

Introduction to Parallel FEM in Fortran Parallel Data Structure

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Programming for Parallel Computing (616-2057)
Seminar on Advanced Computing (616-4009)

Parallel Computing

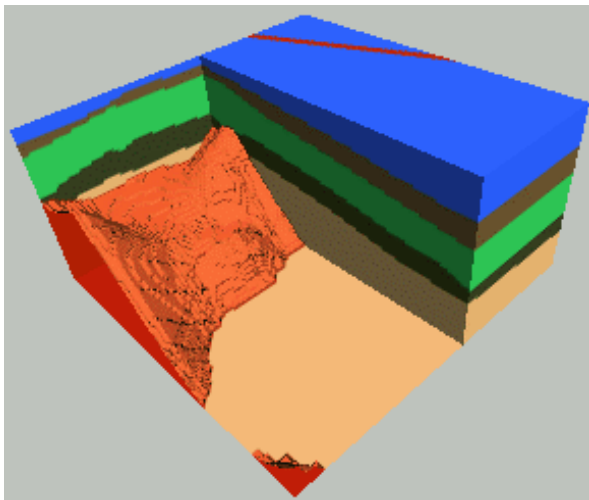
- Faster, Larger & More Complicated
- Scalability
 - Solving N^x scale problem using N^x computational resources during same computation time
 - for large-scale problems: Weak Scaling
 - e.g. CG solver: more iterations needed for larger problems
 - Solving a problem using N^x computational resources during $1/N$ computation time
 - for faster computation: Strong Scaling

What is Parallel Computing ? (1/2)

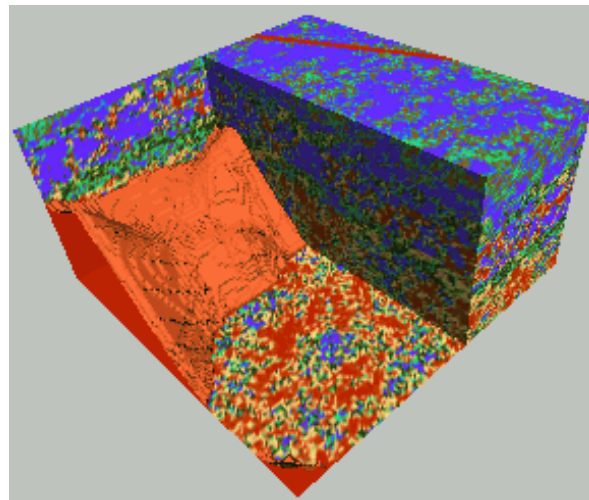
- to solve larger problems faster

Homogeneous/Heterogeneous Porous Media

Lawrence Livermore National Laboratory



Homogeneous

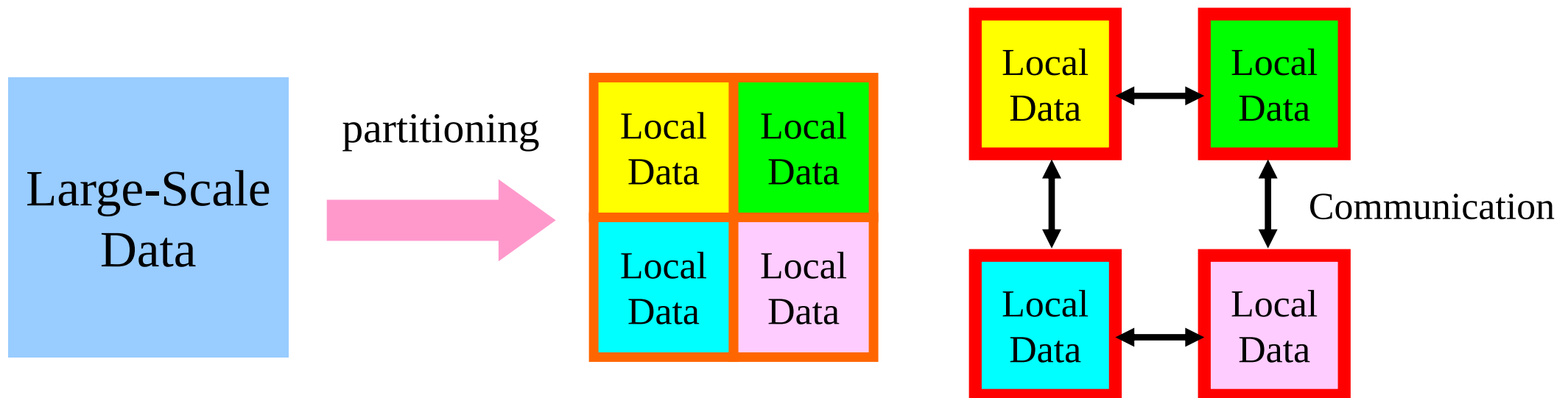


Heterogeneous

very fine meshes are required for simulations of heterogeneous field.

What is Parallel Computing ? (2/2)

- PC with 1GB memory : 1M meshes are the limit for FEM
 - Southwest Japan with 1,000km x 1,000km x 100km in 1km mesh
-> 10^8 meshes
- Large Data -> Domain Decomposition -> Local Operation
- Inter-Domain Communication for Global Operation



What is Communication ?

- Parallel Computing -> Local Operations
- Communications are required in Global Operations for Consistency.

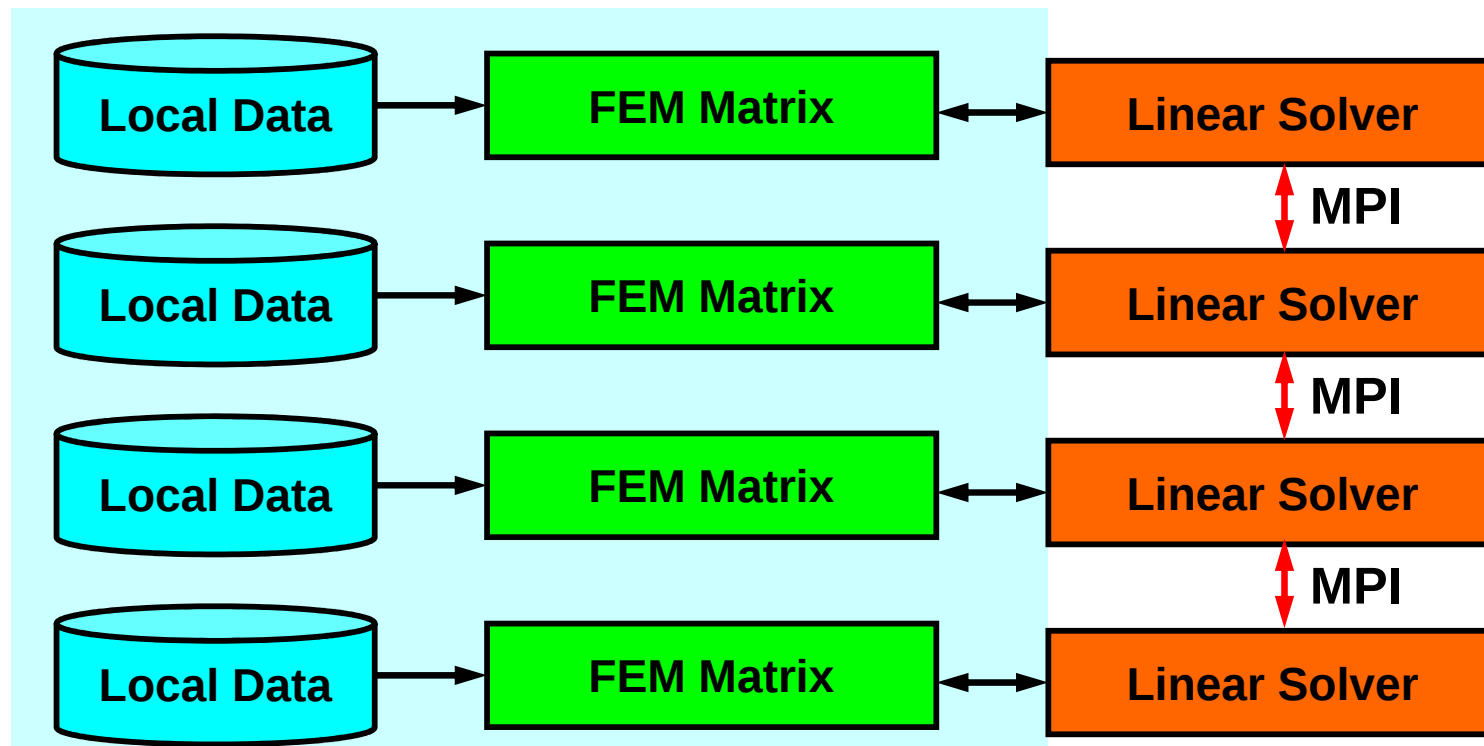
Operations in Parallel FEM

SPMD: Single-Program Multiple-Data

Large Scale Data -> partitioned into Distributed Local Data Sets.

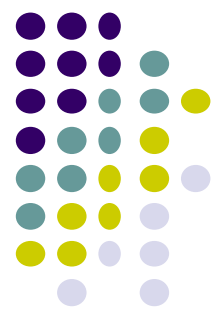
FEM code can assemble coefficient matrix for each local data set :
this part could be completely local, same as serial operations

Global Operations & Communications happen only in Linear Solvers
dot products, matrix-vector multiply, preconditioning



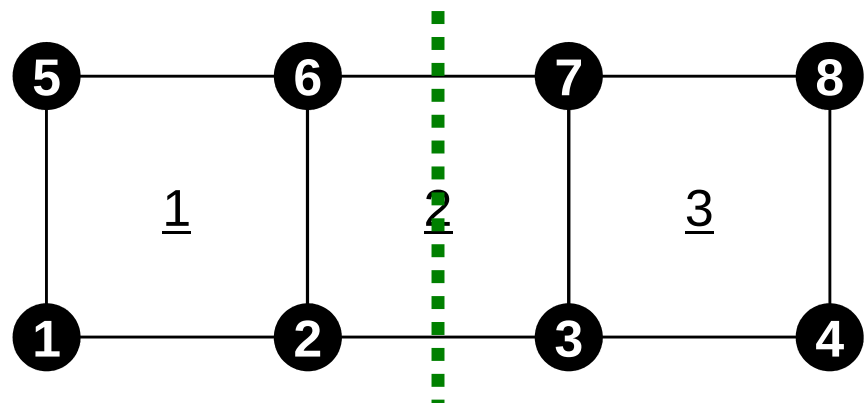
Parallel FEM Procedures

- Design on “Local Data Structure” is important
 - for SPMD-type operations in the previous page
- Matrix Generation
- Preconditioned Iterative Solvers for Linear Equations



Bi-Linear Square Elements

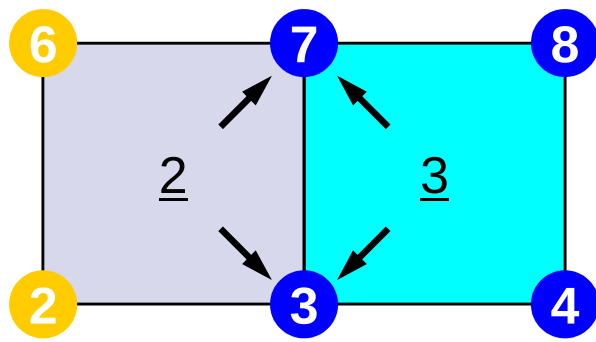
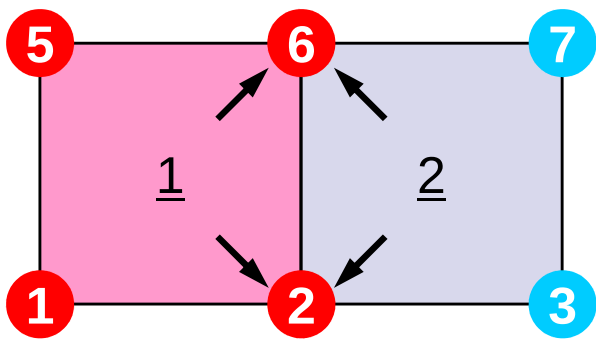
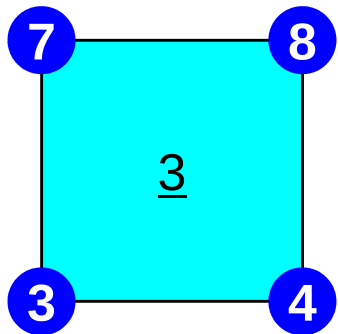
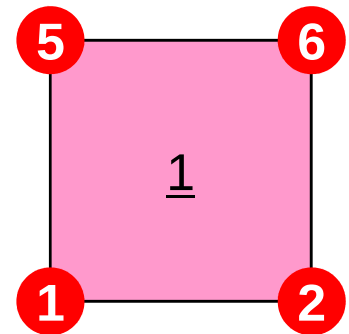
Values are defined on each node

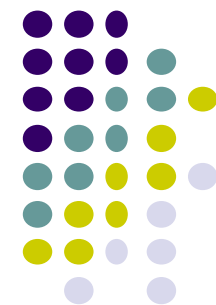


divide into two domains by “node-based” manner, where number of “nodes (vertices)” are balanced.

Local information is not enough for matrix assembling.

Information of overlapped elements and connected nodes are required for matrix assembling on boundary nodes.



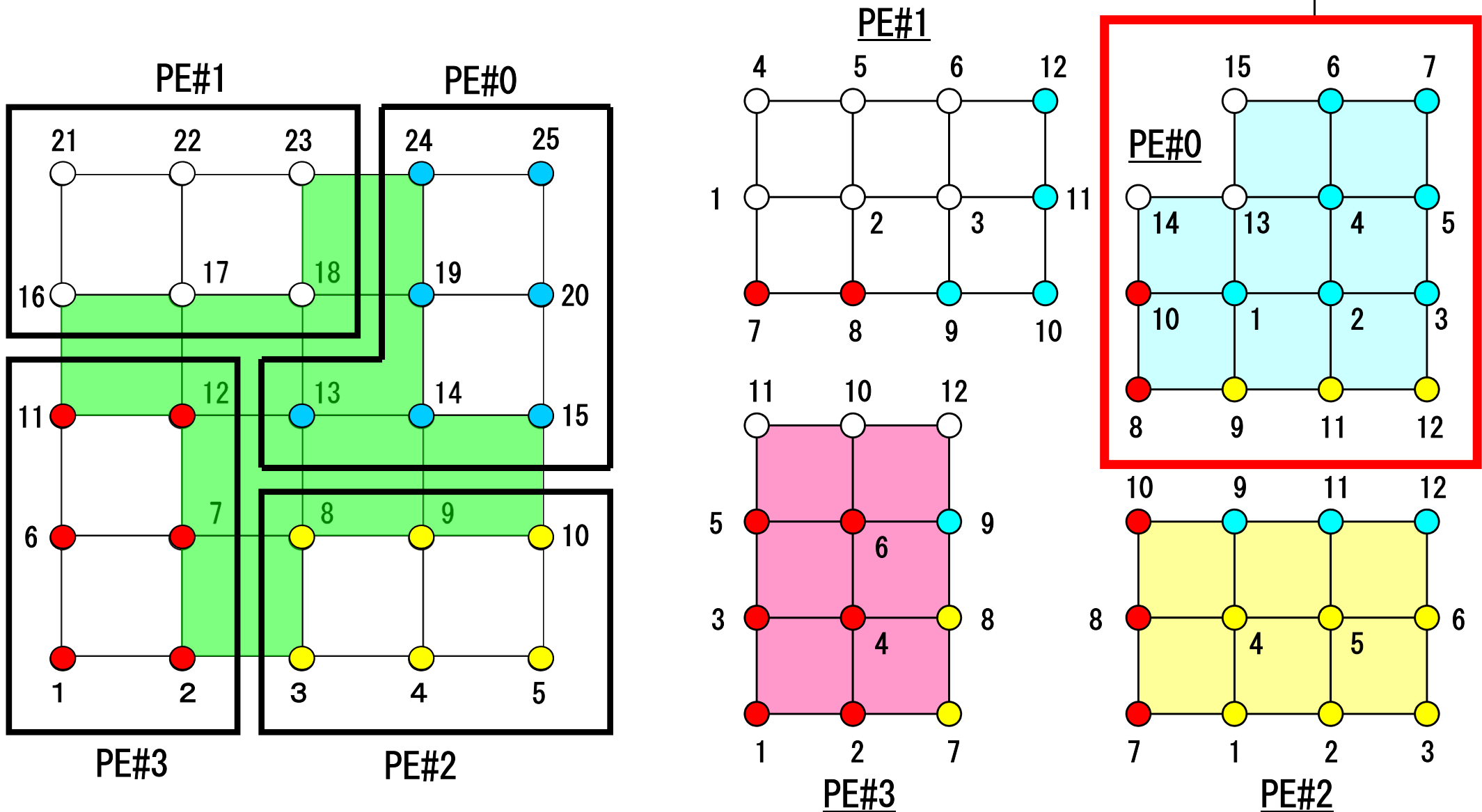


Local Data of Parallel FEM

- **Node-based partitioning for IC/ILU type preconditioning methods**
- Local data includes information for :
 - Nodes originally assigned to the partition/PE
 - Elements which include the nodes : Element-based operations (Matrix Assemble) are allowed for fluid/structure subsystems.
 - All nodes which form the elements but out of the partition
- Nodes are classified into the following 3 categories from the viewpoint of the message passing
 - **Internal nodes** originally assigned nodes
 - **External nodes** in the overlapped elements but out of the partition
 - **Boundary nodes** *external nodes* of other partition
- Communication table between partitions
- NO global information required except partition-to-partition connectivity

