hierarchical Interface Decomposition, based on global nested dissection) [Henon & Saad 2007], ext. HID [KN 2010]



Communication/Synchronization Avoiding/Reducing/Hiding for Parallel Preconditioned Krylov Iterative Methods

- SpMV
 - Overlapping of Computations & Communications
 - Matrix Powers Kernel
- Dot Products
 - Pipelined Methods
 - Gropp's Algorithm,

Algorithm 1 Preconditioned CG

```

1:  $r_0 := b - Ax_0; \quad u_0 := M^{-1}r_0; \quad p_0 := u_0$ 
2: for  $i = 0, \dots$  do
3:    $s := Ap_i$ 
4:    $\alpha := (r_i, u_i) / (s, p_i)$ 
5:    $x_{i+1} := x_i + \alpha p_i$ 
6:    $r_{i+1} := r_i - \alpha s$ 
7:    $u_{i+1} := M^{-1}r_{i+1}$ 
8:    $\beta := (r_{i+1}, u_{i+1}) / (r_i, u_i)$ 
9:    $p_{i+1} := u_{i+1} + \beta p_i$ 
10: end for
  
```

Communication Avoiding/Reducing Algorithms for Sparse Linear Solvers utilizing Matrix Powers Kernel

- Matrix Powers Kernel: Ax , A^2x , A^3x ...
- Krylov Iterative Method without Preconditioning
 - Demmel, Hoemmen, Mohiyuddin etc. (UC Berkeley)
- *s-step* method
 - Just one P2P communication for each Mat-Vec during s iterations. Convergence may become unstable for large s .
- Communication Avoiding ILU0 (CA-ILU0) [Moufawad & Grigori, 2013]
 - First attempt to CA preconditioning
 - Nested dissection reordering for limited geometries (2D FDM)
- Generally, it is difficult to apply Matrix Powers Kernel to preconditioned iterative solvers

- Communication/Synchronization
Avoiding/Reducing in Krylov Iterative
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- **Pipelined CG: Background**
- Pipelined CG: Results
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Hiding Overhead by Collective Comm. in Krylov Iterative Solvers

- Dot Products in Krylov Iterative Solvers
 - MPI_Allreduce: Collective Communications
 - Large overhead with many nodes
- Pipelined CG [Ghysels et al. 2014]
 - Utilization of asynchronous collective communications (e.g. MPI_Iallreduce) supported in MPI-3 for hiding such overhead.
 - Algorithm is kept, but order of computations is changed
 - [Reference] P. Ghysels et al., Hiding global synchronization latency in the preconditioned Conjugate Gradient algorithm, Parallel Computing 40, 2014
 - When I visited LBNL in September 2013, Dr. Ghysels asked me to evaluate his idea in my parallel multigrid solvers

4 Algorithms [Ghysels et al. 2014]

- Alg.1 Original Preconditioned CG
- Alg.2 Chronopoulos/Gear
 - 2 dot products are combined in a single reduction
- Alg.3 Pipelined CG (MPI_allreduce)
- Alg.4 Gropp's asynchronous CG (MPI_allreduce)
- Algorithm itself is not different from the original one
 - Recurrence Relations: 漸化式
 - Order of computation changed -> Rounding errors are propagated differently
 - Convergence may be affected (not happened in my case)
 - update of $r = b - Ax$ needed at every 50 iterations (original paper)

Original Preconditioned CG (Alg.1)

Original Preconditioned CG (Alg.1)

```
1:  $r_0 := b - Ax_0; \quad u_0 := M^{-1}r_0; \quad p_0 := u_0$   
2: for  $i = 0, \dots$  do  
3:    $s := Ap_i$   
4:    $\alpha := (r_i, u_i) / (s, p_i)$   
5:    $x_{i+1} := x_i + \alpha p_i$   
6:    $r_{i+1} := r_i - \alpha s$   
7:    $u_{i+1} := M^{-1}r_{i+1}$   
8:    $\beta := (r_{i+1}, u_{i+1}) / (r_i, u_i)$   
9:    $p_{i+1} := u_{i+1} + \beta p_i$   
10: end for
```

$$s_i = Ap_i$$

$$x_{i+1} = x_i + \alpha_i p_i$$

$$\begin{aligned} r_{i+1} &= b - Ax_{i+1} = b - Ax_i - \alpha_i Ap_i \\ &= r_i - \alpha_i Ap_i = r_i - \alpha_i s_i \end{aligned}$$

Chronopoulos/Gear CG (Alg.2)

2 dot products are combined into a single reduction
s-step algorithm with s=1

Chronopoulos/Gear CG (Alg.2)

```

1:  $r_0 := b - Ax_0; \quad u_0 := M^{-1}r_0; \quad w_0 := Au_0$ 
2:  $\alpha_0 := (r_0, u_0) / (w_0, u_0); \quad \beta_0 := 0; \quad \gamma_0 := (r_0, u_0)$ 
3: for  $i = 0, \dots$  do
4:    $p_i := u_i + \beta_i p_{i-1}$ 
5:    $s_i := w_i + \beta_i s_{i-1}$ 
6:    $x_{i+1} := x_i + \alpha_i p_i$ 
7:    $r_{i+1} := r_i - \alpha_i s_i$ 
8:    $u_{i+1} := M^{-1}r_{i+1}$ 
9:    $w_{i+1} := Au_{i+1}$ 
10:   $\gamma_{i+1} := (r_{i+1}, u_{i+1})$ 
11:   $\delta := (w_{i+1}, u_{i+1})$ 
12:   $\beta_{i+1} := \gamma_{i+1} / \gamma_i$ 
13:   $\alpha_{i+1} := \gamma_{i+1} / (\delta - \beta_{i+1} \gamma_{i+1} / \alpha_i)$ 
14: end for

```

$$s_i = Au_i + \beta_i s_{i-1}, \quad p_i = u_i + \beta_i p_{i-1} \Leftrightarrow s_i = Ap_i$$

$$x_{i+1} = x_i + \alpha_i p_i$$

$$r_{i+1} = b - Ax_{i+1} = b - Ax_i - \alpha_i Ap_i$$

$$= r_i - \alpha_i Ap_i = r_i - \alpha_i s_i$$

- 2 dot products combined

- $s_i = Ap_i$ is not computed explicitly: by recurrence

Pipelined Chronopoulos/Gear (No Preconditioning)

Pipelined Chronopoulos/Gear (No Precond.)

```

1:  $r_0 := b - Ax_0; \quad w_0 := Ar_0$ 
2: for  $i = 0, \dots$  do
3:    $\gamma_i := (r_i, r_i)$ 
4:    $\delta := (w_i, r_i)$ 
5:    $q_i := Aw_i$ 
6:   if  $i > 0$  then
7:      $\beta_i := \gamma_i / \gamma_{i-1}; \quad \alpha_i := \gamma_i / (\delta - \beta_i \gamma_i / \alpha_{i-1})$ 
8:   else
9:      $\beta_i := 0; \quad \alpha_i := \gamma_i / \delta$ 
10:  end if
11:   $z_i := q_i + \beta_i z_{i-1}$ 
12:   $s_i := w_i + \beta_i s_{i-1}$ 
13:   $p_i := r_i + \beta_i p_{i-1}$ 
14:   $x_{i+1} := x_i + \alpha_i p_i$ 
15:   $r_{i+1} := r_i - \alpha_i s_i$ 
16:   $w_{i+1} := w_i - \alpha_i z_i$ 
17: end for

```

- Global synchronization of dot products are overlapped with SpMV

$$\begin{aligned}
 u_i &= r_i, \quad w_i = Au_i = Ar_i \\
 Ar_{i+1} &= Ar_i - \alpha_i As_i \Rightarrow w_{i+1} = w_i - \alpha_i As_i \\
 As_i &= Aw_i + \beta_i As_{i-1} \Rightarrow z_i (= As_i = A^2 p_i) = Aw_i + \beta_i z_{i-1} \\
 q_i &= Aw_i = A^2 r_i \\
 &= r_i - \alpha_i Ap_i = r_i - \alpha_i s_i
 \end{aligned}$$

Preconditioned Pipelined CG (Alg.3)

Preconditioned Pipelined CG (Alg.3)

```

1:  $r_0 := b - Ax_0; \quad u_0 := M^{-1}r_0; \quad w_0 := Au_0$ 
2: for  $i = 0, \dots$  do
3:    $\gamma_i := (r_i, u_i)$ 
4:    $\delta := (w_i, u_i)$ 
5:    $m_i := M^{-1}w_i$ 
6:    $n_i := Am_i$ 
7:   if  $i > 0$  then
8:      $\beta_i := \gamma_i / \gamma_{i-1}; \quad \alpha_i := \gamma_i / (\delta - \beta_i \gamma_i / \alpha_{i-1})$ 
9:   else
10:     $\beta_i := 0; \quad \alpha_i := \gamma_i / \delta$ 
11:   end if
12:    $z_i := n_i + \beta_i z_{i-1}$ 
13:    $q_i := m_i + \beta_i q_{i-1}$ 
14:    $s_i := w_i + \beta_i s_{i-1}$ 
15:    $p_i := u_i + \beta_i p_{i-1}$ 
16:    $x_{i+1} := x_i + \alpha_i p_i$ 
17:    $r_{i+1} := r_i - \alpha_i s_i$ 
18:    $u_{i+1} := u_i - \alpha_i q_i$ 
19:    $w_{i+1} := w_i - \alpha_i z_i$ 
20: end for

```

- Global synchronization of dot products are overlapped with SpMV and Preconditioning

$$\begin{aligned}
 \gamma_i &= (u_i, u_i)_M = (Mu_i, u_i) = (r_i, u_i) \\
 \delta_i &= (M^{-1}Au_i, u_i)_M = (Au_i, u_i) = (w_i, u_i) \\
 M^{-1}r_{i+1} &= M^{-1}r_i - \alpha_i M^{-1}s_i \Rightarrow u_{i+1} = u_i - \alpha_i q_i \quad (q_i = M^{-1}s_i) \\
 M^{-1}s_{i+1} &= M^{-1}w_i + \beta_i M^{-1}s_i \Rightarrow q_{i+1} = M^{-1}w_i + \beta_i q_i \\
 Au_{i+1} &= Au_i - \alpha_i Aq_i \Rightarrow w_{i+1} = w_i - \alpha_i Aq_i \\
 Aq_i &= AM^{-1}w_i + \beta_i Aq_{i-1} \Rightarrow z_i = Am_i + \beta_i z_{i-1} \\
 (m_i &= M^{-1}w_i = M^{-1}Au_i = M^{-1}AM^{-1}r_i, z_i = Aq_i)
 \end{aligned}$$

Gropp's Asynchronous CG (Alg.4)

Smaller Computations than Alg.3

Gropp's Asynchronous CG (Alg.4)

```

1:  $r_0 := b - Ax_0; \quad u_0 := M^{-1}r_0; \quad p_0 := u_0; \quad s_0 := Ap_0; \quad \gamma_0 := (r_0, u_0)$ 
2: for  $i = 0, \dots$  do
3:    $\delta := (p_i, s_i)$ 
4:    $q_i := M^{-1}s_i$ 
5:    $\alpha_i := \gamma_i / \delta$ 
6:    $x_{i+1} := x_i + \alpha_i p_i$ 
7:    $r_{i+1} := r_i - \alpha_i s_i$ 
8:    $u_{i+1} := u_i - \alpha_i q_i$ 
9:    $\gamma_{i+1} := (r_{i+1}, u_{i+1})$ 
10:   $w_{i+1} := Au_{i+1}$ 
11:   $\beta_{i+1} := \gamma_{i+1} / \gamma_i$ 
12:   $p_{i+1} := u_{i+1} + \beta_{i+1} p_i$ 
13:   $s_{i+1} := w_{i+1} + \beta_{i+1} s_i$ 
14: end for

```

- Definition of δ is different from that of Alg.3
- Global synchronization of dot products are overlapped with SpMV and Preconditioning

W. Gropp, Update on Libraries for Blue Waters.

<http://jointlab-pc.ncsa.illinois.edu/events/workshop3/pdf/presentations/Gropp-Update-on-Libraries.pdf>

Presentation Material (not a paper, article)



Implementation (Alg.4)

```

!C
!C +-----+
!C | Delta= (p, s) |
!C +-----+
!C===
      DLO= 0.d0
!$omp parallel do private(i) reduction(+:DLO)
      do i= 1, 3*N
        DLO= DLO + P(i)*S(i)
      enddo
      call MPI_Iallreduce (DLO, Delta, 1, MPI_DOUBLE_PRECISION,
&                          MPI_SUM, MPI_COMM_WORLD, req1, ierr)
!C===

!C
!C +-----+
!C | {q}= [Minv] {s} |
!C +-----+
!C===

      (前处理：省略)

!C===

      call MPI_Wait (req1, stat1, ierr)