Node-based Partitioning

internal nodes - elements - external nodes

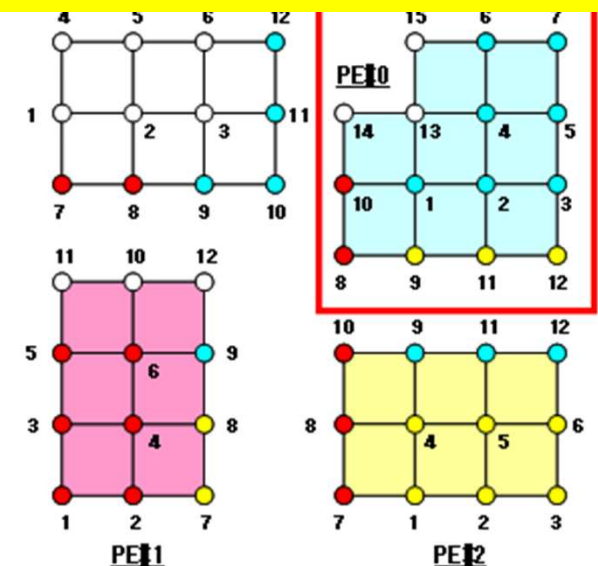
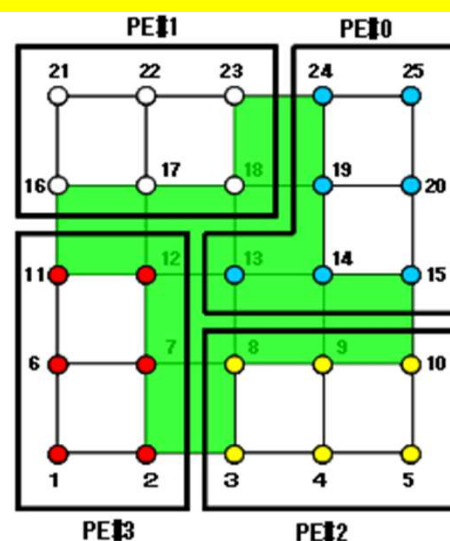
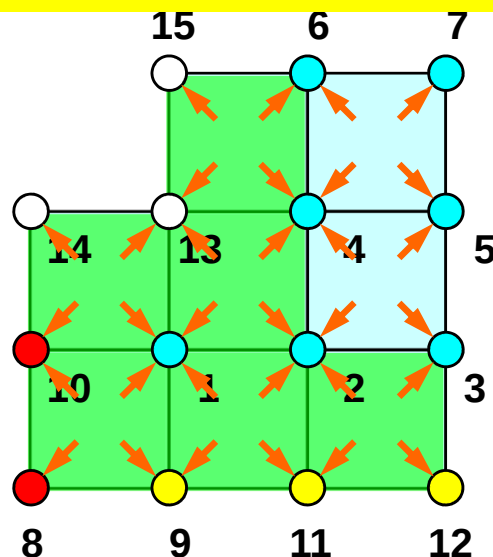


Node-based Partitioning

internal nodes - elements - external nodes

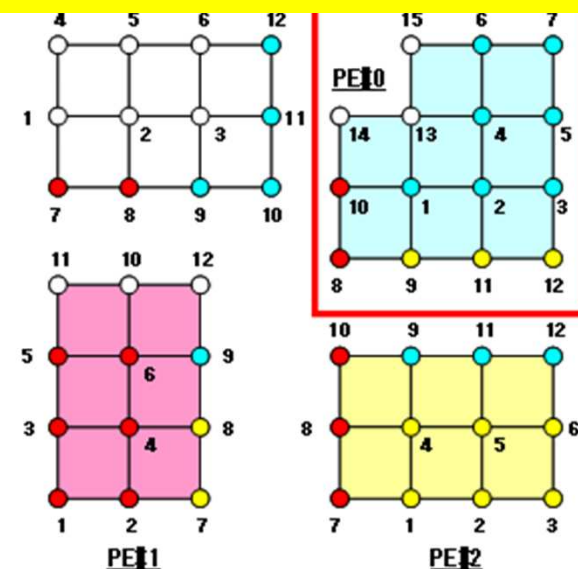
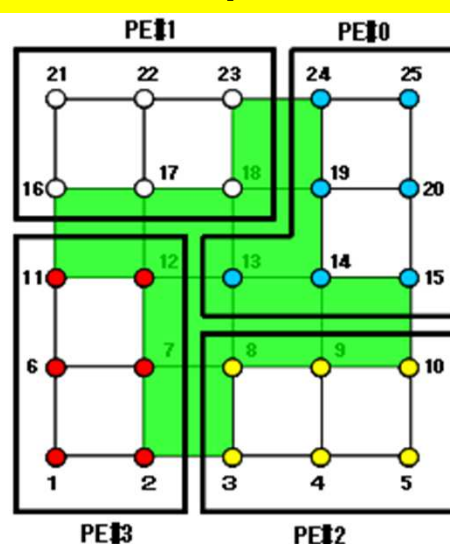
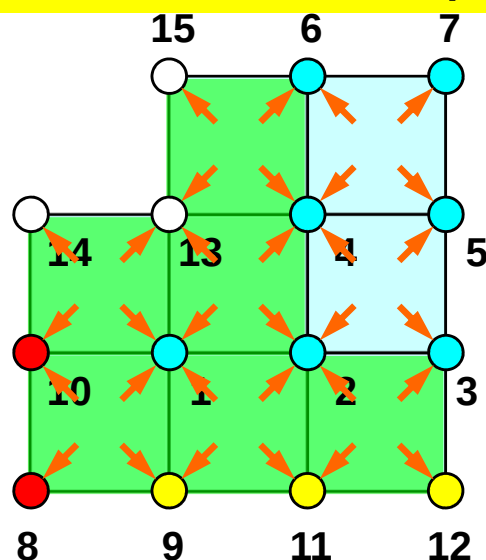


- Partitioned nodes themselves (Internal Nodes) 内点
- Elements which include Internal Nodes 内点を含む要素
- External Nodes included in the Elements 外点
in overlapped region among partitions.
- Info of External Nodes are required for completely local element-based operations on each processor.



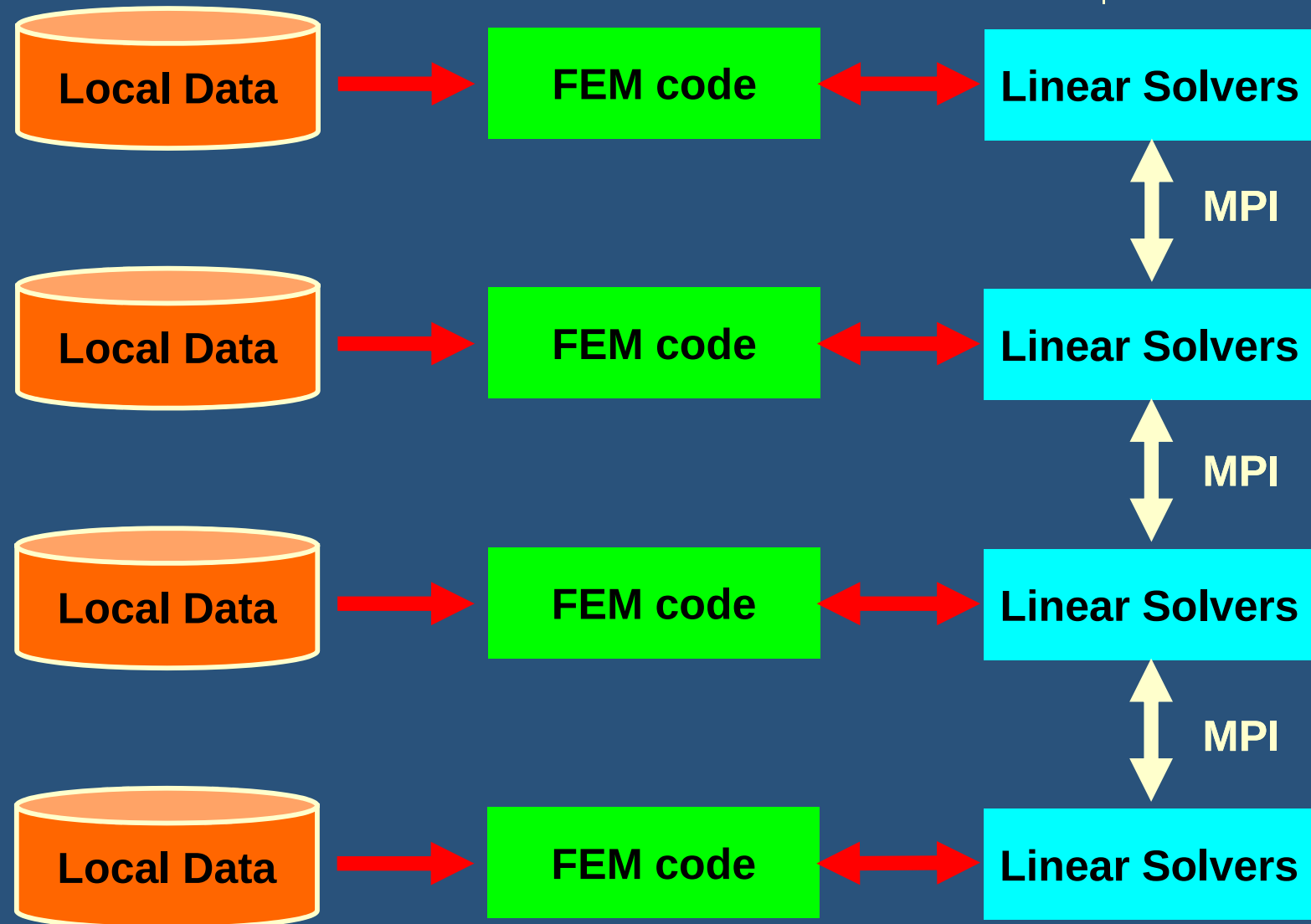
We do not need communication during matrix assemble !!

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- Elements which include Internal Nodes
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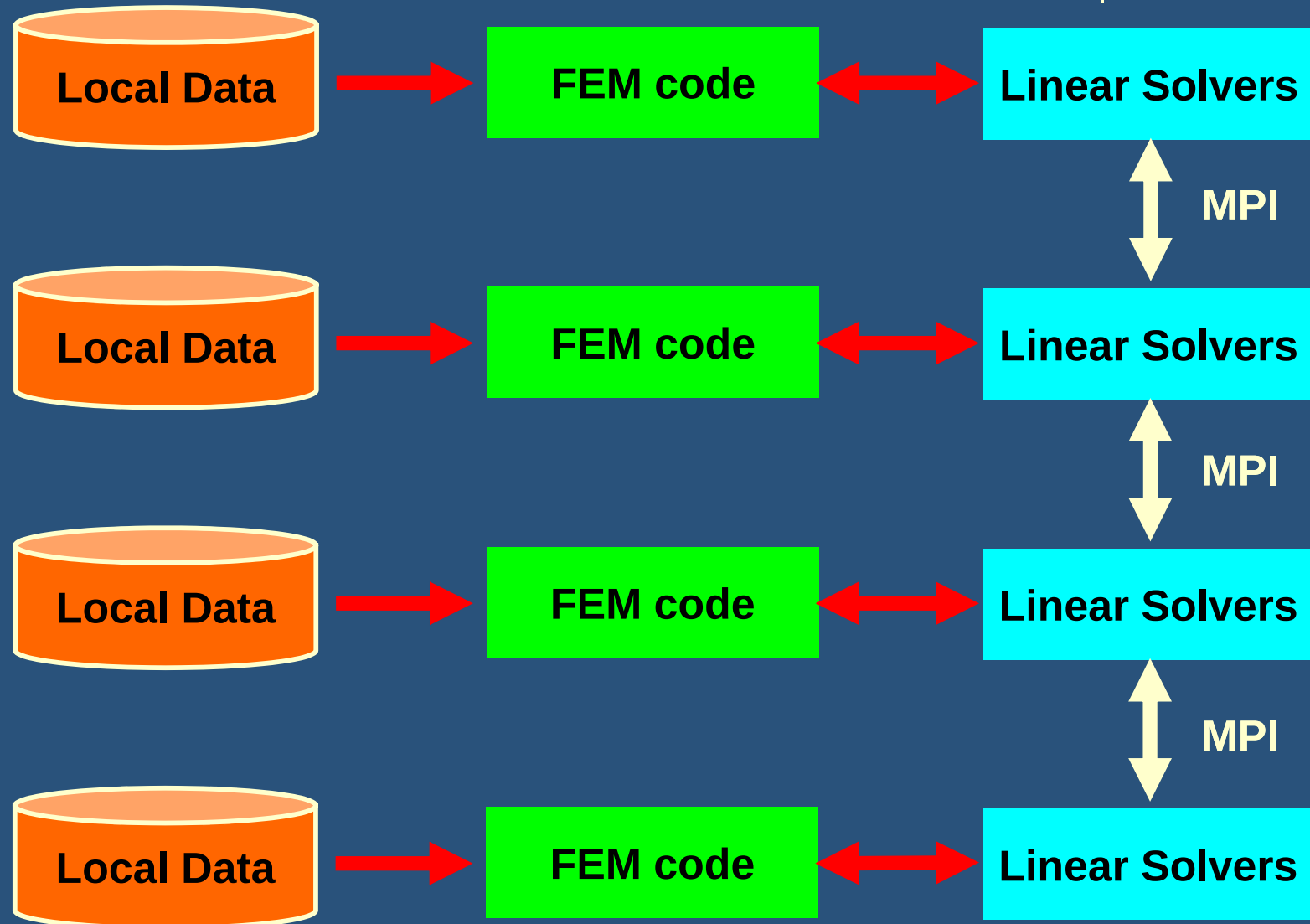
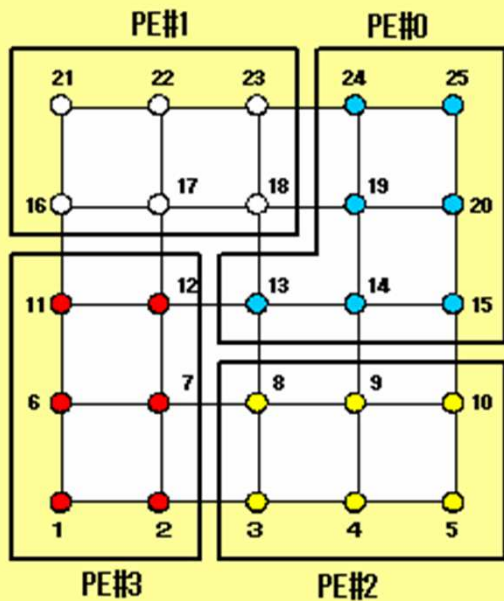
Parallel Computing in FEM