```

Gropp's Asynchronous CG (Alg.4)

```

1:  $r_0 := b - Ax_0$ ;  $u_0 := M^{-1}r_0$ ;  $p_0 := u_0$ ;  $s_0 := Ap_0$ ;  $\gamma_0 := (r_0, u_0)$ 
2: for  $i = 0, \dots$  do
3:    $\delta := (p_i, s_i)$ 
4:    $q_i := M^{-1}s_i$ 
5:    $\alpha_i := \gamma_i / \delta$ 
6:    $x_{i+1} := x_i + \alpha_i p_i$ 
7:    $r_{i+1} := r_i - \alpha_i s_i$ 
8:    $u_{i+1} := u_i - \alpha_i q_i$ 
9:    $\gamma_{i+1} := (r_{i+1}, u_{i+1})$ 
10:   $w_{i+1} := Au_{i+1}$ 
11:   $\beta_{i+1} := \gamma_{i+1} / \gamma_i$ 
12:   $p_{i+1} := u_{i+1} + \beta_{i+1} p_i$ 
13:   $s_{i+1} := w_{i+1} + \beta_{i+1} s_i$ 
14: end for

```

Pipelined CR (Conjugate Residuals)

Preconditioned Pipelined CR

```

1:  $r_0 := b - Ax_0; \quad u_0 := M^{-1}r_0; \quad w_0 := Au_0$ 
2: for  $i = 0, \dots$  do
3:    $m_i := M^{-1}w_i$ 
4:    $\gamma_i := (w_i, u_i)$ 
5:    $\delta := (m_i, w_i)$ 
6:    $n_i := Am_i$ 
7:   if  $i > 0$  then
8:      $\beta_i := \gamma_i / \gamma_{i-1}; \quad \alpha_i := \gamma_i / (\delta - \beta_i \gamma_i / \alpha_{i-1})$ 
9:   else
10:     $\beta_i := 0; \quad \alpha_i := \gamma_i / \delta$ 
11:  end if
12:   $z_i := n_i + \beta_i z_{i-1}$ 
13:   $q_i := m_i + \beta_i q_{i-1}$ 
14:   $p_i := u_i + \beta_i p_{i-1}$ 
15:   $x_{i+1} := x_i + \alpha_i p_i$ 
16:   $u_{i+1} := u_i - \alpha_i q_i$ 
17:   $w_{i+1} := w_i - \alpha_i z_i$ 
18: end for

```

- Global synchronization of dot products are overlapped with SpMV

Amount of Computations

DAXPY could very smaller than (sophisticated) preconditioning

		SpMV	Precond.	Dot Prod.	DAXPY
1	Original CG	1	1	2+1	3
2	Chronopoulos/ Gear	1	1	2+1	4
3	Pipelined CG	1	1	2+1	8
4	Grppp's Algorithm	1	1	2+1	5
	Pipelined CR	1	1	2+1	6

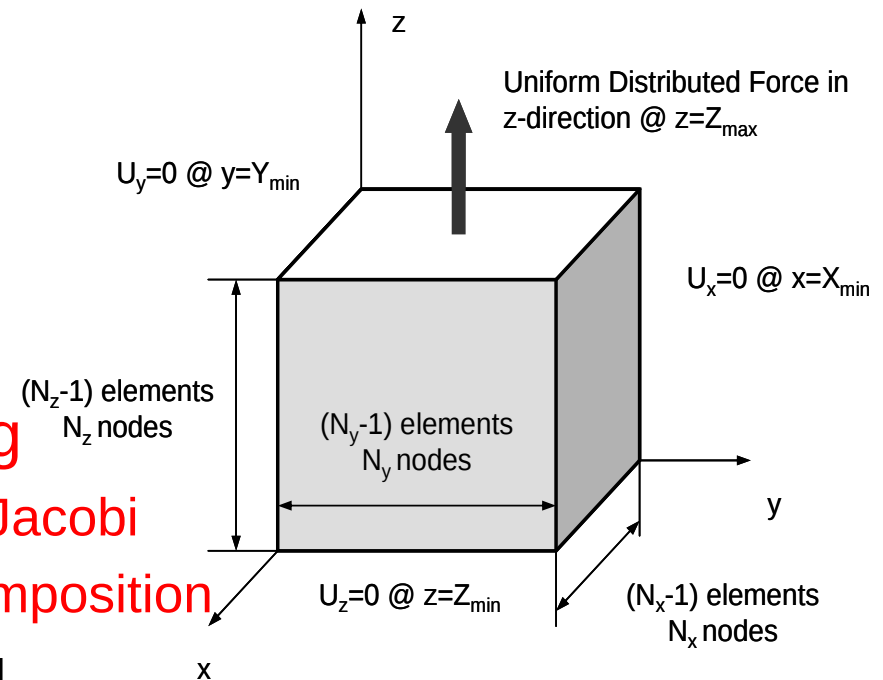
+1 for
residual norm

- Communication/Synchronization
Avoiding/Reducing in Krylov Iterative
Solvers
- Pipelined CG: Background
- **Pipelined CG: Results**
- Communication-Computation
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GeoFEM/Cube

- Parallel FEM Code (& Benchmarks)
- 3D-Static-Elastic-Linear (Solid Mechanics)
- Performance of Parallel Precond. Iterative Solvers
 - 3D Tri-linear Elements
 - SPD matrices: CG solver
 - Fortran90+MPI+OpenMP
 - Distributed Data Structure
 - **Localized SGS Preconditioning**
 - Symmetric Gauss-Seidel, Block Jacobi
 - Additive Schwartz Domain Decomposition
 - Reordering by CM-RCM: RCM
 - MPI, OpenMP, OpenMP/MPI Hybrid



Overlapped Additive Schwarz Domain Decomposition Method



Stabilization of Localized Preconditioning: ASDD

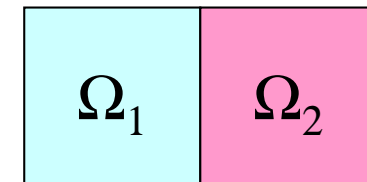
Global Operation

$$Mz = r$$



Local Operation

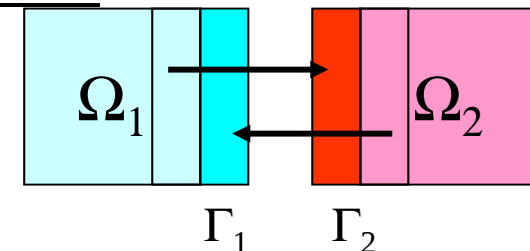
$$z_{\Omega_1} = M_{\Omega_1}^{-1} r_{\Omega_1}, \quad z_{\Omega_2} = M_{\Omega_2}^{-1} r_{\Omega_2}$$