SPMD: Single-Program Multiple-Data



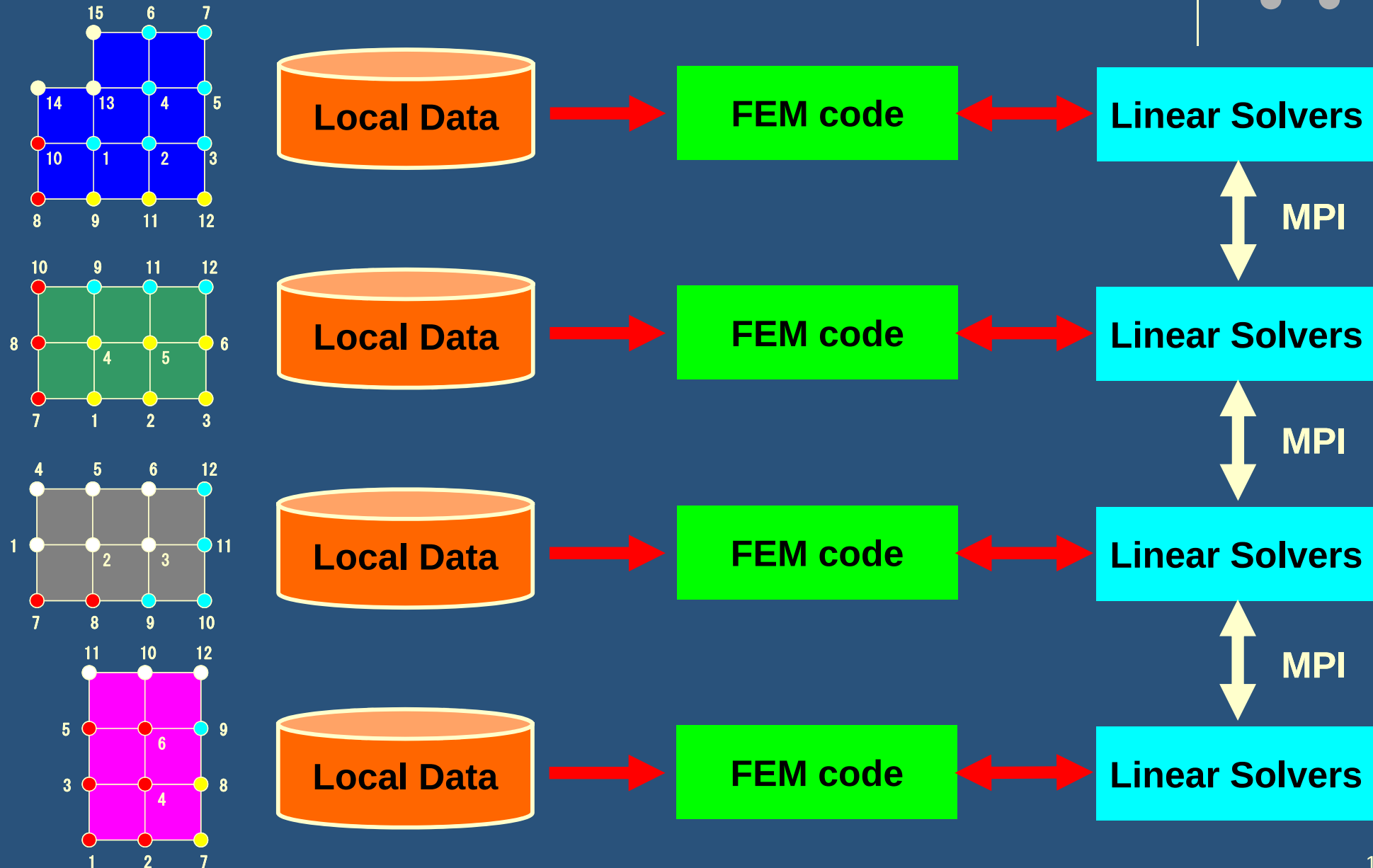
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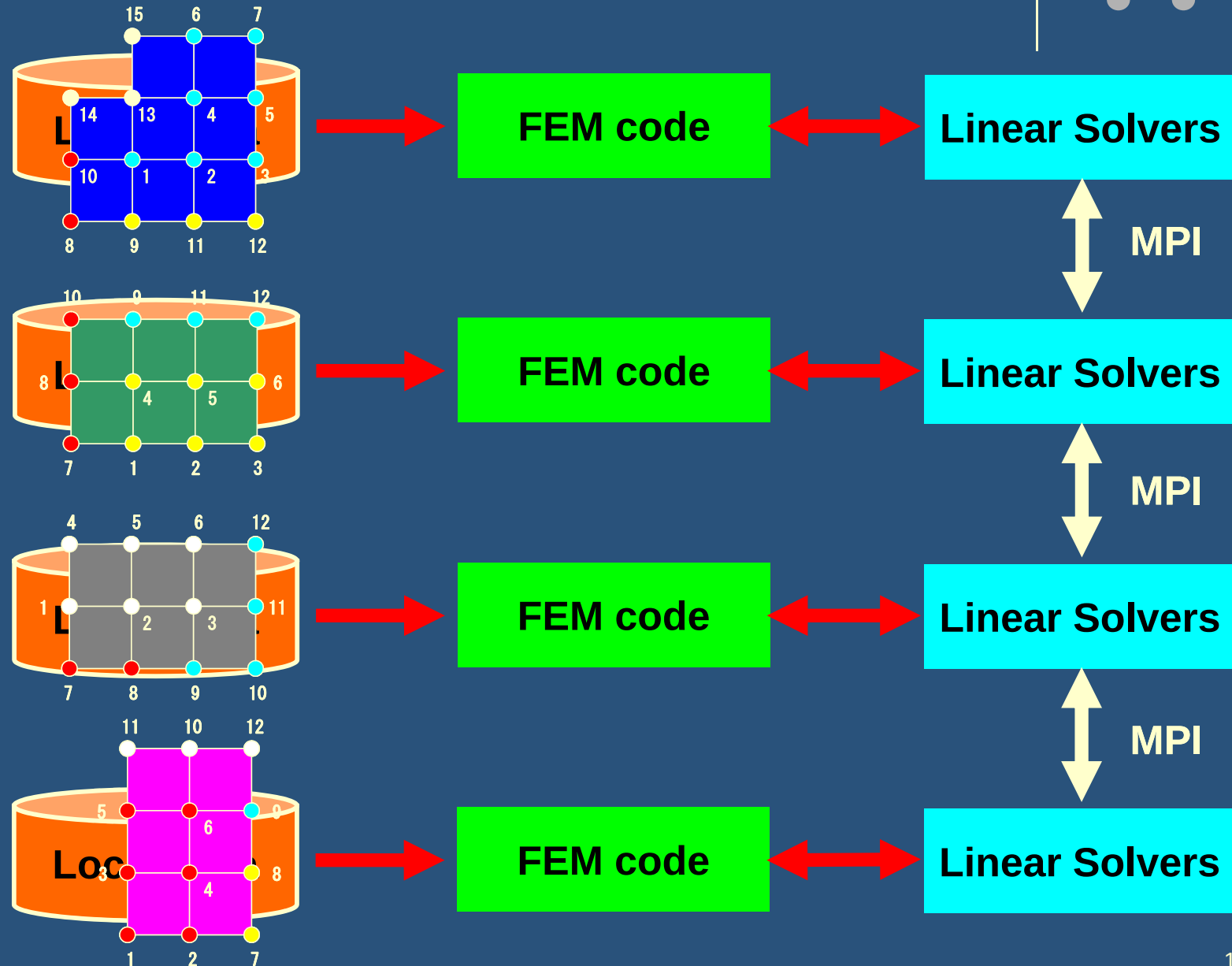
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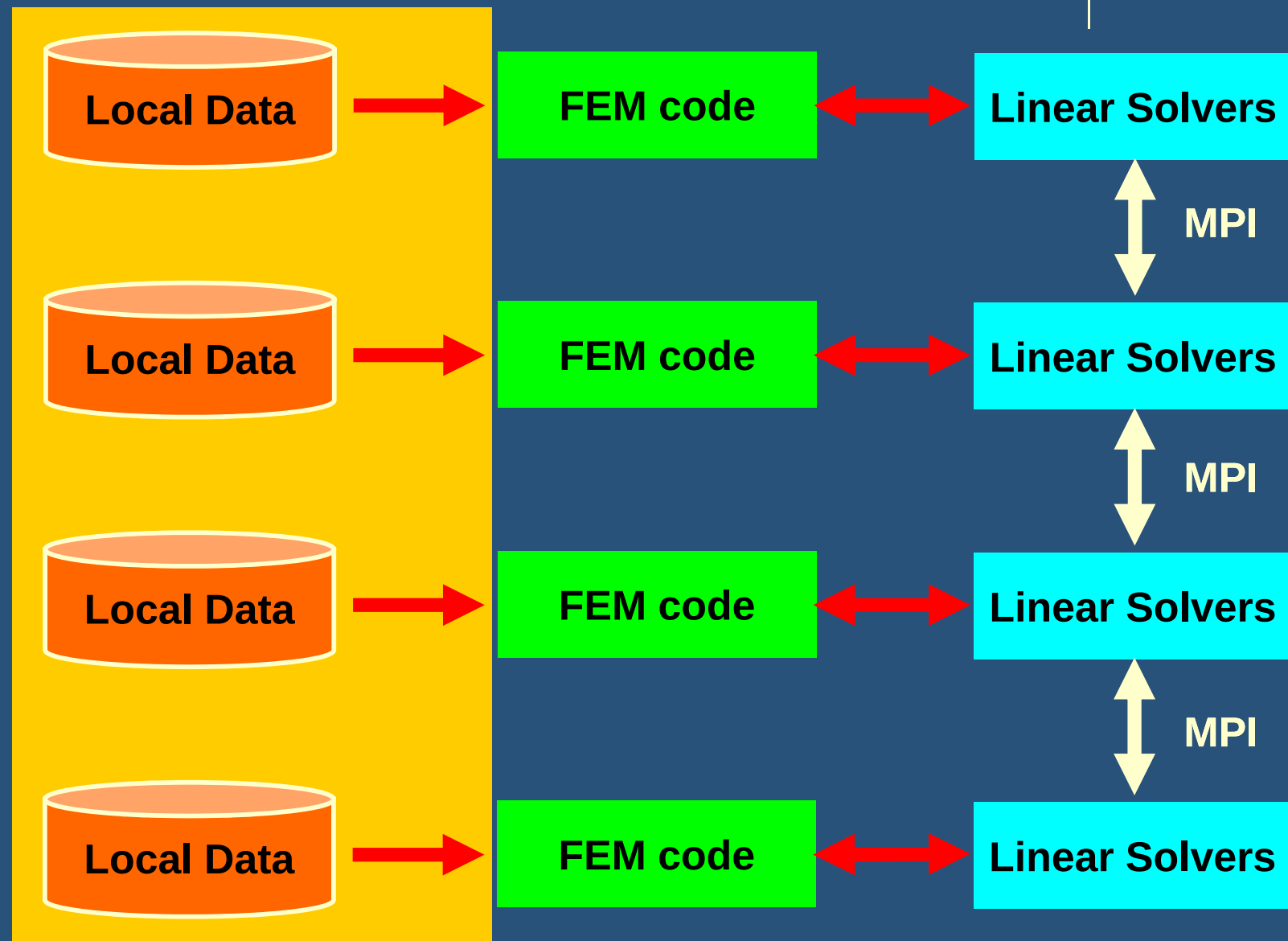
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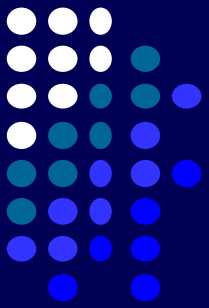


Parallel Computing in FEM

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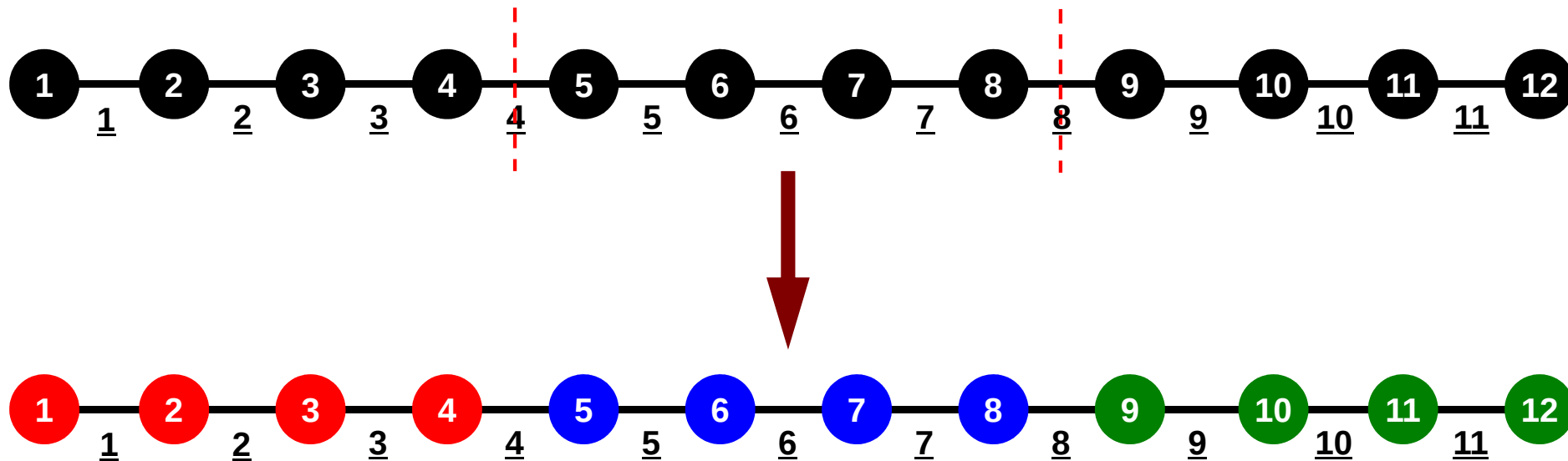


What is Communications ?

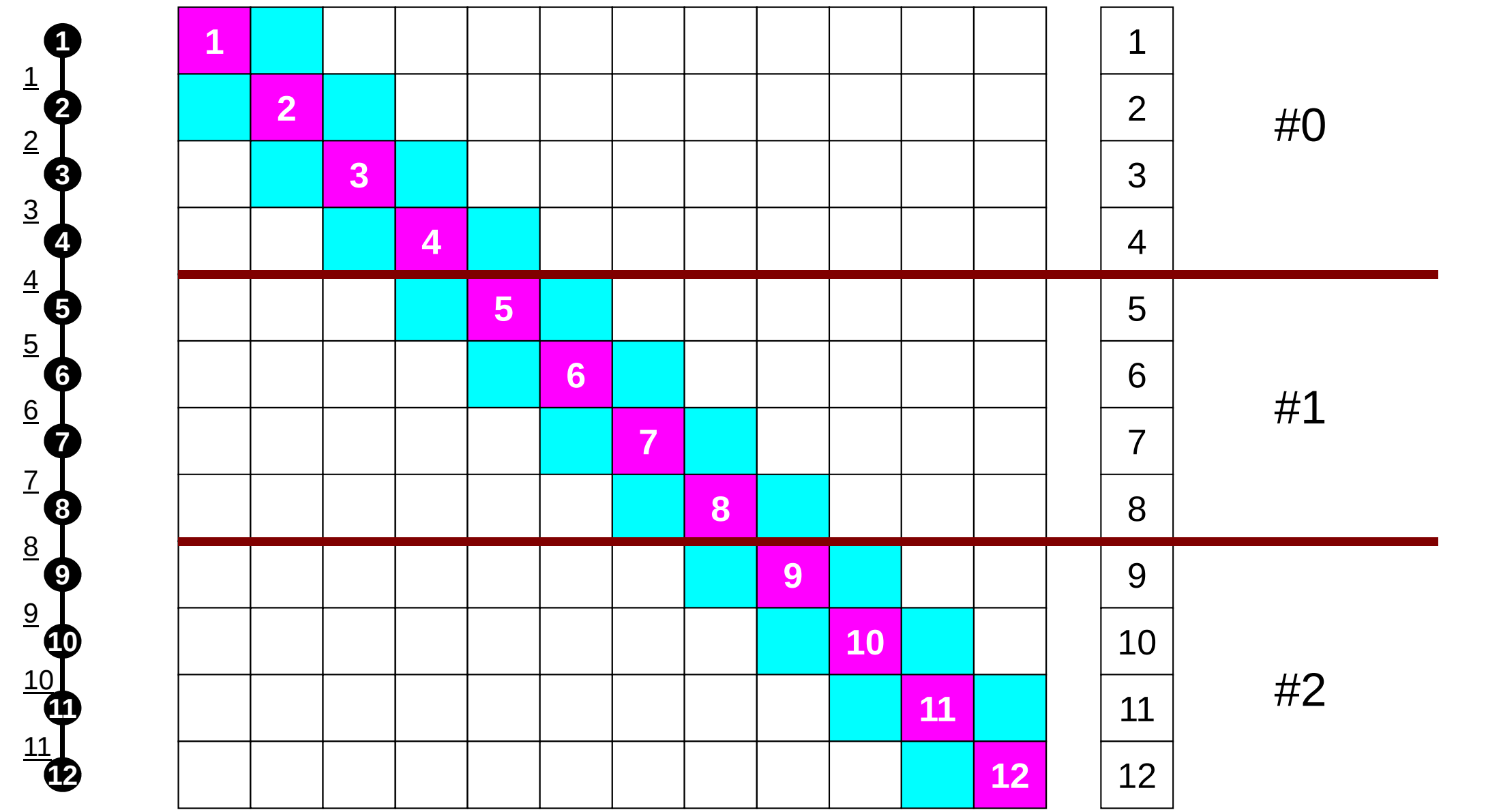


- to get information of “external nodes” from external partitions (local data)
- “Communication tables” contain the information

1D FEM: 12 nodes/11 elem's/3 domains

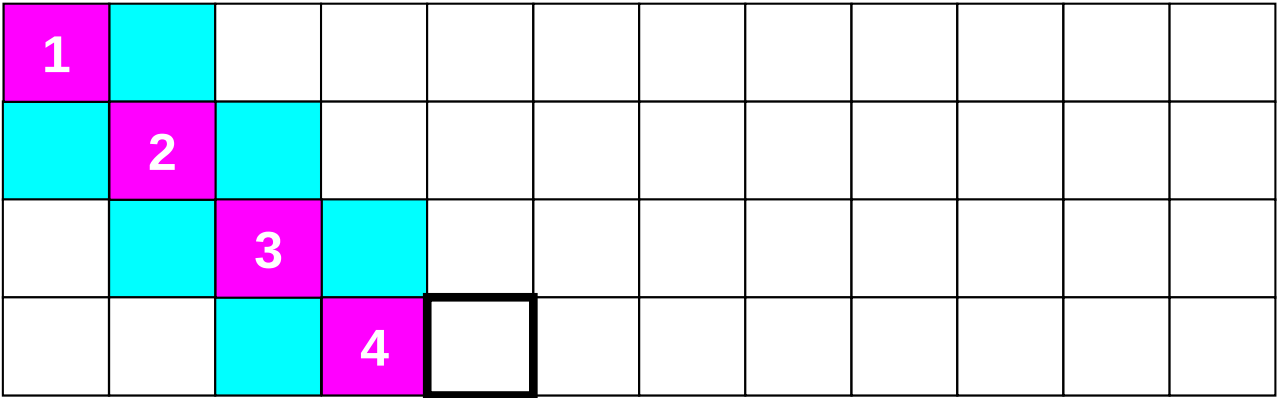


“Internal Nodes” should be balanced



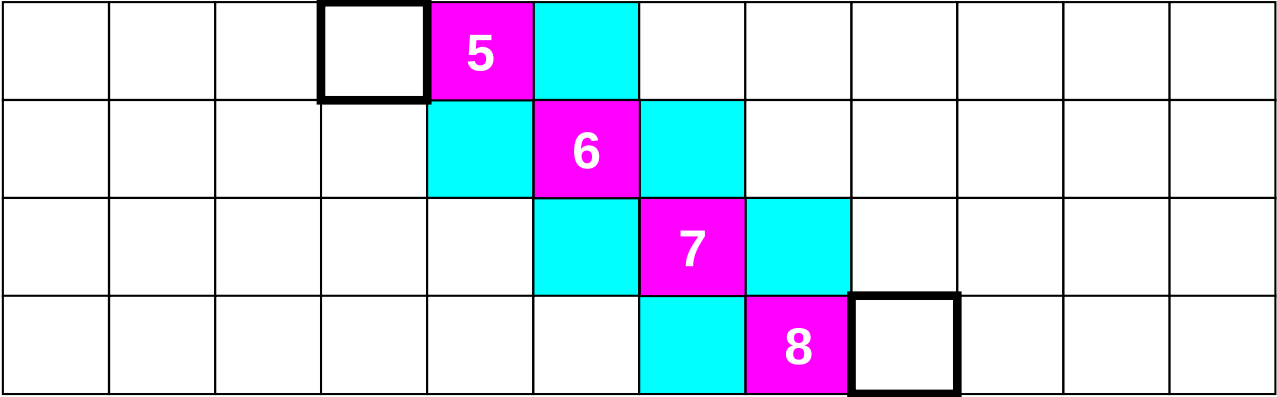
Matrices are incomplete !

- 1
- 2
- 3
- 4



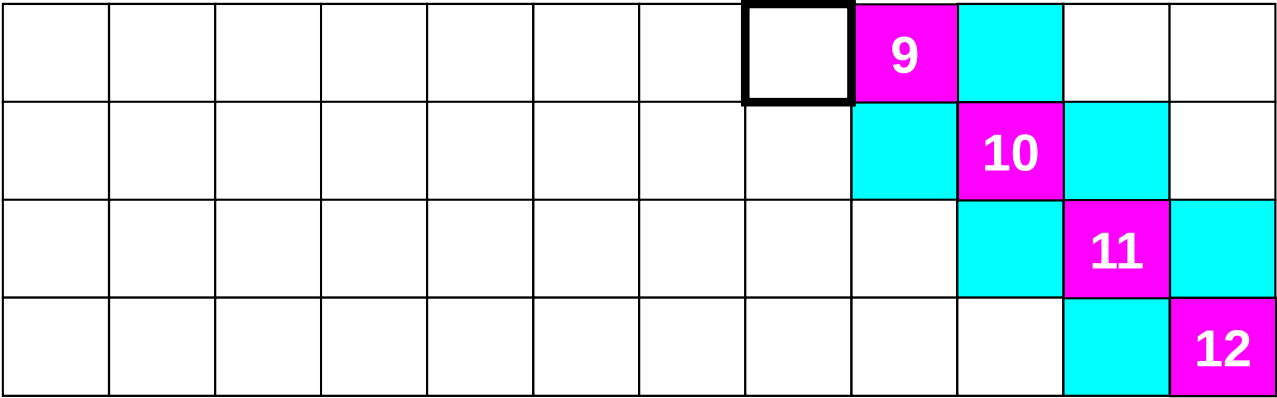
#0

- 5
- 6
- 7
- 8



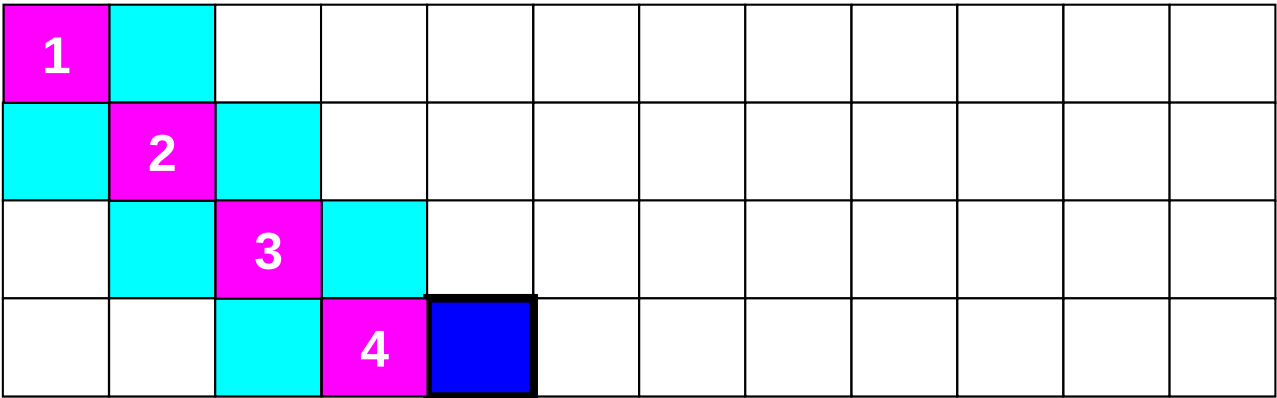
#1

- 9
- 10
- 11
- 12

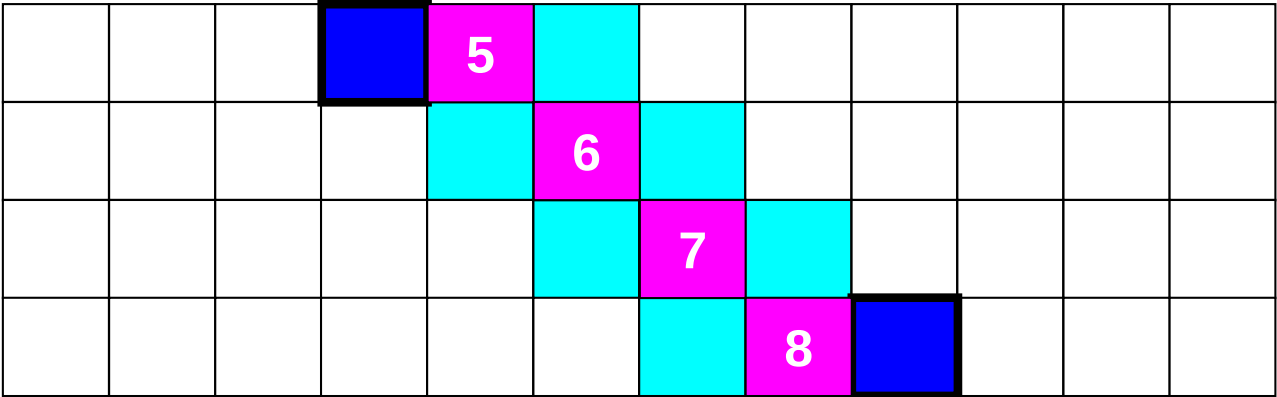


#2

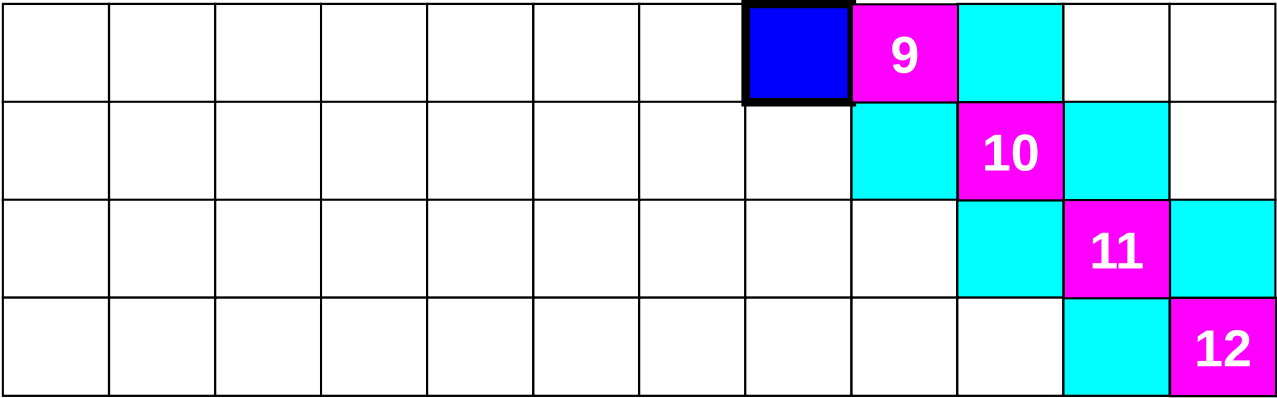
Connected Elements + External Nodes



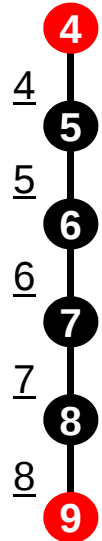
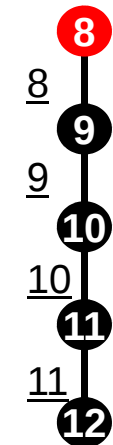
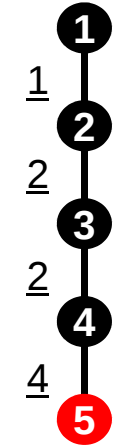
#0



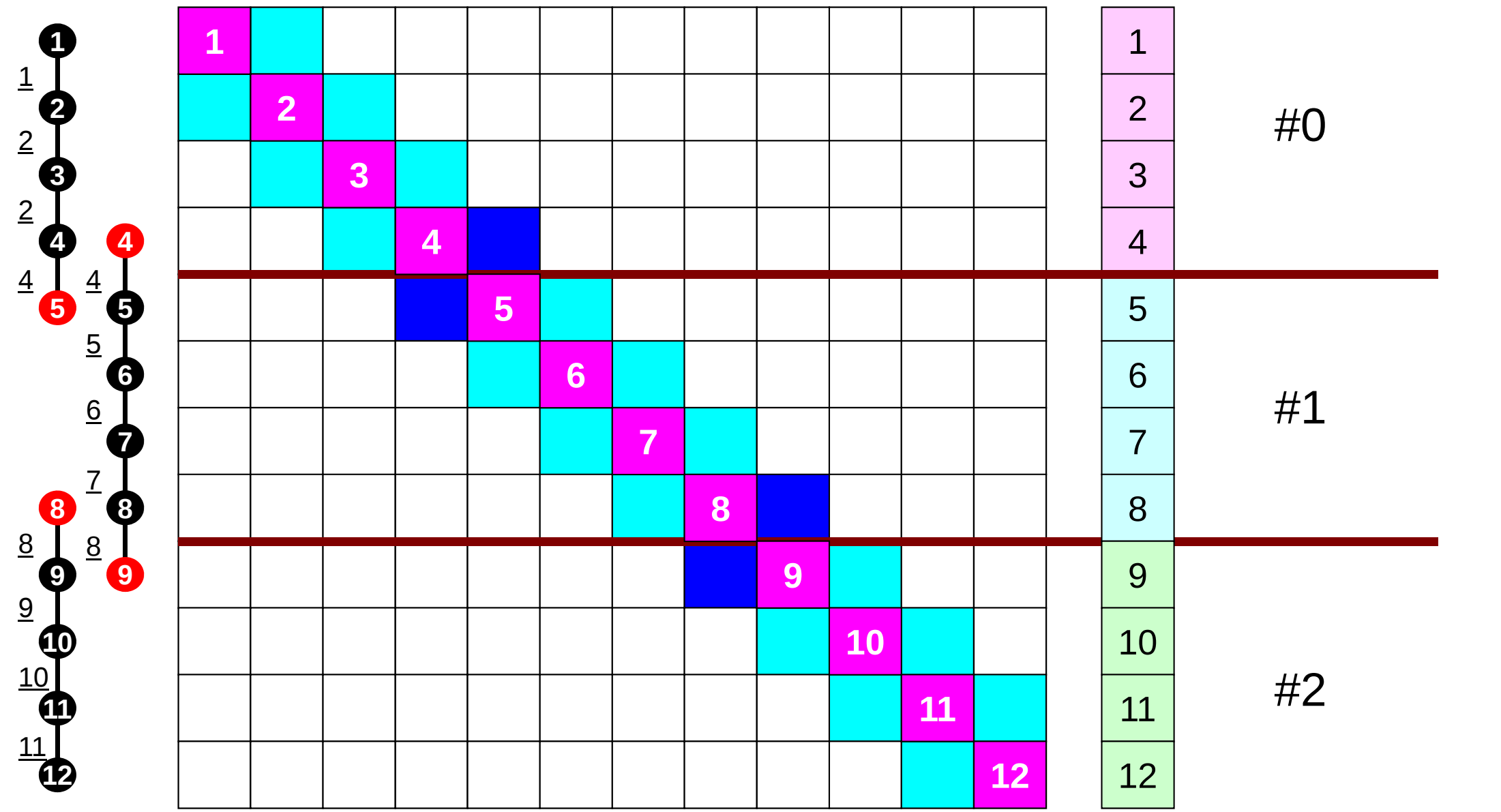
#1



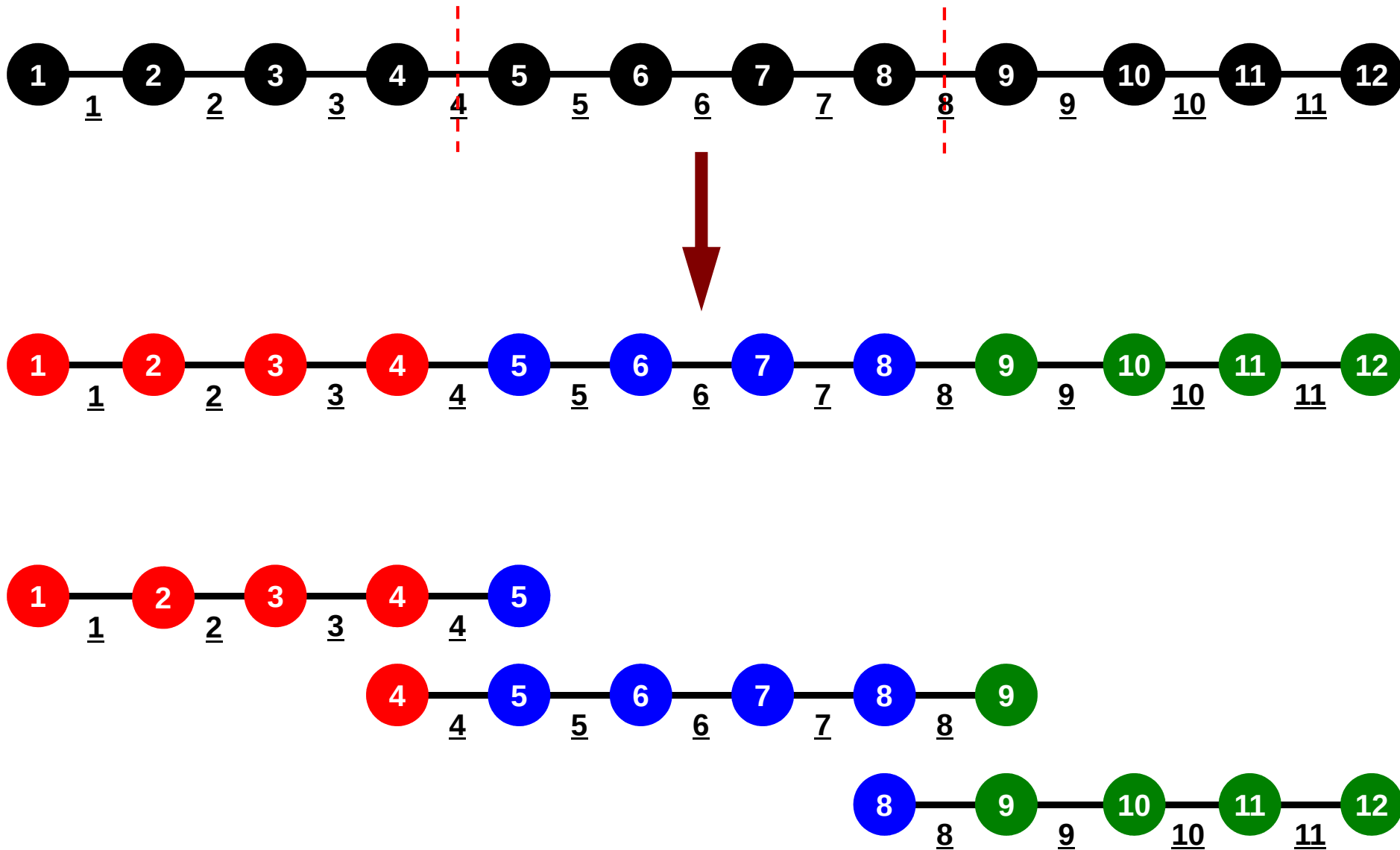
#2



1D FEM: 12 nodes/11 elem's/3 domains

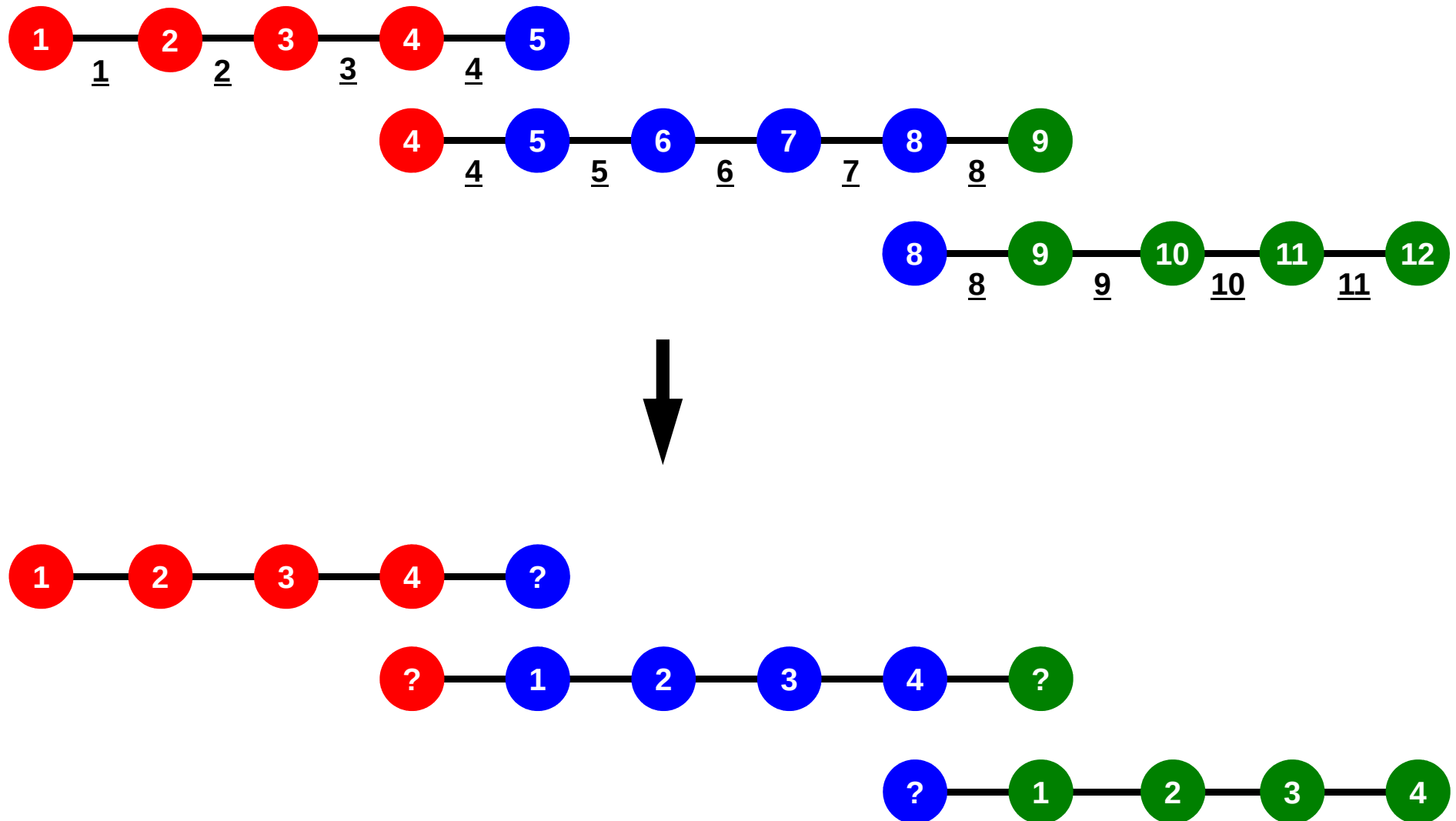


1D FEM: 12 nodes/11 elem's/3 domains



Local Numbering for SPMD

Numbering of internal nodes is 1-N (0-N-1), same operations in serial program can be applied. How about numbering of external nodes ?