Global Nesting Correction: Repeating -> Stable

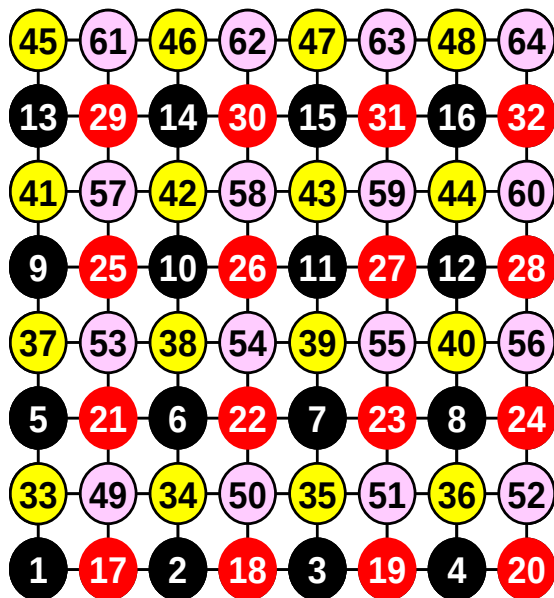
$$z_{\Omega_1}^n = z_{\Omega_1}^{n-1} + M_{\Omega_1}^{-1} (r_{\Omega_1} - M_{\Omega_1} z_{\Omega_1}^{n-1} - M_{\Gamma_1} z_{\Gamma_1}^{n-1})$$

$$z_{\Omega_2}^n = z_{\Omega_2}^{n-1} + M_{\Omega_2}^{-1} (r_{\Omega_2} - M_{\Omega_2} z_{\Omega_2}^{n-1} - M_{\Gamma_2} z_{\Gamma_2}^{n-1})$$

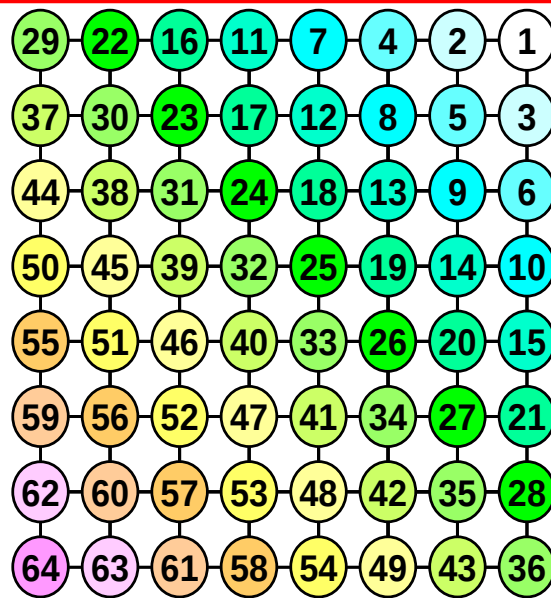


Reordering for avoiding data dependency in IC/ILU computations on each MPI process

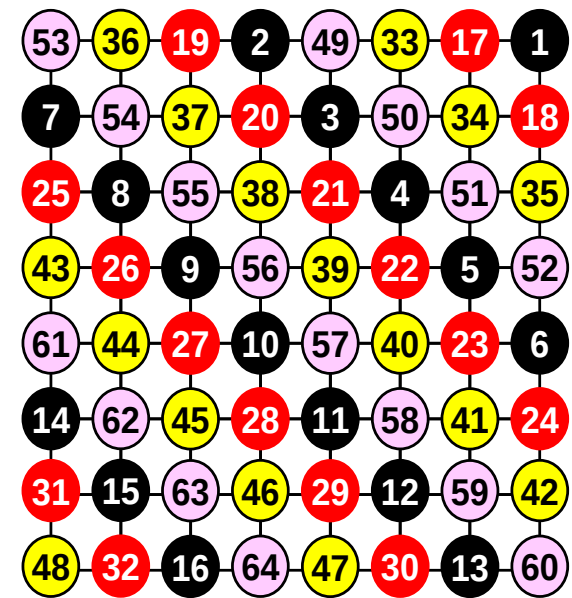
Elements in “same color” are independent: to be parallelized by OpenMP on each MPI process.



**MC (Color#=4)
Multicoloring**



**RCM
Reverse Cuthill-McKee**



**CM-RCM (Color#=4)
Cyclic MC + RCM**

Results on Reedbush-U (RDB)

- 4 Types of Algorithms
 - Alg.1 Original Preconditioned CG
 - Alg.2 Chronopoulos/Gear
 - Alg.3 Pipelined CG (MPI_Allreduce, MPI_Iallreduce,)
 - Alg.4 Gropp's asynchronous CG (MPI_Allreduce, MPI_Iallreduce)
- Flat MPI, OpenMP/MPI Hybrid with Reordering
- Platform
 - Integrated Supercomputer System for Data Analyses & Scientific Simulations (Reedbush-U)
 - Intel Broadwell-EP 18 cores x 2 sockets x 420 nodes
 - Intel Fortran + Intel MPI
 - 16 of 18 cores/socket, up to 384 nodes (= 768 sockets, 12,288 cores)

Results: Number of Iterations

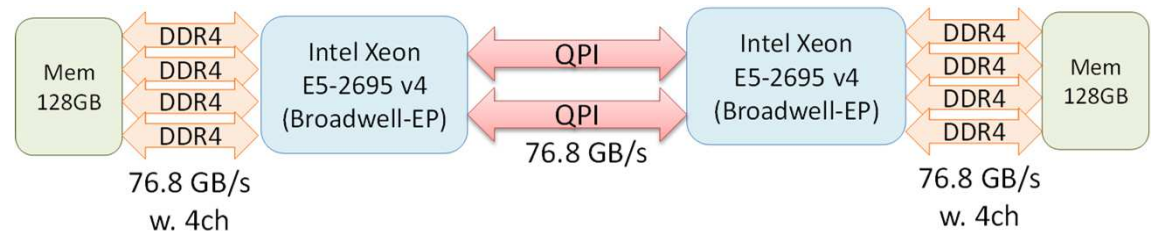
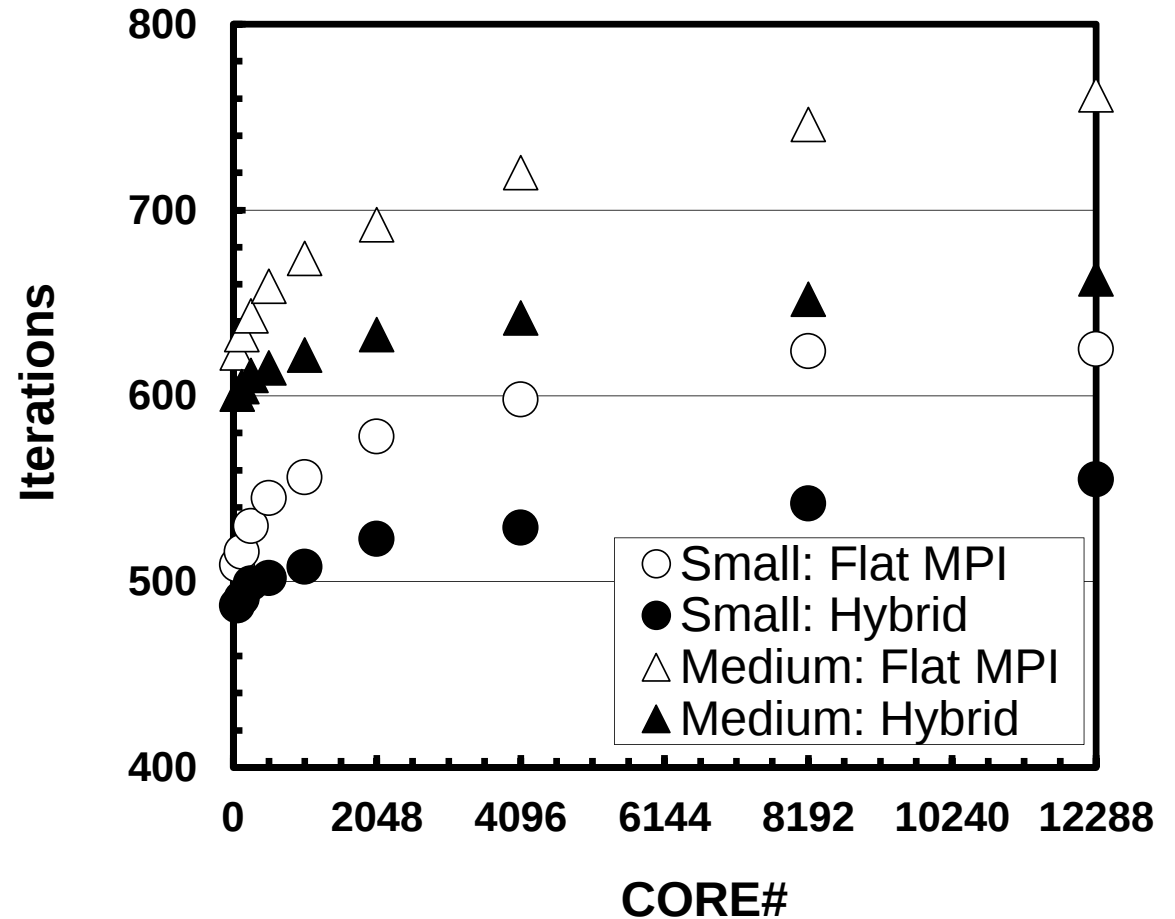
- Strong Scaling

- Small

- $256 \times 128 \times 144$ nodes (=4,718,592)
- 14,155,776 DOF
- at 384 nodes (12,288 cores)
 - $8 \times 8 \times 6 = 384$ nodes/core
 - 1,152 DOF/core

- Medium

- $256 \times 128 \times 288$ nodes (=9,437,184)
- 28,311,552 自由度



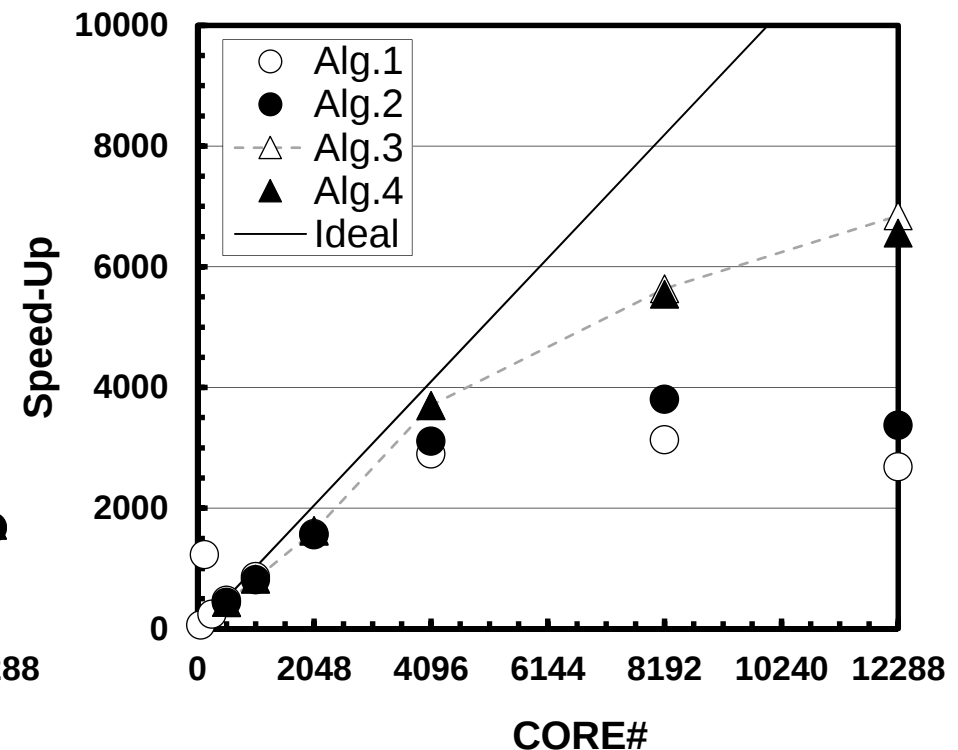
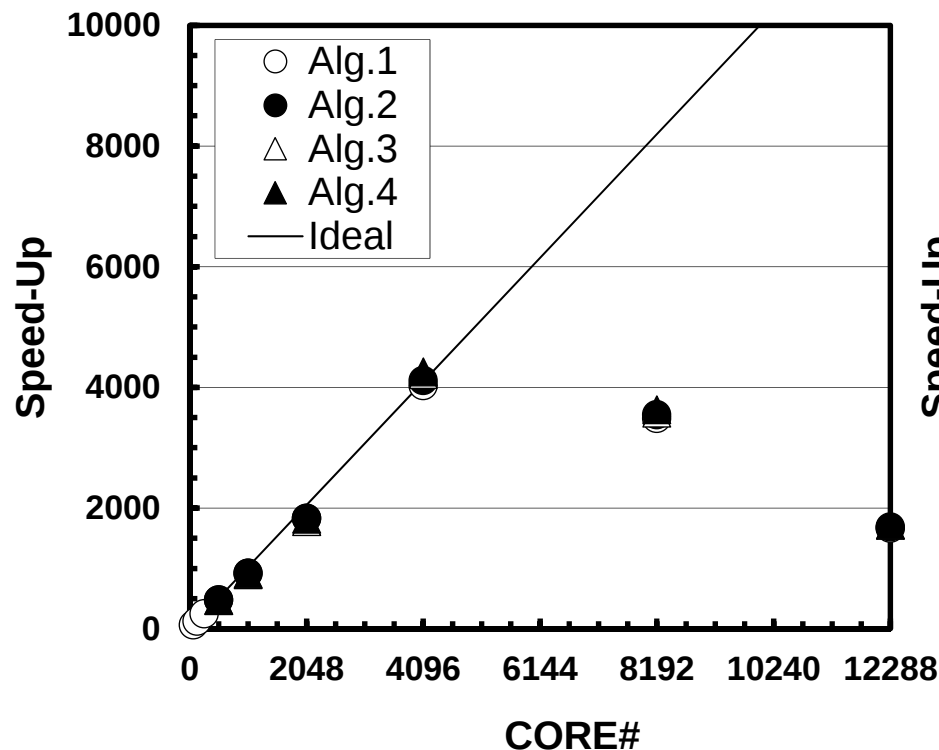
Results: Speed-Up: Small