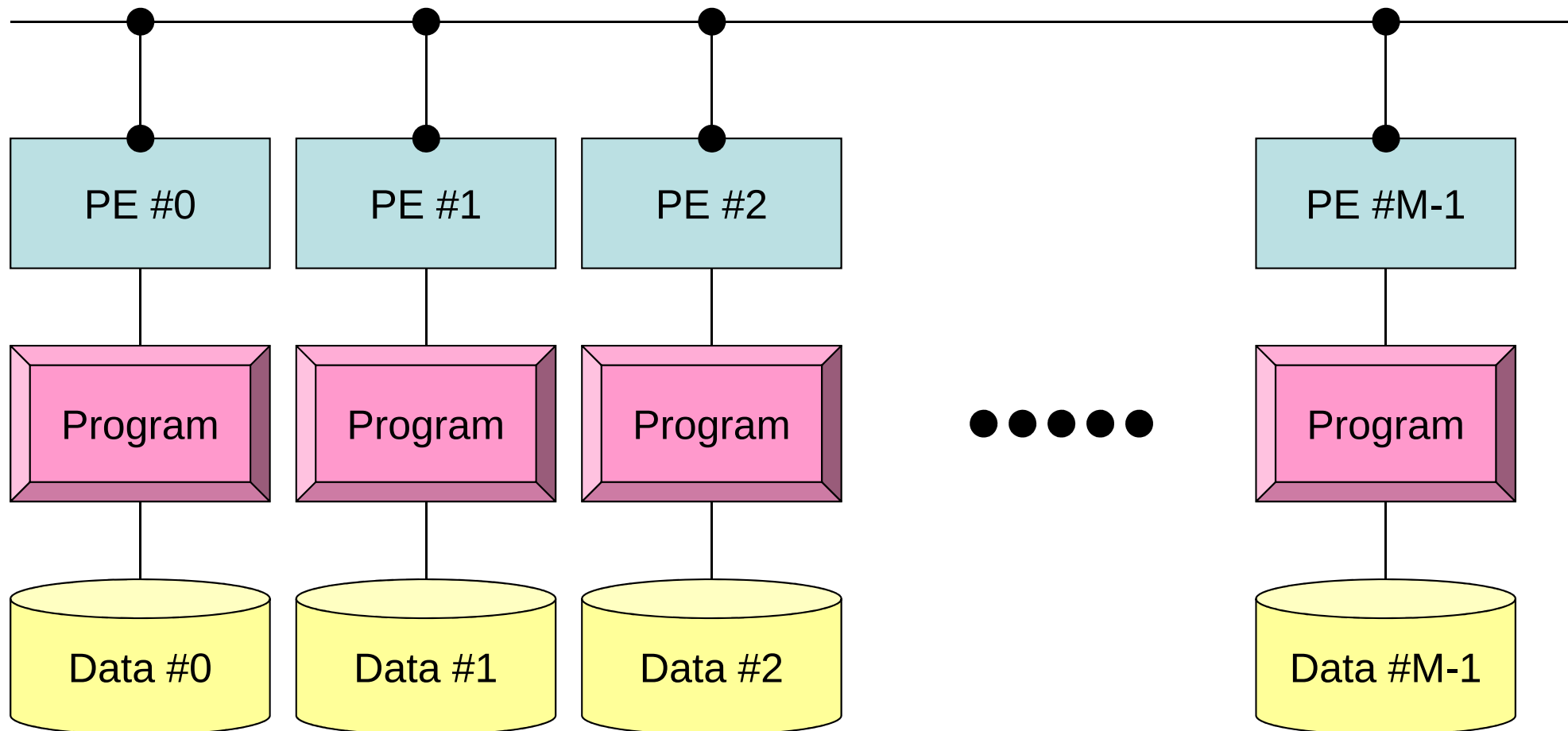


PE: Processing Element
Processor, Domain, Process

SPMD:

Single Program Multiple Data

```
mpirun -np M <Program>
```



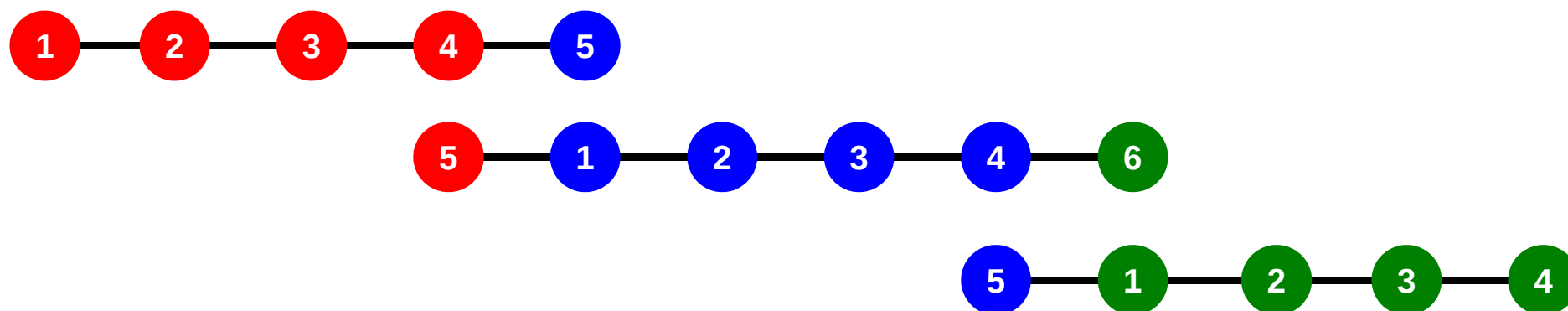
Each process does same operation for different data

Large-scale data is decomposed, and each part is computed by each process

It is ideal that parallel program is not different from serial one except communication.

Local Numbering for SPMD

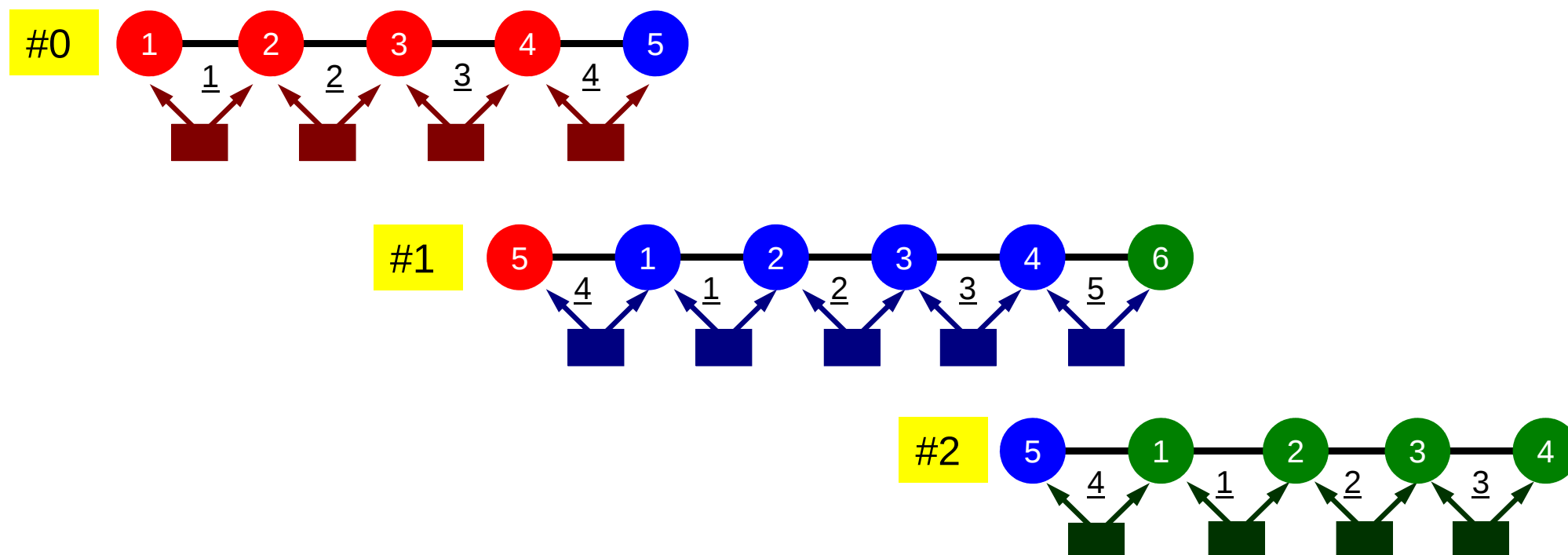
Numbering of external nodes: $N+1$, $N+2$ ($N, N+1$)



1D FEM: 12 nodes/11 elem's/3 domains

Integration on each element, element matrix $\rightarrow$ global matrix

Operations can be done by info. of internal/external nodes and elements which include these nodes



Finite Element Procedures

- Initialization
 - Control Data
 - Node, Connectivity of Elements (N: Node#, NE: Elem#)
 - Initialization of Arrays (Global/Element Matrices)
 - Element-Global Matrix Mapping (Index, Item)
- Generation of Matrix
 - Element-by-Element Operations (do icel= 1, NE)
 - Element matrices
 - Accumulation to global matrix
 - Boundary Conditions
- Linear Solver
 - Conjugate Gradient Method

Preconditioned CG Solver

```

Compute  $\mathbf{r}^{(0)} = \mathbf{b} - [\mathbf{A}]\mathbf{x}^{(0)}$ 
for i= 1, 2, ...
  solve  $[\mathbf{M}]\mathbf{z}^{(i-1)} = \mathbf{r}^{(i-1)}$ 
   $\rho_{i-1} = \mathbf{r}^{(i-1)} \cdot \mathbf{z}^{(i-1)}$ 
  if i=1
     $\mathbf{p}^{(1)} = \mathbf{z}^{(0)}$ 
  else
     $\beta_{i-1} = \rho_{i-1} / \rho_{i-2}$ 
     $\mathbf{p}^{(i)} = \mathbf{z}^{(i-1)} + \beta_{i-1} \mathbf{p}^{(i-1)}$ 
  endif
   $\mathbf{q}^{(i)} = [\mathbf{A}]\mathbf{p}^{(i)}$ 
   $\alpha_i = \rho_{i-1} / \mathbf{p}^{(i)} \cdot \mathbf{q}^{(i)}$ 
   $\mathbf{x}^{(i)} = \mathbf{x}^{(i-1)} + \alpha_i \mathbf{p}^{(i)}$ 
   $\mathbf{r}^{(i)} = \mathbf{r}^{(i-1)} - \alpha_i \mathbf{q}^{(i)}$ 
  check convergence  $|\mathbf{r}|$ 
end

```

- Preconditioning
 - Diagonal Scaling/Point Jacobi
- Parallel operations are required in
 - Dot Products
 - Mat-Vec. Multiplication
 - SpMV: Sparse Mat-Vec. Mult.

$$[\mathbf{M}] = \begin{bmatrix} D_1 & 0 & \dots & 0 & 0 \\ 0 & D_2 & & 0 & 0 \\ \dots & & \dots & & \dots \\ 0 & 0 & & D_{N-1} & 0 \\ 0 & 0 & \dots & 0 & D_N \end{bmatrix}$$

Preconditioning, DAXPY

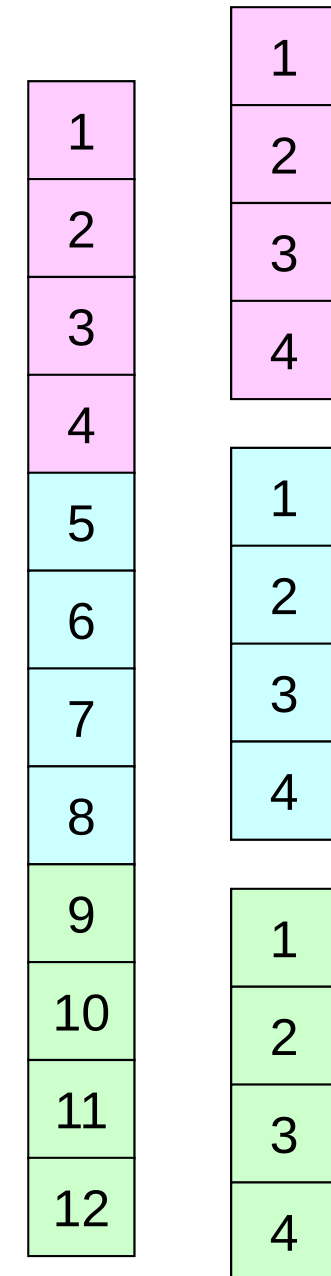
Local Operations by Only Internal Points: Parallel Processing is possible

```
!C
!C-- {z}= [Minv]{r}

do i= 1, N
  W(i,Z)= W(i,DD) * W(i,R)
enddo
```

```
!C
!C-- {x}= {x} + ALPHA*{p}      DAXPY: double a{x} plus {y}
!C  {r}= {r} - ALPHA*{q}

do i= 1, N
  PHI(i)= PHI(i) + ALPHA * W(i,P)
  W(i,R)= W(i,R) - ALPHA * W(i,Q)
enddo
```

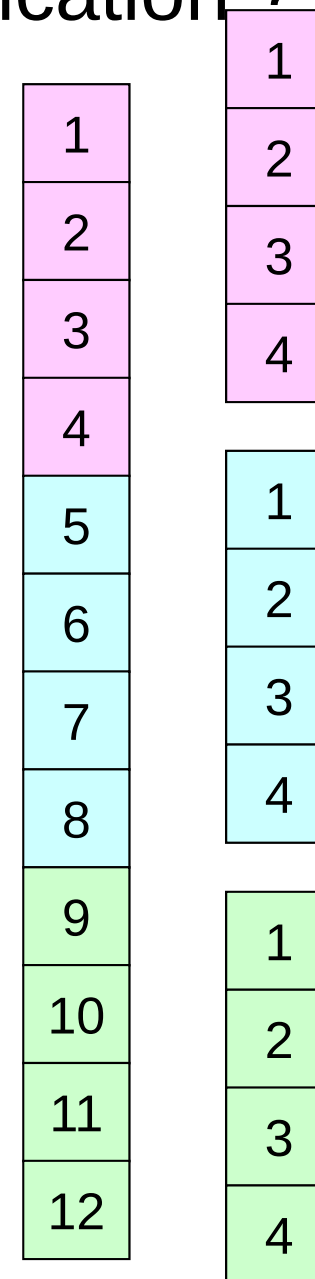


Dot Products

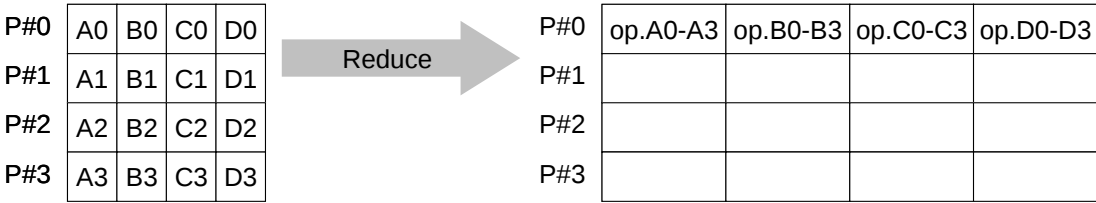
Global Summation needed: Communication ?

```
!C
!C-- ALPHA= RHO / {p} {q}

C1= 0. d0
do i= 1, N
  C1= C1 + W(i, P)*W(i, Q)
enddo
ALPHA= RHO / C1
```



MPI_REDUCE



- Reduces values on all processes to a single value
 - Summation, Product, Max, Min etc.
- **call MPI_REDUCE**
(sendbuf, recvbuf, count, datatype, op, root, comm, ierr)
 - **sendbuf** choice I starting address of send buffer
 - **recvbuf** choice O starting address receive buffer
type is defined by "**datatype**"
 - **count** I I number of elements in send/receive buffer
 - **datatype** I I data type of elements of send/recive buffer
 - FORTTRAN MPI_INTEGER, MPI_REAL, MPI_DOUBLE_PRECISION, MPI_CHARACTER etc.
 - C MPI_INT, MPI_FLOAT, MPI_DOUBLE, MPI_CHAR etc
 - **op** I I reduce operation
 - MPI_MAX, MPI_MIN, MPI_SUM, MPI_PROD, MPI_LAND, MPI_BAND etc
 - Users can define operations by **MPI_OP_CREATE**
 - **root** I I rank of root process
 - **comm** I I communicator
 - **ierr** I O completion code