Performance of 2 nodes of Flat MPI = 64.0 (4 sockets, 64 cores)

Flat MPI

Hybrid

Alg.1	Original PCG
Alg.2	Chronopoulos/Gear
Alg.3	Pipelined CG
Alg.4	Gropp's CG



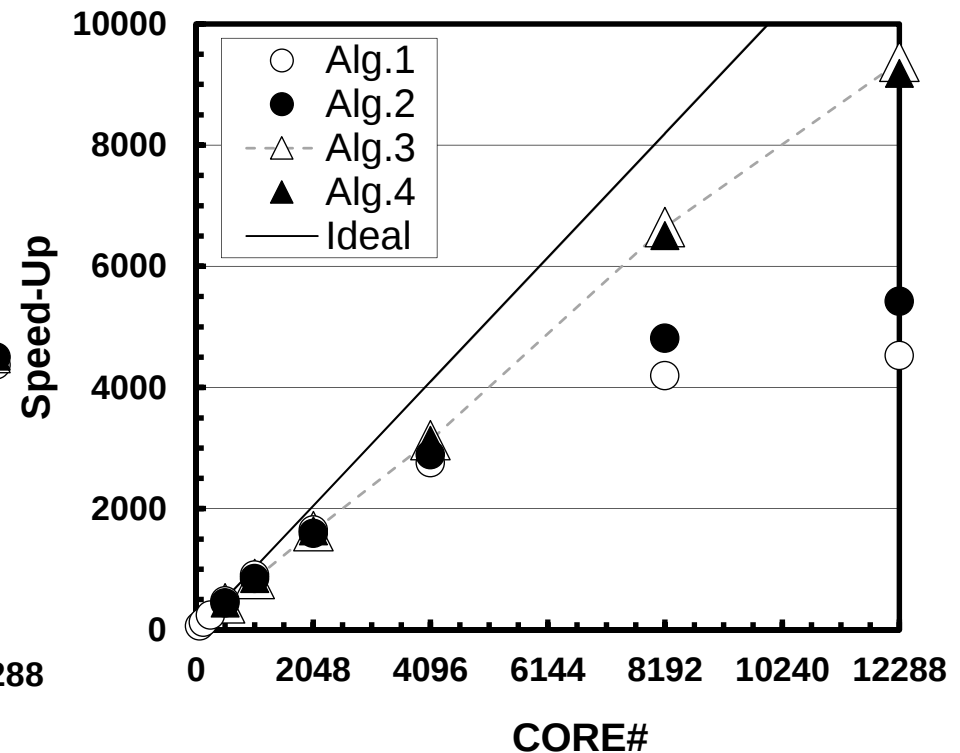
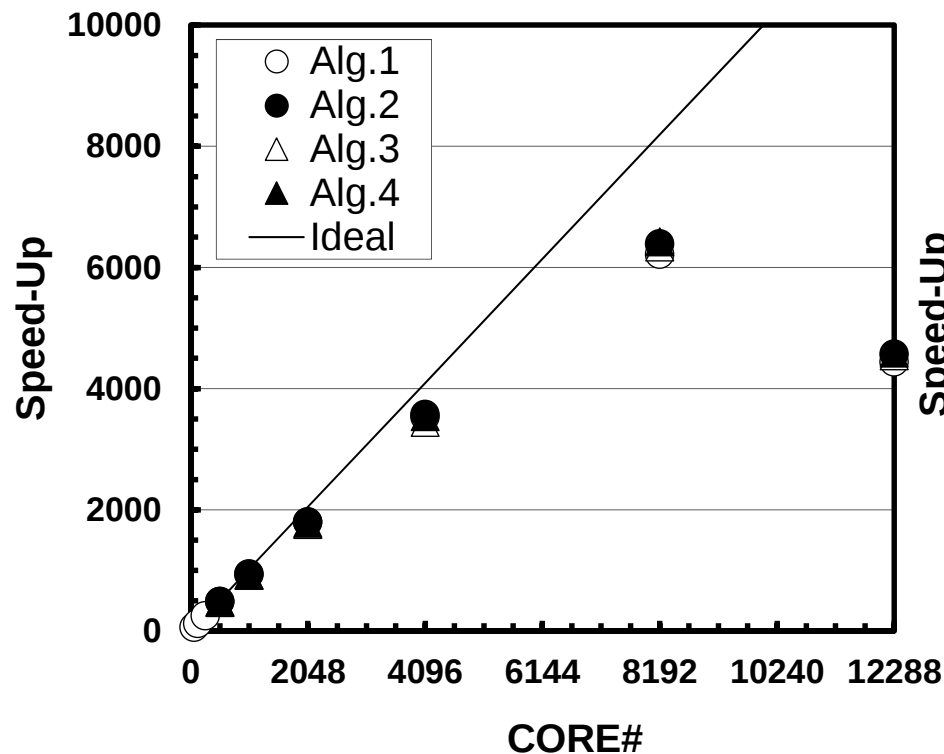
Results: Speed-Up: Medium

Performance of 2 nodes of Flat MPI = 64.0 (4 sockets, 64 cores)

Flat MPI

Hybrid

Alg.1	Original PCG
Alg.2	Chronopoulos/Gear
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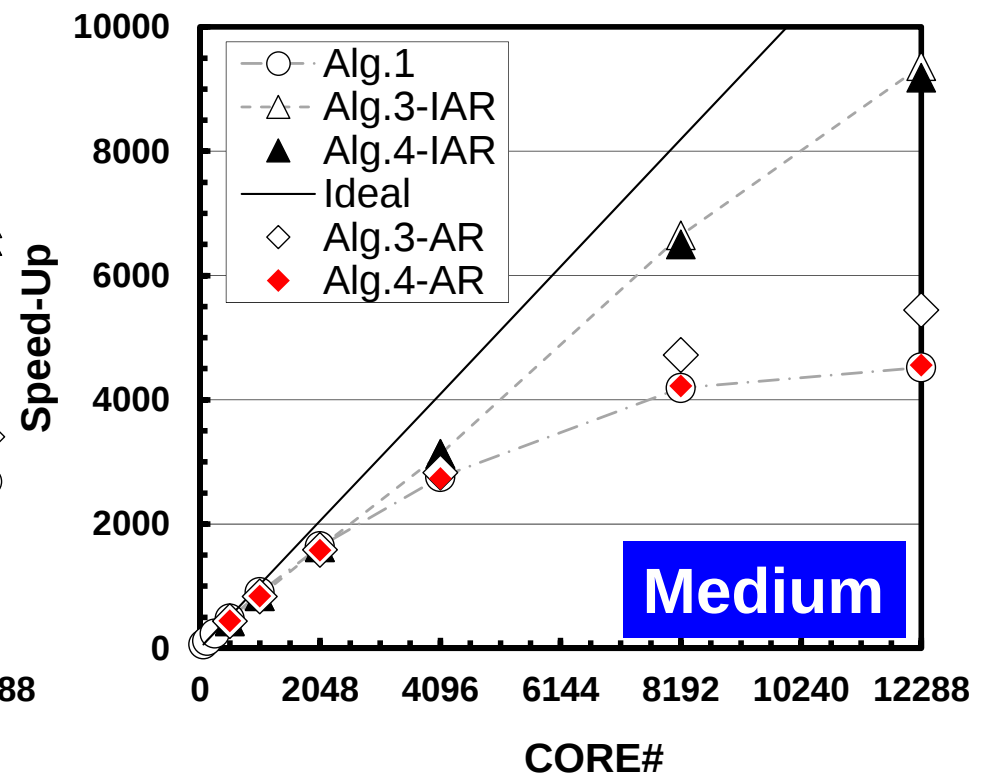
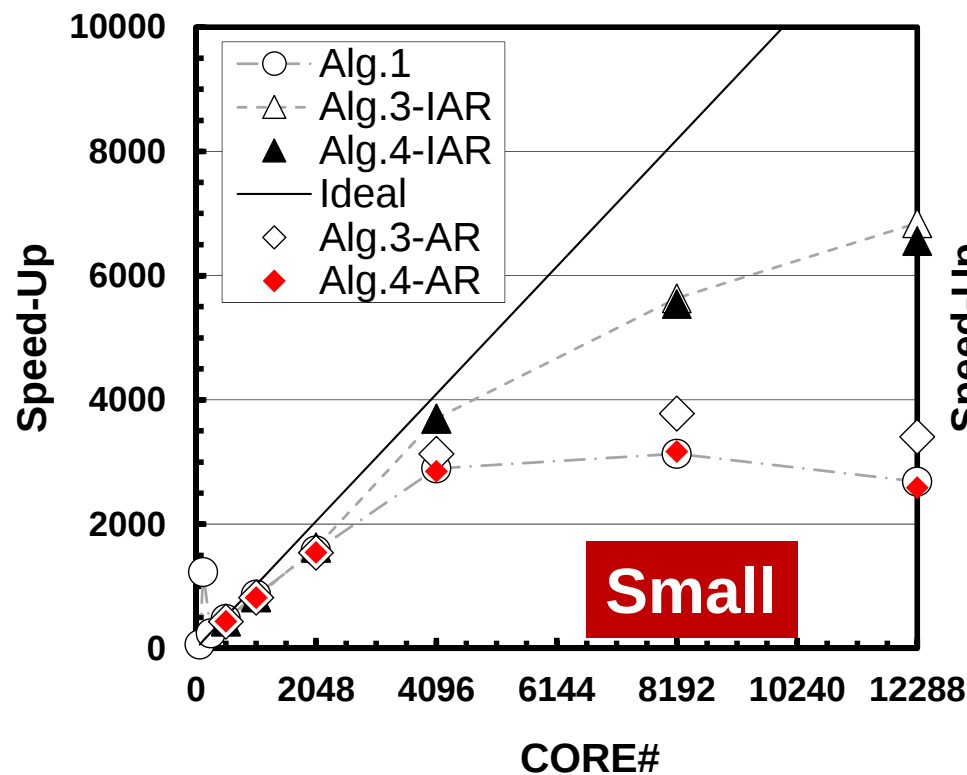


Allreduce vs. lallreduce for Hybrid

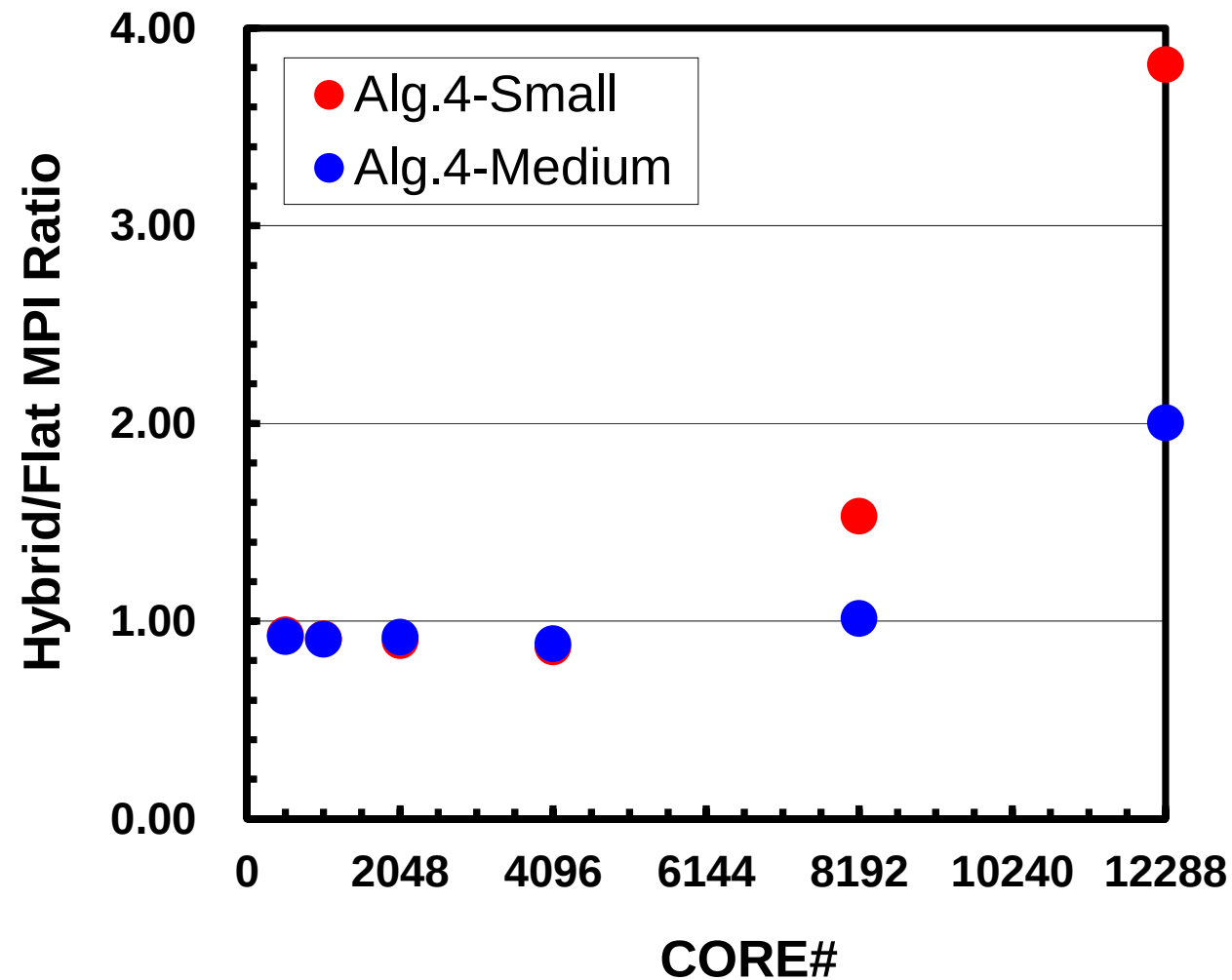
Performance of 2 nodes of Flat MPI = 64.0 (4 sockets, 64 cores)