Diagram illustrating a Broadcast operation:

Initial State (Left):

P#0	A0	B0	C0	D0
P#1				
P#2				
P#3				

Broadcast Operation (Arrow):

Final State (Right):

P#0	A0	B0	C0	D0
P#1	A0	B0	C0	D0
P#2	A0	B0	C0	D0
P#3	A0	B0	C0	D0

- # Fortran

“op” of MPI_Reduce/Allreduce

```
call MPI_REDUCE
```

```
(sendbuf, recvbuf, count, datatype, op, root, comm, ierr)
```

- `MPI_MAX, MPI_MIN` Max, Min
- `MPI_SUM, MPI_PROD` Summation, Product
- `MPI_LAND` Logical AND

Preconditioned CG Solver

```

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```

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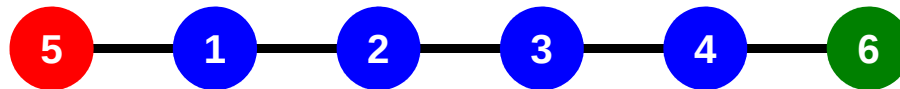
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Matrix-Vector Products

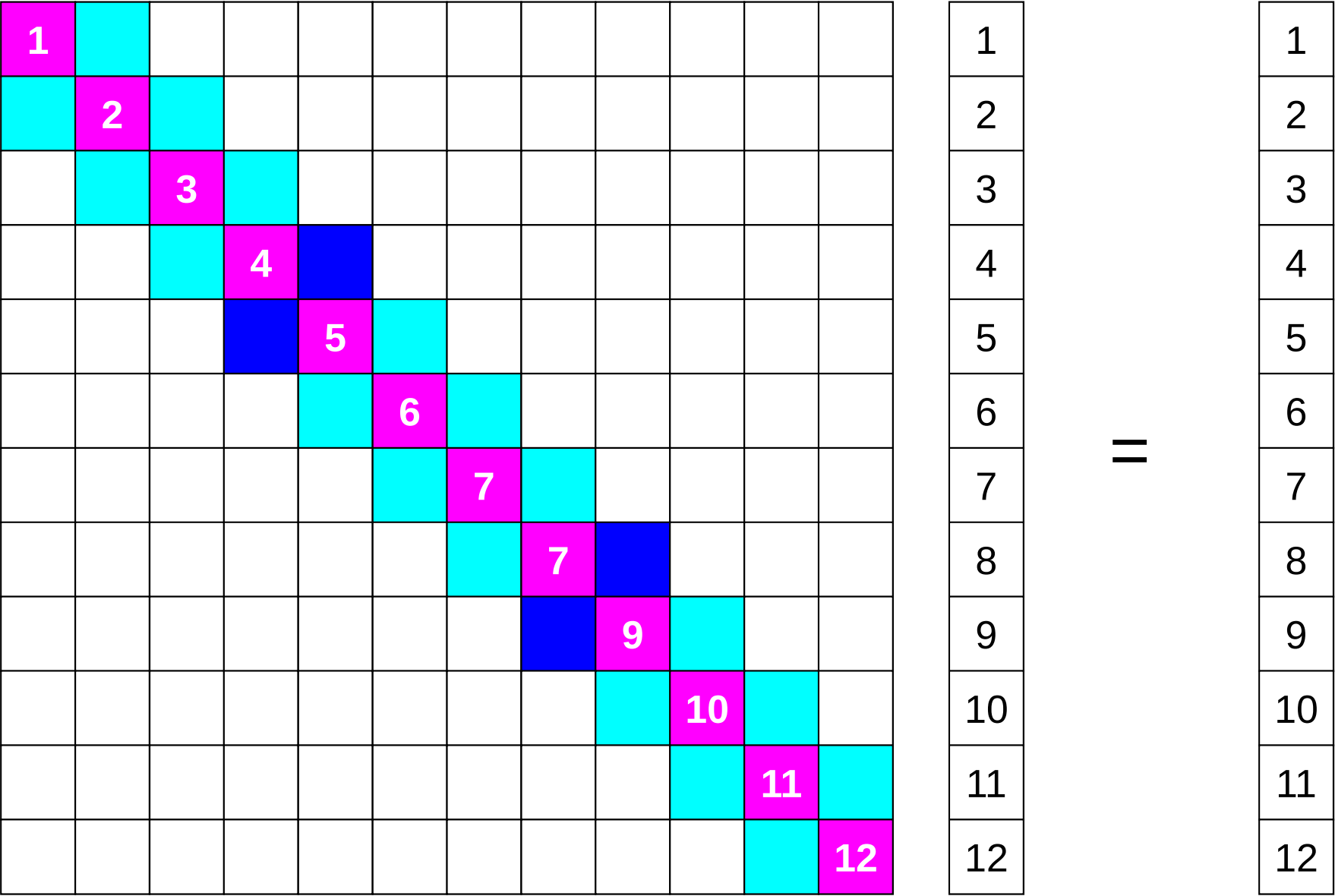
Values at External Points: P-to-P Communication

```
!C
!C-- {q} = [A] {p}

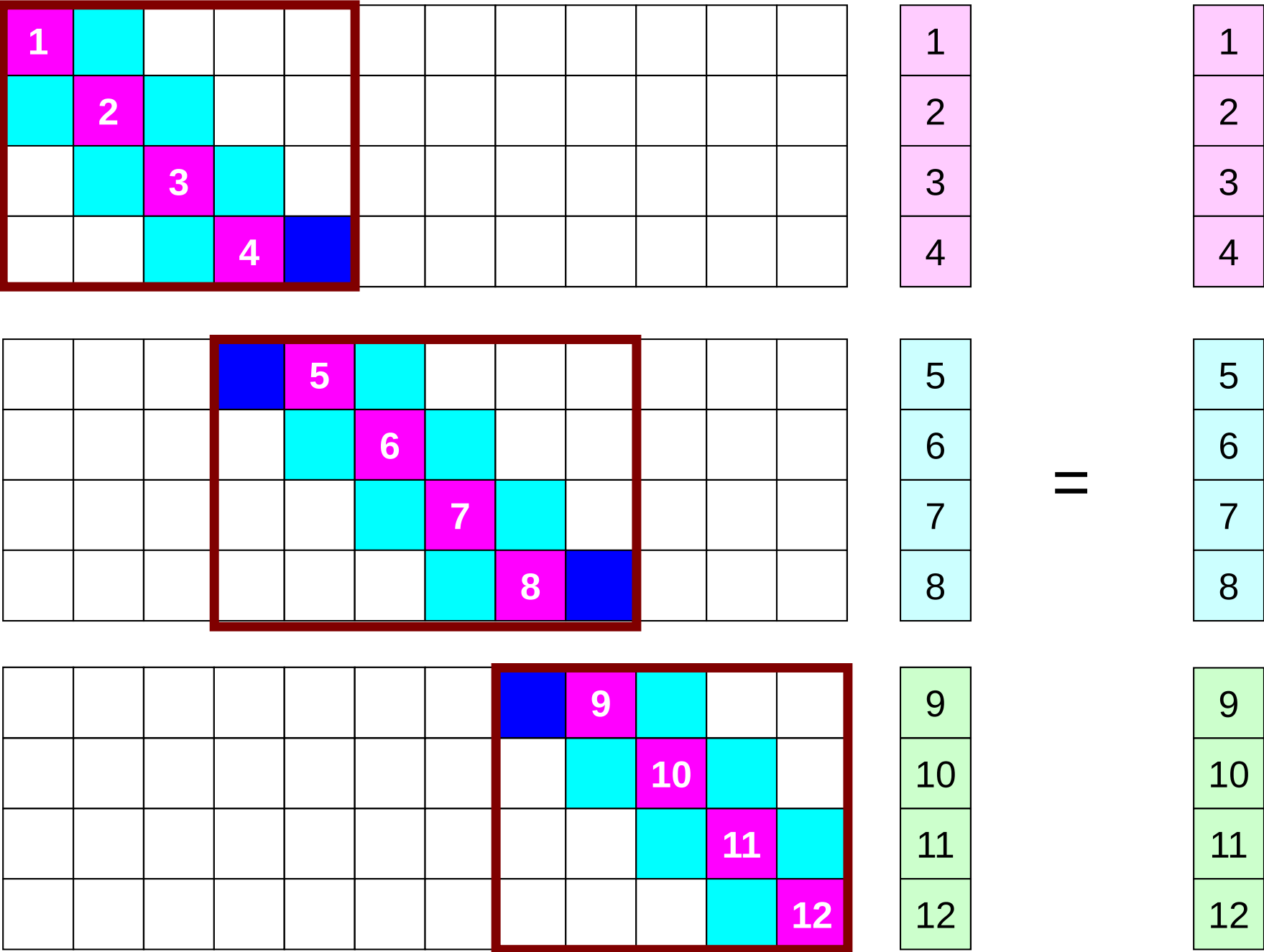
do i= 1, N
  W(i, Q) = DIAG(i)*W(i, P)
  do j= INDEX(i-1)+1, INDEX(i)
    W(i, Q) = W(i, Q) + AMAT(j)*W(ITEM(j), P)
  enddo
enddo
```



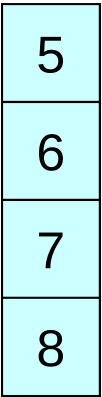
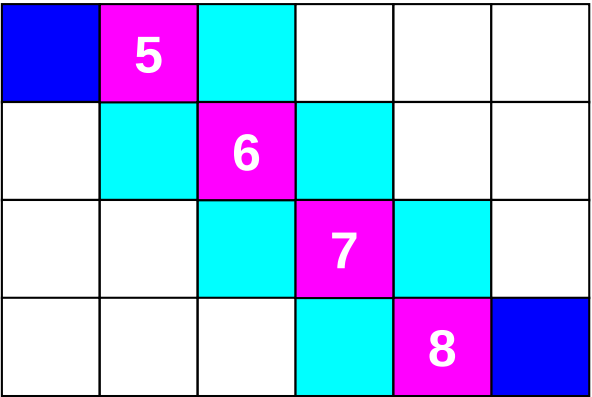
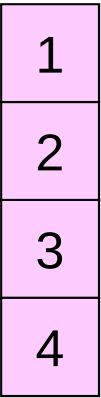
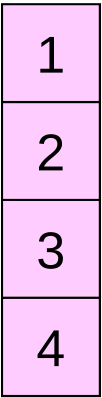
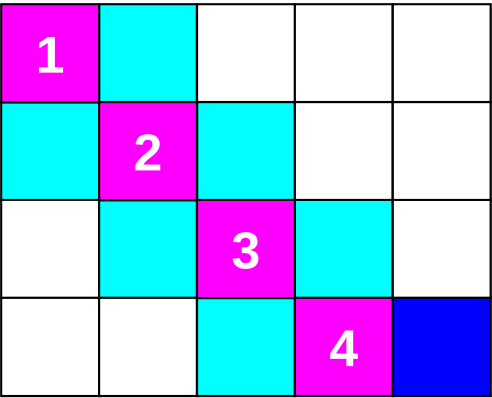
Mat-Vec Products: Local Op. Possible



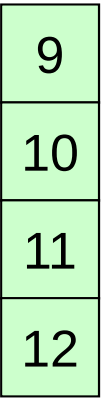
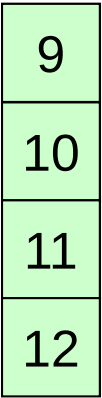
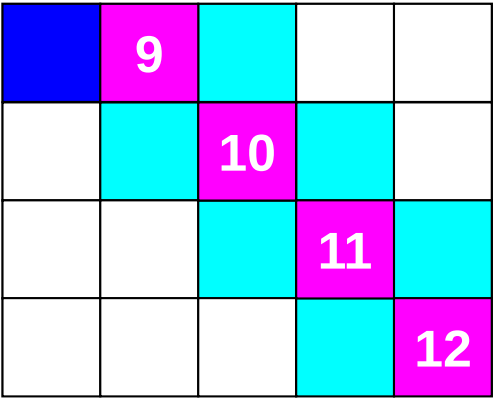
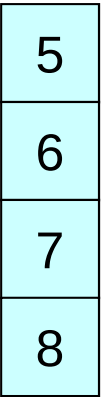
Mat-Vec Products: Local Op. Possible



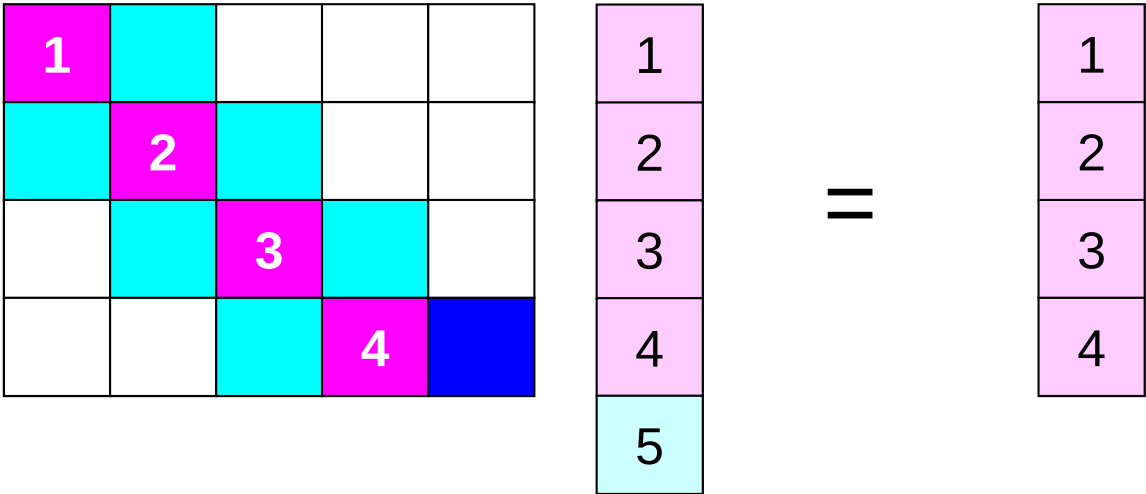
Mat-Vec Products: Local Op. Possible



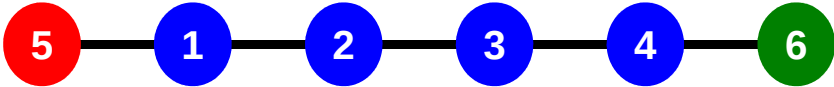
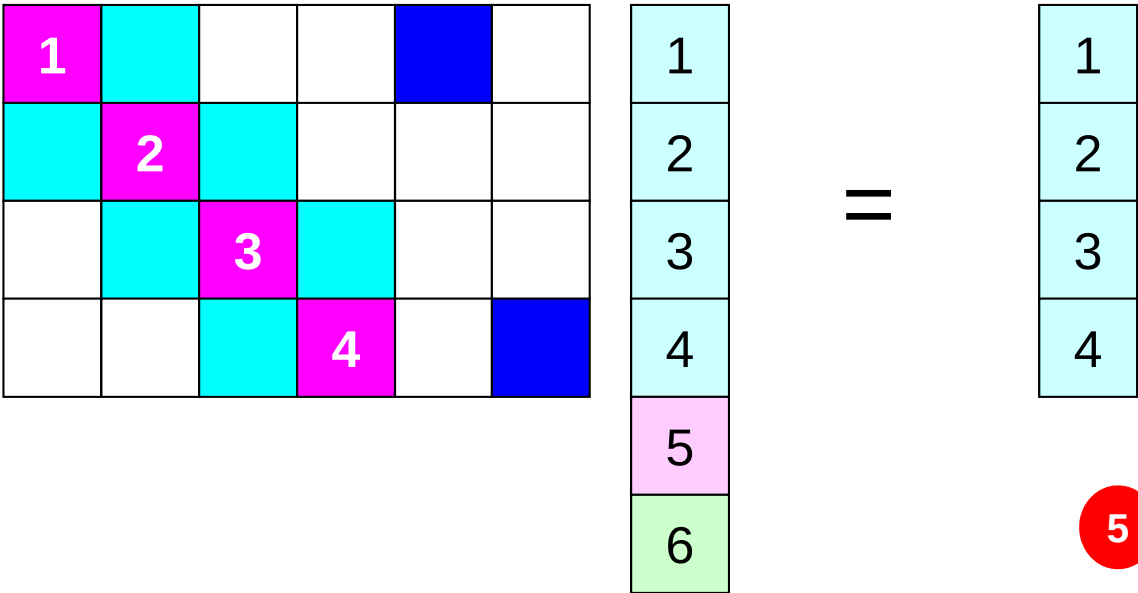
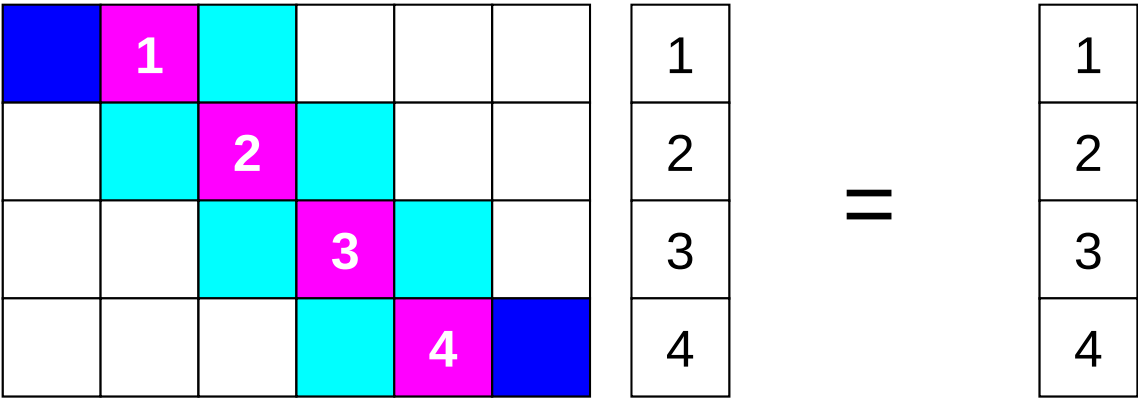
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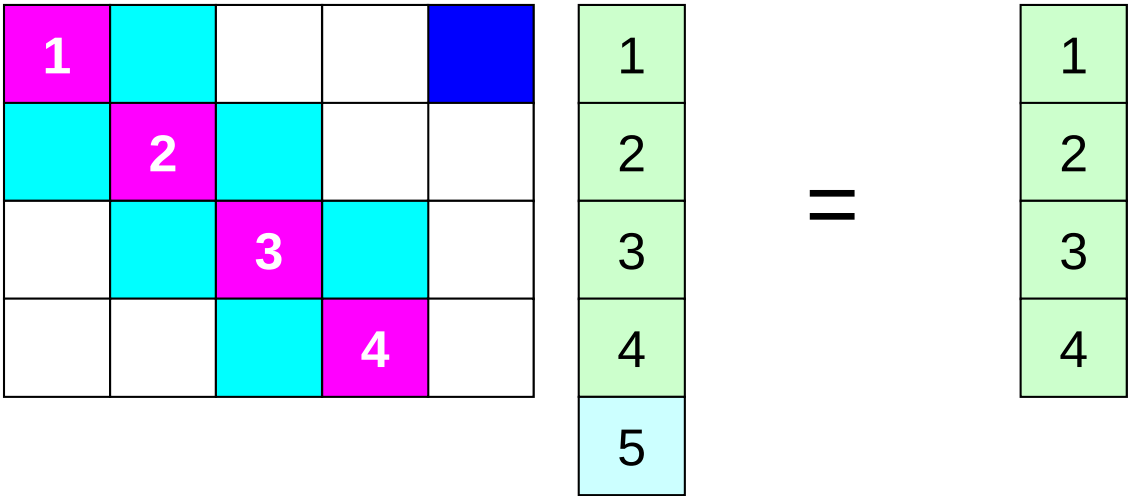
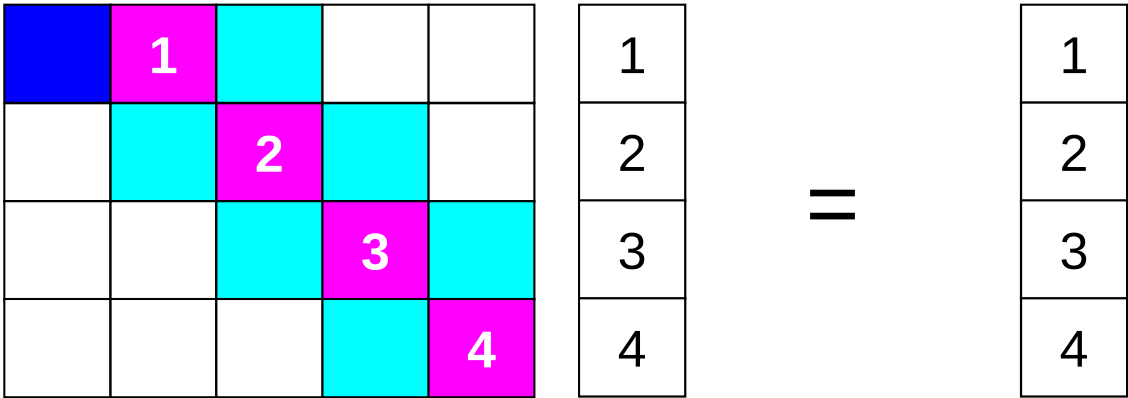
Mat-Vec Products: Local Op. #0



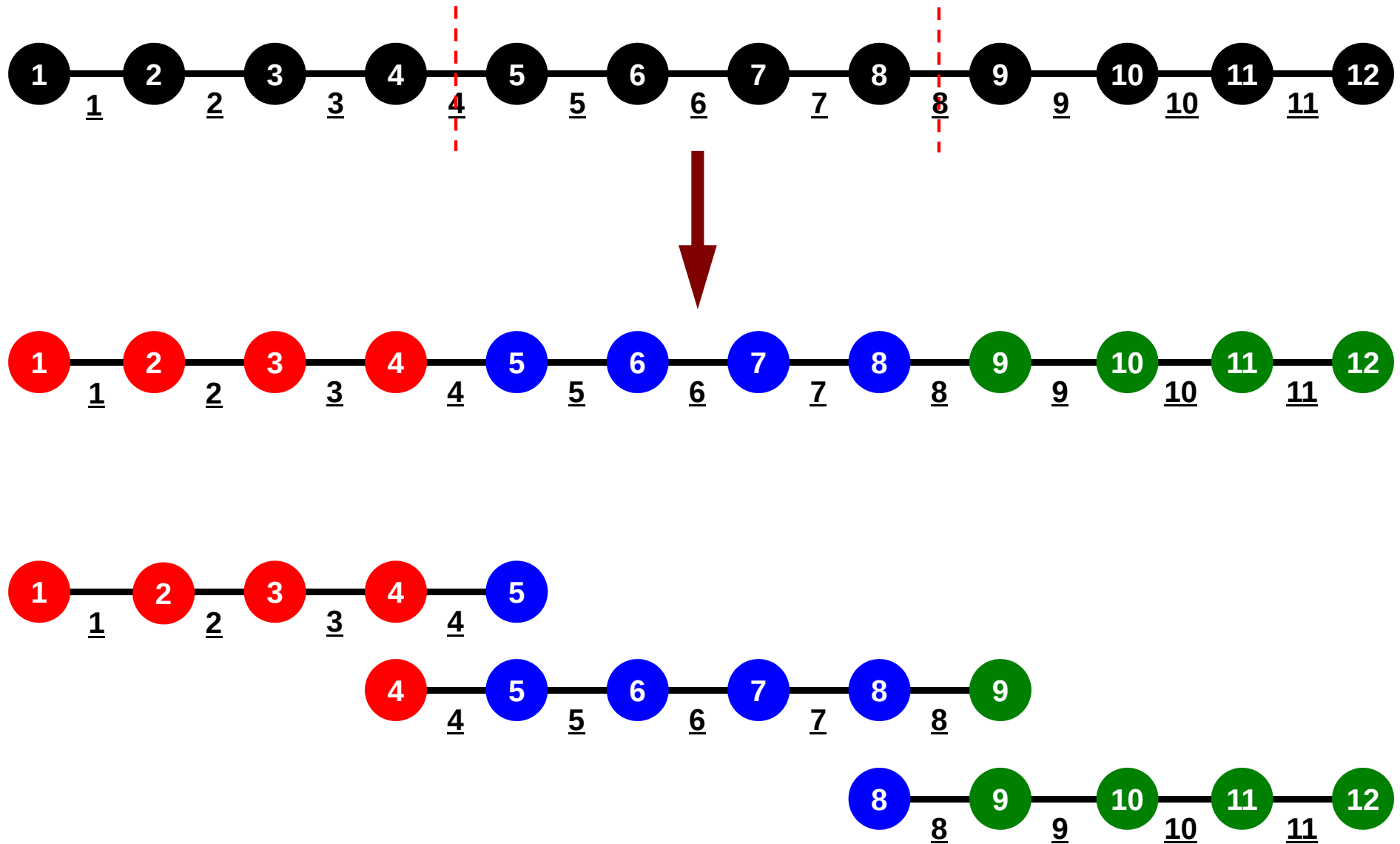
Mat-Vec Products: Local Op. #1



Mat-Vec Products: Local Op. #2

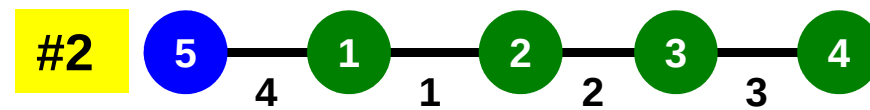
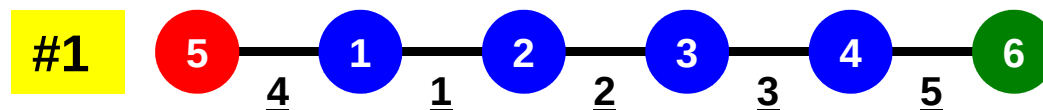
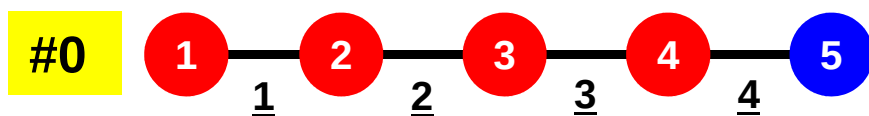


1D FEM: 12 nodes/11 elem's/3 domains



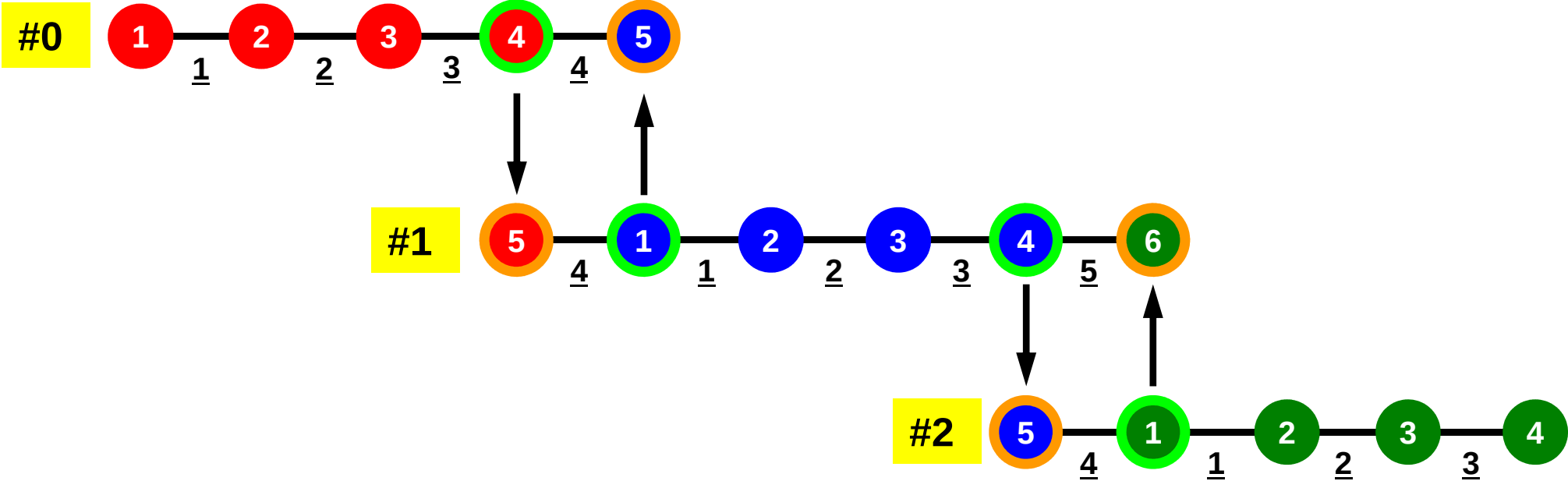
1D FEM: 12 nodes/11 elem's/3 domains

Local ID: Starting from 1 for node and elem at each domain



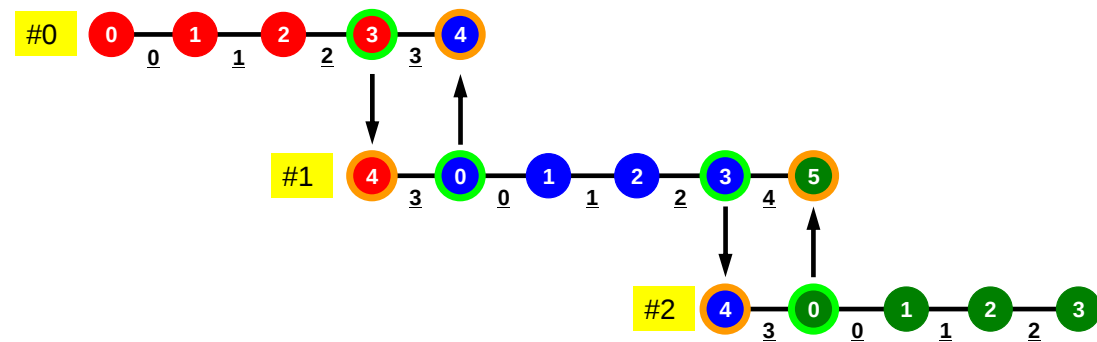
1D FEM: 12 nodes/11 elem's/3 domains

Internal/External Nodes



What is Peer-to-Peer Communication ?

- Collective Communication
 - MPI_Reduce, MPI_Scatter/Gather etc.
 - Communications with all processes in the communicator
 - Application Area
 - BEM, Spectral Method, MD: global interactions are considered
 - Dot products, MAX/MIN: Global Summation & Comparison
- Peer-toPeer/Point-to-Point
 - MPI_Send, MPI_Receive
 - Communication with limited processes
 - Neighbors
 - Application Area
 - FEM, FDM: Localized Method



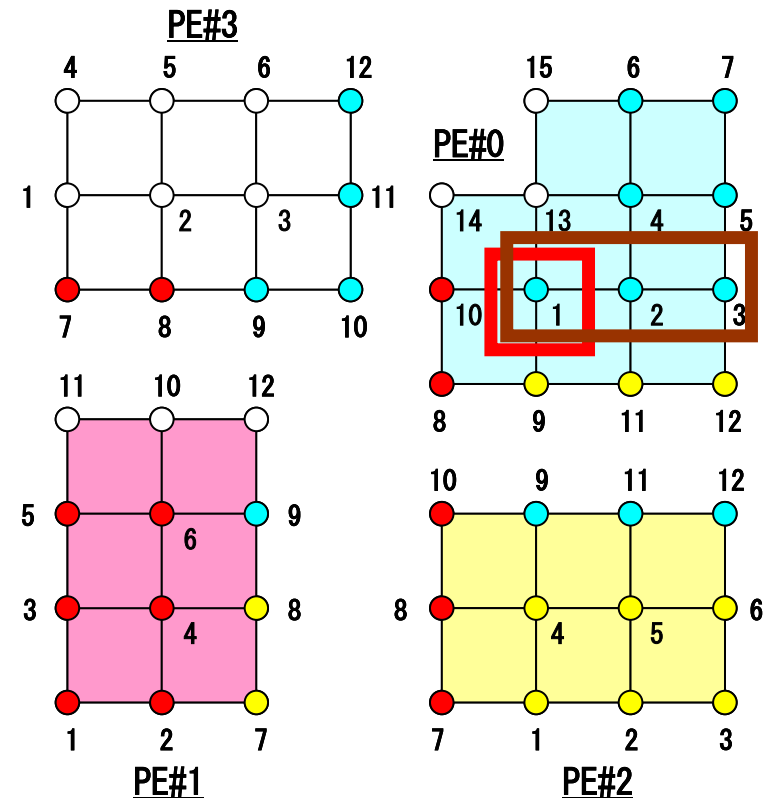
SEND: sending from boundary nodes

Send continuous data to send buffer of neighbors

- MPI_Isend**

(sendbuf, count, datatype, dest, tag, comm, request)

- **sendbuf** choice I starting address of sending buffer
- **count** I I number of elements sent to each process
- **datatype** I I data type of elements of sending buffer
- **dest** I I rank of destination



- Begins a non-blocking send
 - Send the contents of sending buffer (starting from **sendbuf**, number of messages: **count**) to **dest** with **tag** .
 - Contents of sending buffer cannot be modified before calling corresponding **MPI_Waitall**.

- | | | | |
|-------------------------|---|---|---|
| - <u>comm</u> | I | I | communicator |
| - <u>request</u> | I | O | communication request array used in MPI_Waitall |
| - ierr | I | O | completion code |

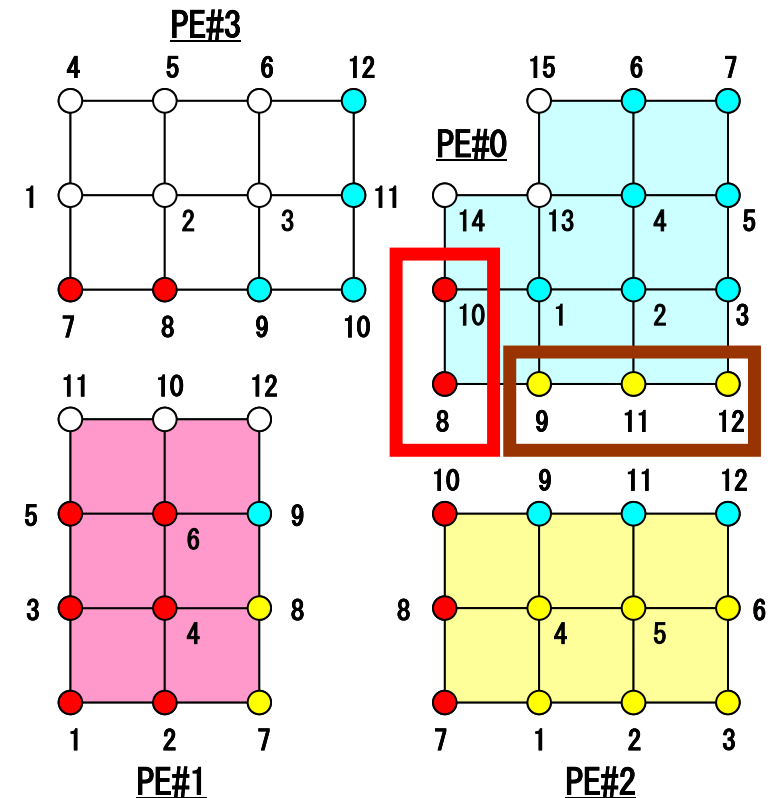
RECV: receiving to external nodes

Recv. continuous data to recv. buffer from neighbors

- MPI_Irecv**

(recvbuf, count, datatype, dest, tag, comm, request)

- **recvbuf** choice I starting address of receiving buffer