IAR: MPI_lallreduce
AR : MPI_Allreduce

Alg.1 Original PCG
Alg.3 Pipelined CG
Alg.4 Gropp's CG



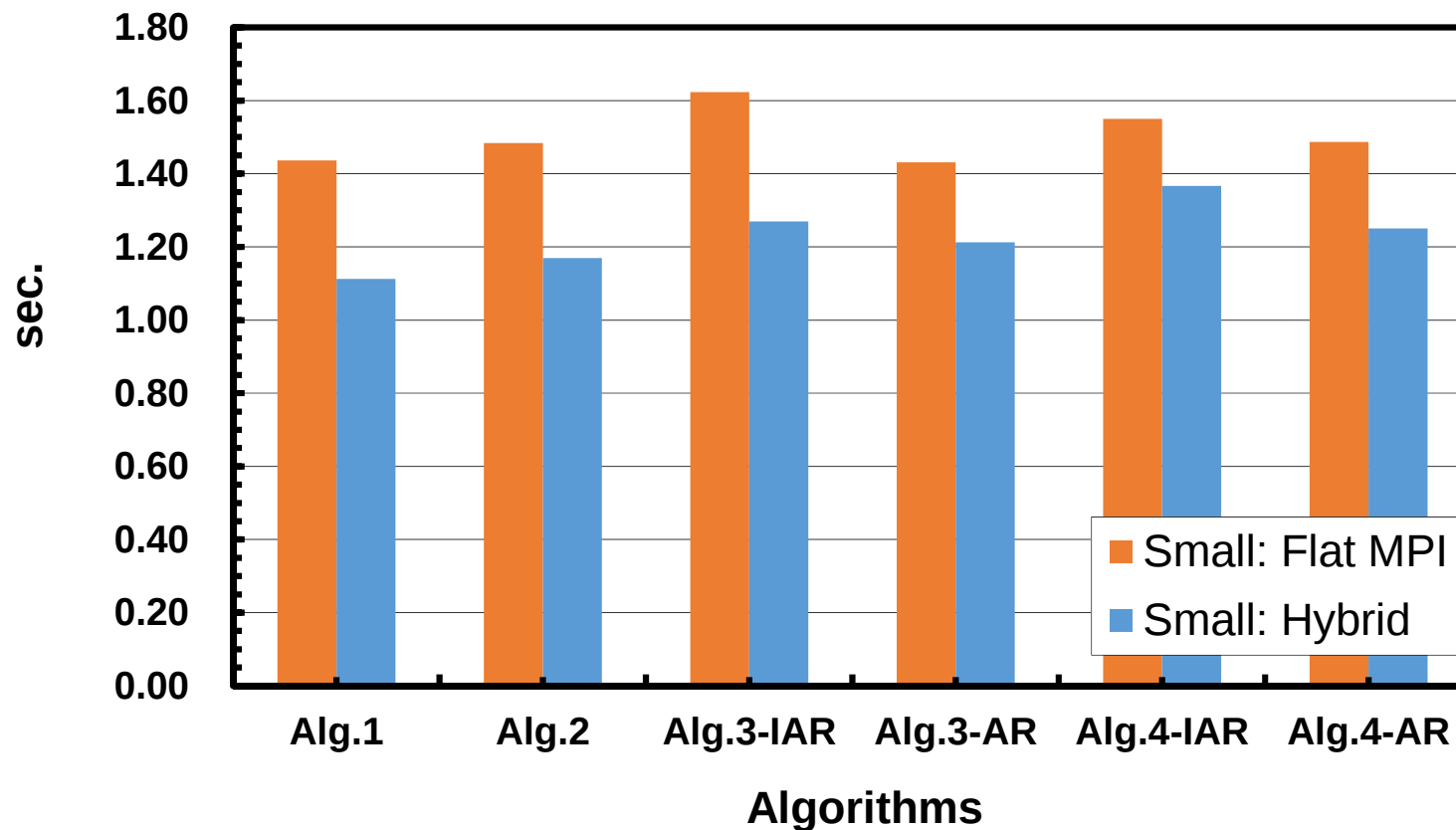
Hybrid vs. Flat MPI for Alg.4



Results on 768 nodes (12,288 cores) of Fujitsu FX10 (Oakleaf-FX)

MPI-3 is not optimized

FX100 has special HW for communication



Summary

- Pipelined CG, Gropp's CG
 - Effect of hiding collective communication by MPI_allreduce is significant, especially for strong scaling
 - Alg.3 ~ Alg.4
 - Future works
 - Pipelined CR should be also evaluated
 - Dr. Ghysels's recommendation
 - Application to Multigrid, HID (Hierarchical Interface Decomposition)
 - Evaluation on FX100, Oakforest-PACS (KNL Cluster) with OmniPATH
- Loop Scheduling for OpenMP
 - Effect is significant on FX10
 - Detailed profiling needed
 - Good target for AT (already done)