- **count** I I number of elements in receiving buffer
- **datatype** I I data type of elements of receiving buffer
- **source** I I rank of source



MPI_IRecv

- Begins a non-blocking receive
 - Receiving the contents of receiving buffer (starting from **recvbuf**, number of messages: **count**) from **source** with **tag** .
 - Contents of receiving buffer cannot be used before calling corresponding **MPI_Waitall**.

• **call MPI_IRecv**
(recvbuf, count, datatype, dest, tag, comm, request, ierr)

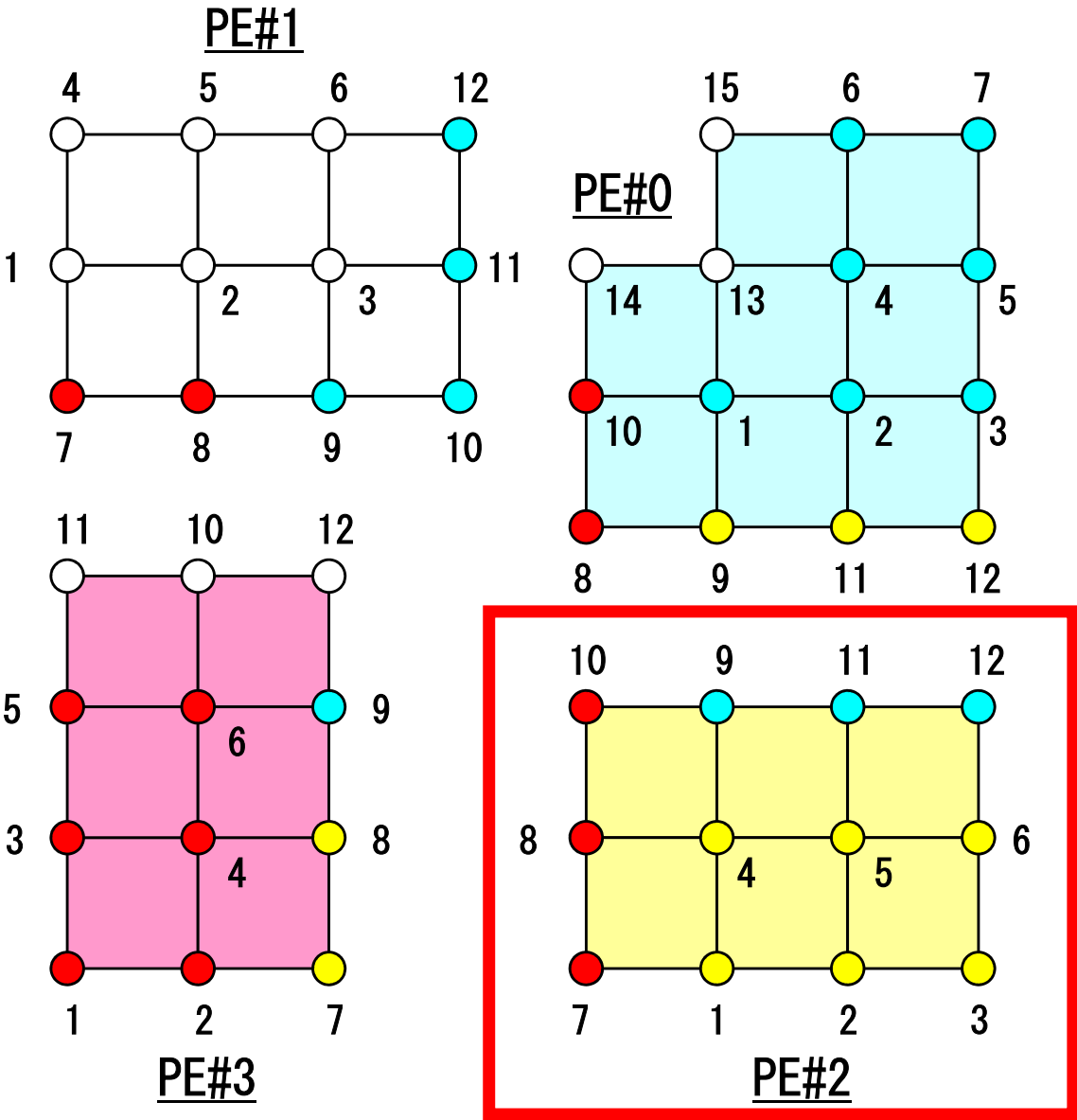
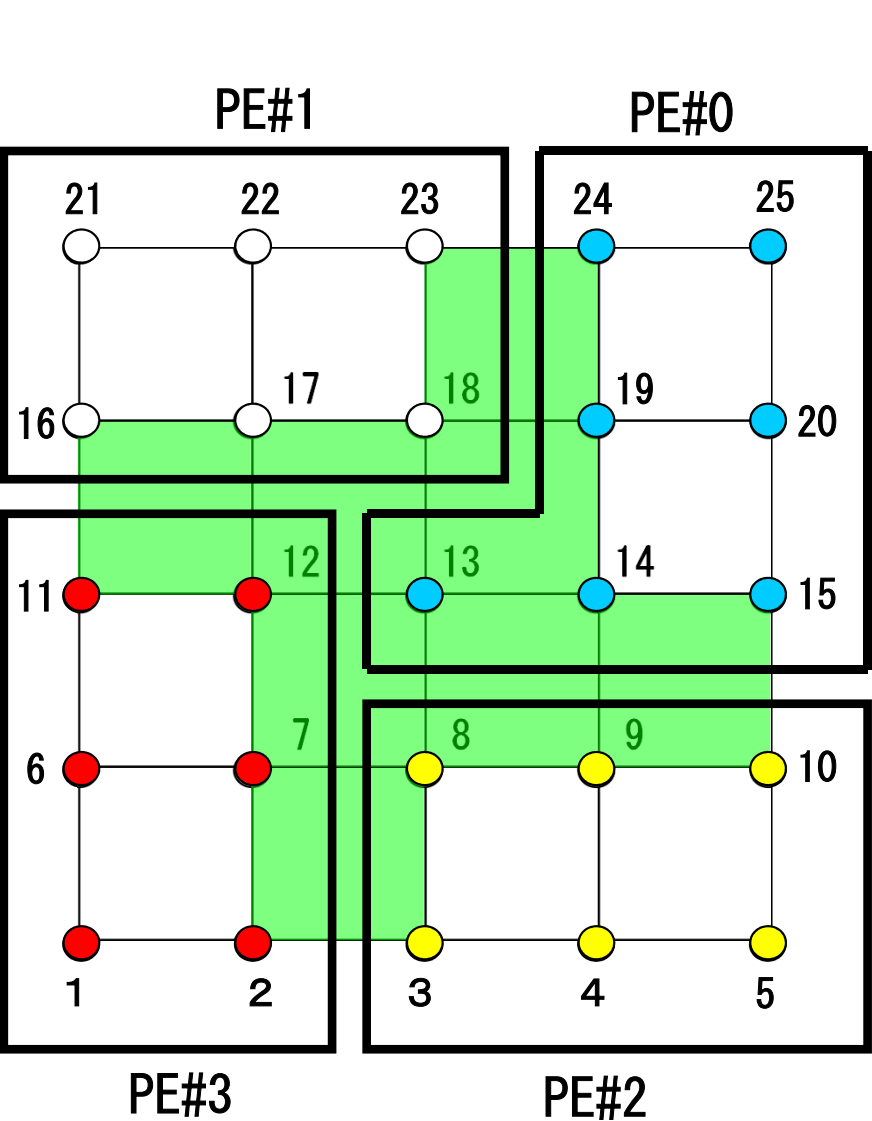
– <u>recvbuf</u>	choice	I	starting address of receiving buffer
– <u>count</u>	I	I	number of elements in receiving buffer
– <u>datatype</u>	I	I	data type of elements of receiving buffer
– <u>source</u>	I	I	rank of source
– <u>tag</u>	I	I	message tag
This integer can be used by the application to distinguish messages. Communication occurs if tag' s of MPI_Isend and MPI_Irecv are matched.			
Usually tag is set to be “0” (in this class),			
– <u>comm</u>	I	I	communicator
– <u>request</u>	I	O	communication request used in MPI_Waitall
– <u>ierr</u>	I	O	completion code

MPI_WAITALL

- **MPI_Waitall** blocks until all comm's, associated with request in the array, complete. It is used for synchronizing MPI_Isend and MPI_Irecv in this class.
- At sending phase, contents of sending buffer cannot be modified before calling corresponding **MPI_Waitall**. At receiving phase, contents of receiving buffer cannot be used before calling corresponding **MPI_Waitall**.
- MPI_Isend and MPI_Irecv can be synchronized simultaneously with a single **MPI_Waitall** if it is consistent.
 - Same request should be used in MPI_Isend and MPI_Irecv.
- Its operation is similar to that of **MPI_Barrier** but, **MPI_Waitall** can not be replaced by **MPI_Barrier**.
 - Possible troubles using **MPI_Barrier** instead of **MPI_Waitall**: Contents of **request** and **status** are not updated properly, very slow operations etc.
- **call MPI_WAITALL (count, request, status, ierr)**
 - count I I number of processes to be synchronized
 - request I I/O comm. request used in MPI_Waitall (array size: count)
 - status I O array of status objects
MPI_STATUS_SIZE: defined in 'mpif.h', 'mpi.h'
 - ierr I O completion code

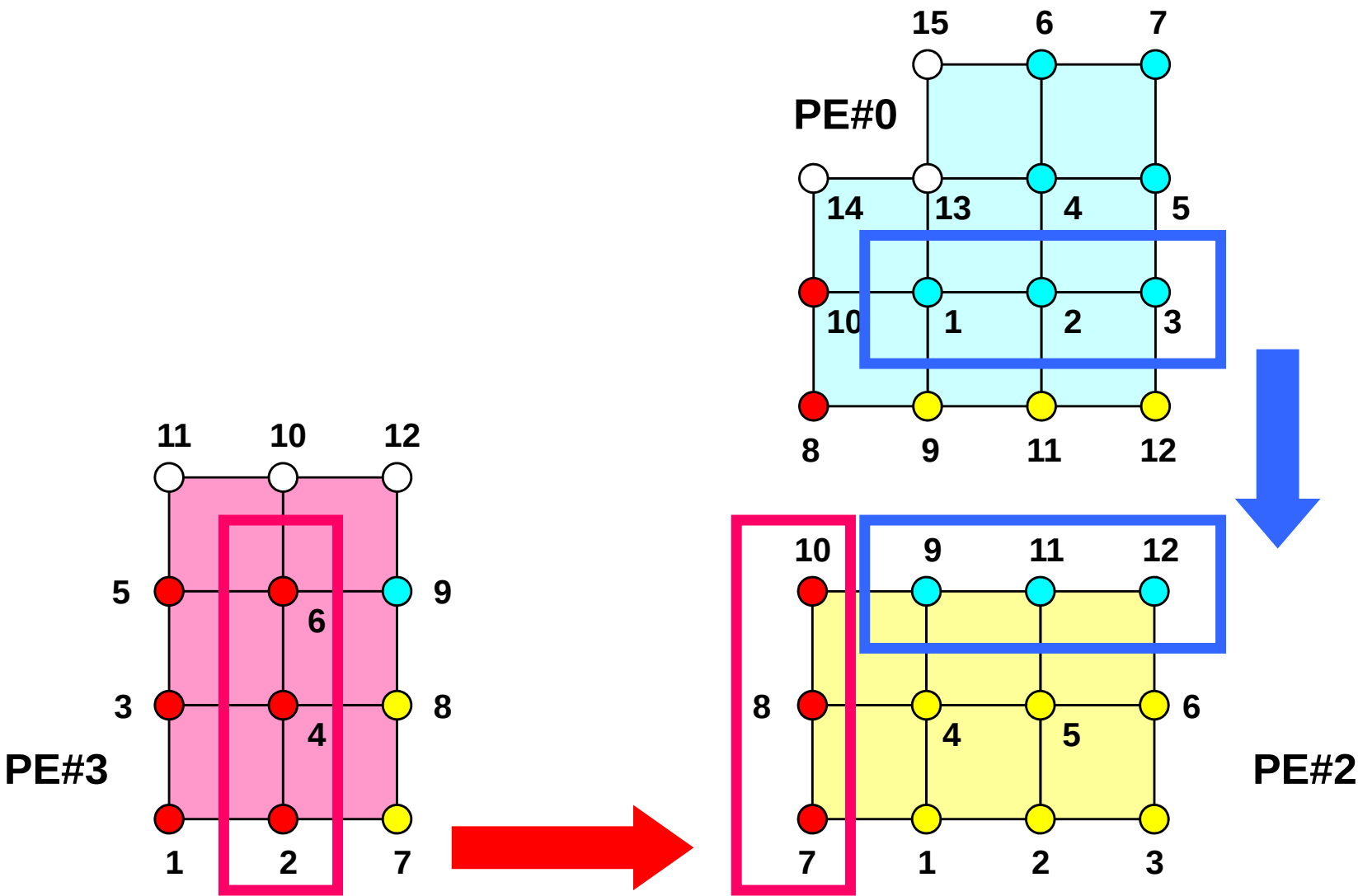
Node-based Partitioning

internal nodes - elements - external nodes



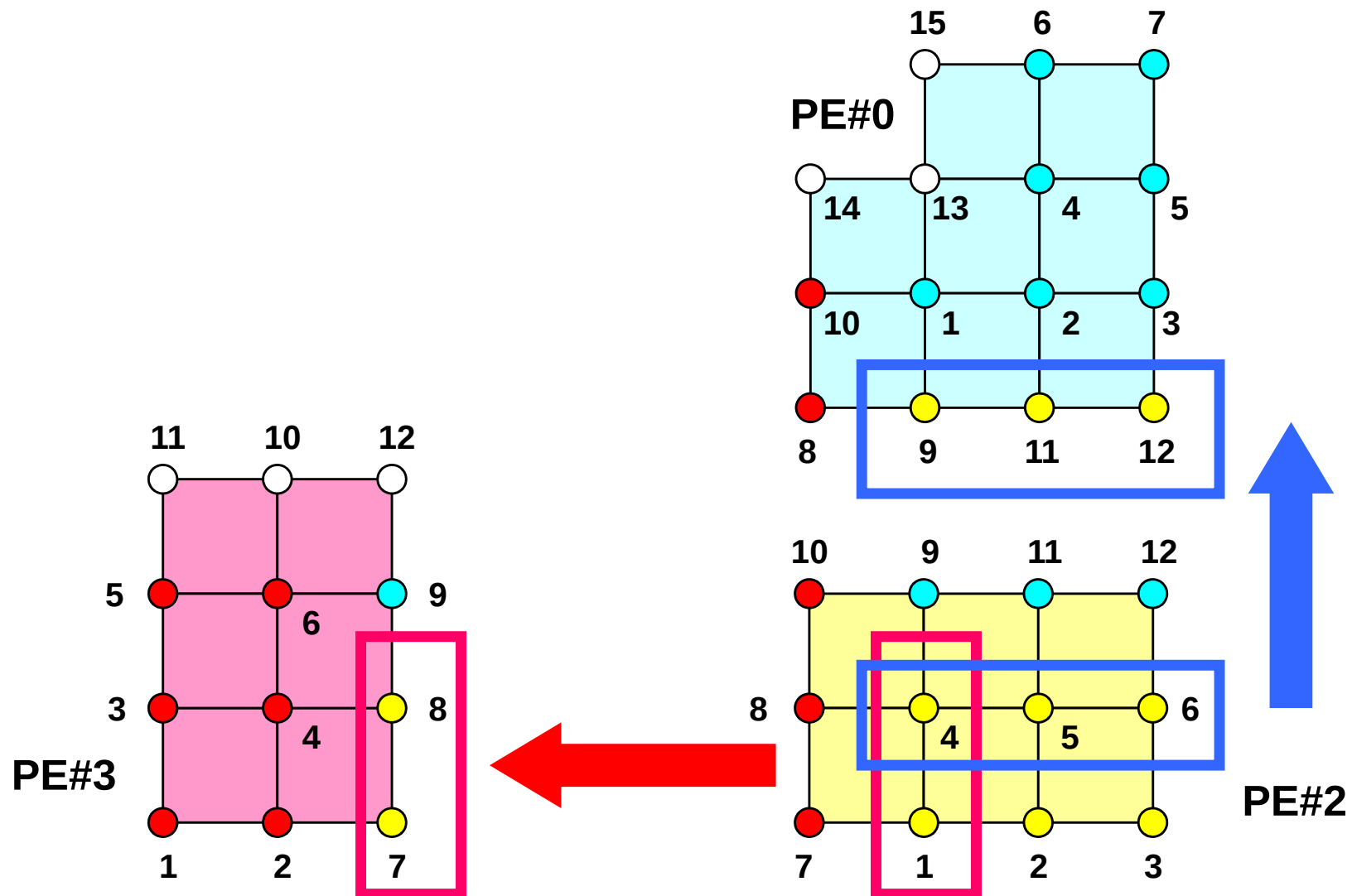
External Nodes (外点) : RECEIVE

PE#2 : receive information for “external nodes”



Boundary Nodes (境界点) : SEND

PE#2 : send information on “boundary nodes”



Distributed Local Data Structure for Parallel Computation

- Distributed local data structure for domain-to-domain communications has been introduced, which is appropriate for such applications with sparse coefficient matrices (e.g. FDM, FEM, FVM etc.).
 - SPMD
 - Local Numbering: Internal pts to External pts
 - Generalized communication table
- Everything is easy, if proper data structure is defined:
 - Values at boundary pts are copied into sending buffers
 - Send/Recv
 - Values at external pts are updated through receiving buffers